#### Research Article

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# Some results on the total proper k-connection number

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**Abstract:** In this paper, we first investigate the total proper connection number of a graph G according to some constraints of  $\overline{G}$ . Next, we investigate the total proper connection numbers of graph G with large clique number  $\omega(G) = n - s$  for  $1 \le s \le 3$ . Finally, we determine the total proper k-connection numbers of circular ladders, Möbius ladders and all small cubic graphs of order 8 or less.

**Keywords:** total coloring, total proper path, total proper *k*-connected, total proper *k*-connection number, complement graph, clique number

MSC 2020: 05C15, 05C35, 05C40

#### 1 Introduction

In this paper, all graphs under our consideration are simple, finite and undirected. We follow the notation and terminology of [1]. For a graph G, we denote by V(G), E(G) and  $\operatorname{diam}(G)$  the vertex set, edge set and diameter of G, respectively. The *distance* between two vertices u and v in a connected graph G, denoted by  $\operatorname{dist}(u,v)$ , is the length of a shortest path between them in G. The *eccentricity* of a vertex v in G is defined as  $\operatorname{ecc}_G(v) = \max_{x \in V(G)} \operatorname{dist}(v,x)$ . For convenience, a set of internally pairwise vertex disjoint paths will be called *disjoint*.

In recent years, colored notions of connectivity in graphs becomes a new and active subject in graph theory. Stating from rainbow connection [2], rainbow vertex connection [3] and total rainbow connection [4,5] appeared later. Many researchers are working in this field, and a lot of papers have been published in journals, see [6–16] for details. The reader can also see [17] for a survey, [18] for a dynamic survey and [19] for a new monograph on this topic.

In 2012, Borozan et al. [20] introduced the concept of proper k-connection number. A path in an edge-colored graph is a *proper path* if any two adjacent edges on the path differ in color. An edge-colored graph is *proper k-connected* if any two distinct vertices of the graph are connected by k disjoint proper paths. The proper k-connection number of a k-connected graph G, denoted by  $pc_k(G)$ , is defined as the smallest number of colors that are needed in order to make G proper k-connected. For more results, the reader can see [21–24] for details.

As a natural generalization, Jiang et al. [25] presented the concept of proper vertex k-connection number. A path in a vertex-colored graph is a *vertex proper path* if any two internal adjacent vertices of the path differ in color. A vertex-colored graph is *proper vertex k-connected* if any two distinct vertices of the graph are connected by k disjoint vertex proper paths. For a k-connected graph G, the proper vertex

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*k-connection number* of G, denoted by  $pvc_k(G)$ , is defined as the smallest number of colors required to make *G* proper vertex *k*-connected.

Motivated by the concept of total chromatic number of graph, now for proper connection and proper vertex connection, the concept of total proper connection was introduced by Jiang et al. [26]. A total coloring of a graph G is a mapping from the set  $V(G) \cup E(G)$  to some finite set of colors. A path in a total-colored graph is a total proper path if the coloring of the edges and internal vertices is proper, that is, any two adjacent or incident elements of edges and internal vertices on the path differ in color. A total-colored graph is total proper k-connected if any two distinct vertices of the graph are connected by k disjoint total proper paths. For a connected graph G, the total proper k-connection number of a k-connected graph G, denoted by  $\operatorname{tpc}_k(G)$ , is defined as the smallest number of colors that are needed in order to make G total proper k-connected. For convenience, we write tpc(G) for  $tpc_1(G)$ . Obviously,  $tpc(G) \le tpc_2(G) \le tpc_3(G)$ . By [26], if G is complete, then tpc(G) = 1; if G has a Hamiltonian path that is not complete, then tpc(G) = 3. Note that if G is a nontrivial connected graph and H is a connected spanning subgraph of G, then  $tpc(G) \le tpc(H)$ .

In this paper, we investigate the total proper connection number of a graph G under some constraints on its complement graph  $\overline{G}$ .

**Theorem 1.1.** Let G be a connected graph of order  $n \ge 3$ , if diam( $\overline{G}$ ) does not belong to the following two cases: (i) diam( $\overline{G}$ ) = 2, 3, (ii)  $\overline{G}$  contain exactly two components and one of them is trivial, then tpc(G)  $\leq 4$ .

For the remaining cases, tpc(G) can be very large as discussed in Section 2. Then we add a constraint, i.e., we let  $\overline{G}$  be triangle-free. Hence, G is claw-free, and we can derive our next main result:

**Theorem 1.2.** For a connected graph G, if  $\overline{G}$  is triangle-free, then tpc(G) = 3.

Recall that a clique of a graph is a set of mutually adjacent vertices, and that a maximum clique is a clique of the largest possible size in a given graph. The clique number  $\omega(G)$  of a graph G is the number of vertices in a maximum clique in *G*. Let *G* be a connected graph, and let *X* be a maximum clique of *G*. We say that  $N_X(u)$  is the set of neighbors of u in G[X] and  $d_X(u) = |N_X(u)|$ . Let  $F = G[V(G) \setminus X]$ . Kemnitz and Schiermeyer [9] considered graphs with rc(G) = 2 and large clique number. In this paper, we characterize graphs with small total proper connection number with respect to their large clique number. If  $\omega(G) = n$ , then G is a complete graph, which implies  $\operatorname{tpc}(G) = 1$ . If G is connected and  $\omega(G) = n - 1$ , then G has a Hamiltonian path, and so tpc(G) = 3. For the cases  $\omega(G) = n - 2$ , n - 3, we obtain the following three main results.

**Theorem 1.3.** Let G be a connected graph of order n. If  $\omega(G) = n - 2$  and X is a maximum clique of G with  $V(G)\backslash X = \{u_1, u_2\}, \text{ then } \mathrm{tpc}(G) = 3.$ 

**Theorem 1.4.** Let G be a connected graph of order n, diam(G) = 2. If  $\omega(G) = n - 3$  and X is a maximum clique of G with  $V(G)\backslash X=\{u_1,u_2,u_3\}$ , then tpc(G)=3 or tpc(G)=4 for the following case  $F\cong 3K_3$ ,  $|N_X(u_1) \cap N_X(u_2) \cap N_X(u_3)| = 1$  and  $d_X(u_1) = d_X(u_2) = d_X(u_3) = 1$ .

**Theorem 1.5.** Let G be a connected graph of order n,  $diam(G) \ge 3$ . If  $\omega(G) = n - 3$  and X is a maximum clique of G with  $V(G)\backslash X = \{u_1, u_2, u_3\}$ , then tpc(G) = 3, or tpc(G) = 4 and one of the following holds.

- (i)  $F \cong 3K_3$ ,  $N_X(u) \cap N_X(v) \neq \emptyset$  and  $d_X(u) = d_X(v) = 1$ , where u and v are any two distinct vertices in  $V(G)\backslash X$ .
- (ii)  $F \cong 3K_3$ ,  $d_X(u_1) = d_X(u_2) = d_X(u_3) = 1$  and for any two vertices in  $V(G)\backslash X$ , there is no common neighbor in G[X].

For an integer  $n \ge 3$ , the circular ladder  $CL_{2n}$  of order 2n is a cubic graph constructed by taking two copies of the cycle  $C_n$  on disjoint vertex sets  $(u_1, u_2, \dots, u_n)$  and  $(v_1, v_2, \dots, v_n)$ , then joining the corresponding vertices  $u_i v_i$  for  $1 \le i \le n$ . The Möbius ladder  $M_{2n}$  of order 2n is obtained from the ladder by deleting the edges  $u_1u_n$  and  $v_1v_n$ , and then inserting two edges  $u_1v_n$  and  $u_nv_1$ . Subscripts are considered modulo n, and we can derive our next main result:

**Theorem 1.6.** Let n be an integer with  $n \ge 3$ . Then

- (i)  $tpc(CL_{2n}) = tpc_2(CL_{2n}) = 3$ ,  $tpc_3(CL_{2n}) = 4$ .
- (ii)  $tpc(M_{2n}) = tpc_2(M_{2n}) = 3$ ,  $tpc_3(M_{2n}) = 4$ .

In [7], Fujie-Okamoto et al. investigated the rainbow k-connection numbers of all small cubic graphs of order 8 or less. In this paper, we determine the total proper k-connection numbers of all small cubic graphs of order 8 or less. We can easily verify that all such cubic graphs have orders 4, 6, or 8, and those with orders 4 or 6 are  $K_4$ ,  $K_{3,3}$ , and  $K_3 \square K_2$  (where  $\square$  denotes Cartesian product). In [27], it was shown that all connected cubic graphs of order 8 are  $Q_3$ ,  $M_8$ ,  $F_1$ ,  $F_2$ , and  $F_3$ , and these graphs are depicted in Figure 1. Our last main result is stated as follows:

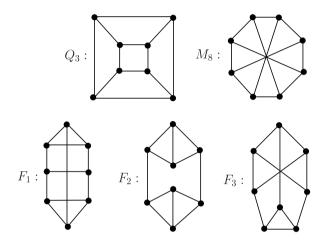


Figure 1: All connected cubic graphs of order 8.

#### Theorem 1.7.

- (i)  $tpc(K_4) = 1$ ,  $tpc_2(K_4) = 3$ ,  $tpc_3(K_4) = 4$ .
- (ii)  $tpc(K_{3,3}) = tpc_2(K_{3,3}) = 3$ ,  $tpc_3(K_{3,3}) = 4$ .
- (iii)  $tpc(K_3 \square K_2) = tpc_2(K_3 \square K_2) = 3$ ,  $tpc_3(K_3 \square K_2) = 4$ .
- (iv)  $tpc(Q_3) = tpc_2(Q_3) = 3$ ,  $tpc_3(Q_3) = 4$ .
- (v)  $tpc(M_8) = tpc_2(M_8) = 3$ ,  $tpc_3(M_8) = 4$ .
- (vi)  $tpc(F_1) = tpc_2(F_1) = 3$ ,  $tpc_3(F_1) = 4$ .
- (*vii*)  $tpc(F_2) = tpc_2(F_2) = 3$ .
- (*viii*)  $tpc(F_3) = tpc_2(F_3) = 3$ ,  $tpc_3(F_3) = 4$ .

#### 2 Proof of Theorem 1.1

In order to prove Theorem 1.1, we need the following lemmas.

**Lemma 2.1.** [26] *For*  $2 \le m \le n$ , *we have*  $tpc(K_{m,n}) = 3$ .

**Lemma 2.2.** [26] If G is a complete multipartite graph that is neither a complete graph nor a tree, then tpc(G) = 3.

Let  $N_{\overline{G}}^i(x) = \{v : \operatorname{dist}_{\overline{G}}(x, v) = i\}$ , where  $0 \le i \le 3$ , and  $N_{\overline{G}}^4(x) = \{v : \operatorname{dist}(x, v) \ge 4\}$ . In this paper, we use  $N_{\overline{G}}^i$  instead of  $N_{\overline{G}}^i(x)$  for convenience. Then  $N_{\overline{G}}^0 = \{x\}$  and  $N_{\overline{G}}^1 = N_{\overline{G}}(x)$ .

**Lemma 2.3.** For a connected graph G, if  $\overline{G}$  is connected and  $\operatorname{diam}(\overline{G}) \geq 4$ , then  $\operatorname{tpc}(G) \leq 4$ .

**Proof.** Choose a vertex x with  $\operatorname{ecc}_{\overline{G}}(x) = \operatorname{diam}(\overline{G})$ . By the definition of  $N_{\overline{G}}^i$ , we know  $uv \in E(G)$  for any  $u \in N_{\overline{G}}^i$ ,  $v \in N_{\overline{G}}^j$  with  $|i-j| \ge 2$ . Now we define a total coloring of G as follows: assign color 1 to the edges xu for  $u \in N_{\overline{G}}^2$ , all edges between  $N_{\overline{G}}^1$  and  $N_{\overline{G}}^3$ , and all vertices and edges in  $N_{\overline{G}}^4$ ; assign color 2 to the edges xu for  $u \in N_{\overline{G}}^3$ , all edges between  $N_{\overline{G}}^2$  and  $N_{\overline{G}}^4$ , and all vertices and edges in  $N_{\overline{G}}^1$ ; assign color 3 to the edges xu for  $u \in N_{\overline{G}}^4$ , all edges between  $N_{\overline{G}}^1$  and  $N_{\overline{G}}^4$ , and all vertices and edges in  $N_{\overline{G}}^2$ ,  $N_{\overline{G}}^3$ ; assign color 4 to the vertex x.

We prove that there is a total proper path between any two vertices u and v of G. It is trivial when  $uv \in E(G)$ . Thus, we only need to consider the pairs  $u, v \in N_{\overline{G}}^i$  or  $u \in N_{\overline{G}}^i$ ,  $v \in N_{\overline{G}}^{i+1}$ . Note that  $P = xx_3x_1x_4x_2$  is a total proper path, where  $x_i \in N_{\overline{G}}^i$ . By means of the path P, we can find that u and v are connected by some total proper path for any  $u \in N_{\overline{G}}^{i-1}$ . If i = 1, then  $P = ux_3x_2x_4v$  is a total proper path, where  $x_i \in N_{\overline{G}}^i$ . If i = 2, then  $P = ux_4x_4v$  is a total proper path, where  $x_4 \in N_{\overline{G}}^4$ . If i = 3, then  $P = ux_4x_4x_2v$  is a total proper path, where  $x_i \in N_{\overline{G}}^i$ . If i = 4, then  $P = ux_2v$  is a total proper path, where  $x_2 \in N_{\overline{G}}^2$ . Hence,  $\operatorname{tpc}(G) \leq 4$ .

**Proof of Theorem 1.1.** Assume that  $\overline{G}$  is connected. Since  $\operatorname{diam}(\overline{G}) \geq 4$ , we have  $\operatorname{tpc}(G) \leq 4$  by Lemma 2.3. Assume that  $\overline{G}$  is disconnected. By the assumption, we know that there exist either at least three connected components or exactly two nontrivial components. Let  $\overline{G_i}$  be the components of  $\overline{G}$  with  $t_i = |V(\overline{G_i})|$ , where  $1 \leq i \leq h$ . Then G contains a connected spanning subgraph  $K_{t_1,t_2,\ldots,t_h}$ , and we have  $\operatorname{tpc}(G) \leq \operatorname{tpc}(K_{t_1,t_2,\ldots,t_h}) = 3$  by Lemma 2.2. Note that G is not complete. Thus,  $\operatorname{tpc}(G) \geq 3$ , and so  $\operatorname{tpc}(G) = 3$ .

Next, we will give three examples to show that  $\operatorname{tpc}(G)$  can be arbitrarily large if one of the following three conditions holds:  $\operatorname{diam}(\overline{G}) = 2$ ,  $\operatorname{diam}(\overline{G}) = 3$ ,  $\overline{G}$  contains exactly two connected components and one of them is trivial.

**Example 2.4.** For the graph  $\overline{G}$  shown in Figure 2, we choose a vertex x with  $\operatorname{ecc}_{\overline{G}}(x) = \operatorname{diam}(\overline{G})$ . Let  $N_{\overline{G}}^1(x) = \{u_i | 1 \le i \le k\}$ ,  $N_{\overline{G}}^2(x) = \{v_j | 1 \le j \le k\}$ , and let  $E(\overline{G}) = \{xu_i | 1 \le i \le k\} \cup \{u_i u_{i_2} | 1 \le i_1, i_2 \le k\} \cup \{v_{j_1} v_{j_2} | 1 \le j_1, j_2 \le k\}$   $\cup \{u_i v_j | 1 \le i, j \le k\} \setminus \{u_i v_i | 1 \le i \le k\}$ , where  $k \ge 3$ . Obviously,  $\operatorname{diam}(\overline{G}) = 2$  and G is a tree. Then  $\operatorname{tpc}(G) = \Delta(G) + 1 = k + 1$  by [26, Theorem 1]. Observe that  $\operatorname{tpc}(G)$  will be arbitrarily large based on the increase of k.

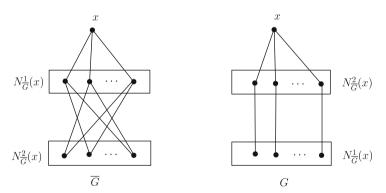


Figure 2: The graph of Example 2.4.

**Example 2.5.** For the graph  $\overline{G}$  shown in Figure 3, we choose a vertex x with  $\operatorname{ecc}_{\overline{G}}(x) = \operatorname{diam}(\overline{G})$ . Let  $N_{\overline{G}}^1(x) = \{u_i | 1 \le i \le k\}$ ,  $N_{\overline{G}}^2(x) = \{v_j | 1 \le j \le k\}$ , and  $N_{\overline{G}}^3(x) = \{w_s | 1 \le s \le k\}$ , where  $k \ge 3$ . Furthermore, let  $E(\overline{G}) = \{xu_i | 1 \le i \le k\} \cup \{u_iv_j | 1 \le i, j \le k\} \cup \{v_jw_s | 1 \le j, s \le k\} \cup \{v_jv_j | 1 \le j_1, j_2 \le k\}$ . Obviously,  $\operatorname{diam}(\overline{G}) = 3$  and G is a connected graph. Note that  $N_{\overline{G}}^2(x)$  is a stable set in  $\overline{G}$ , and each edge between x and  $N_{\overline{G}}^2(x)$  is a cut edge in G. Therefore,  $\operatorname{tpc}(G) \ge k + 1$  by [26, Proposition 2], and so  $\operatorname{tpc}(G)$  will be arbitrarily large based on the increase of k.

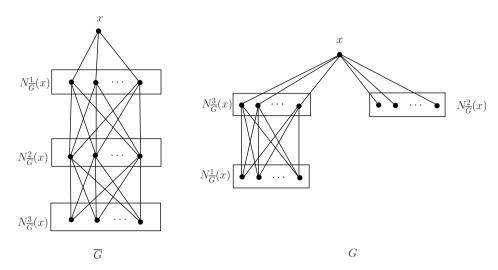


Figure 3: The graph of Example 2.5.

**Example 2.6.** Let  $\overline{G}$  contains exactly two components  $\overline{G}_1$  and  $\overline{G}_2$ , where  $\overline{G}_1$  is trivial and  $\overline{G}_2$  is a clique of  $\overline{G}$ . Clearly, G is a star, and  $\operatorname{tpc}(G) = |V(\overline{G}_2)| + 1$ . Thus,  $\operatorname{tpc}(G)$  can be made arbitrarily large by increasing  $|V(\overline{G}_2)|$ .

# 3 Proof of Theorem 1.2

**Lemma 3.1.** For a connected graph G, if  $\overline{G}$  is connected with  $diam(\overline{G}) \ge 4$ , and  $\overline{G}$  is triangle-free, then tpc(G) = 3.

**Proof.** Choose a vertex x with  $\operatorname{ecc}_{\overline{G}}(x) = \operatorname{diam}(\overline{G})$ . Since  $\overline{G}$  is triangle-free, we know that  $N_{\overline{G}}^1$  is a clique in G. Now we define a total coloring of G as follows: assign color 1 to the edges XU for  $U \in N_{\overline{G}}^2$ , all edges between  $N_{\overline{G}}^1$  and  $N_{\overline{G}}^3$ , and all vertices and edges in  $N_{\overline{G}}^4$ ; assign color 2 to the edges between  $N_{\overline{G}}^2$  and  $N_{\overline{G}}^4$ , all vertices and edges in  $N_{\overline{G}}^1$ , and the vertex X; assign color 3 to the edges XU for  $U \in N_{\overline{G}}^3$ ,  $N_{\overline{G}}^4$ , all edges between  $N_{\overline{G}}^1$  and  $N_{\overline{G}}^4$ , and all vertices and edges in  $N_{\overline{G}}^2$ ,  $N_{\overline{G}}^3$ .

We prove that there is a total proper path between any two distinct vertices u and v in G. Note that  $P = xx_2x_4x_1x_3$  is a total proper path, where  $x_i \in N_{\overline{G}}^i$ . By means of the path P, we can find that u and v are connected by some total proper path for any  $u \in N_{\overline{G}}^i$ ,  $v \in N_{\overline{G}}^{i+1}$ . Thus, we only need to consider the pairs  $u, v \in N_{\overline{G}}^i$ . For i = 2,  $P = uxx_4v$  is a total proper path, where  $x_4 \in N_{\overline{G}}^4$ . For i = 4,  $P = uxx_2v$  is a total proper path, where  $x_2 \in N_{\overline{G}}^i$ . For i = 3,  $P = ux_1x_4x_2xv$  is a total proper path, where  $x_i \in N_{\overline{G}}^i$ . Thus, G is total proper connected with the above coloring, and so  $\operatorname{tpc}(G) = 3$ .

**Lemma 3.2.** Let G be a connected graph. If  $diam(\overline{G}) = 3$  and  $\overline{G}$  is triangle-free, then tpc(G) = 3.

**Proof.** For a vertex x of  $\overline{G}$  satisfying  $\operatorname{ecc}_{\overline{G}}(x) = \operatorname{diam}(\overline{G}) = 3$ , let  $n_i$  represent the number of vertices with distance i from x. If  $n_1 = n_2 = n_3 = 1$ , then  $G \cong P_4$ , and so  $\operatorname{tpc}(G) = 3$ .

**Case 1.** Two of  $n_1$ ,  $n_2$ ,  $n_3$  are equal to 1. Without loss of generality, we may assume  $n_1 = n_2 = 1$ . Since  $\overline{G}$  is triangle-free, we have that  $N_G^3$  is a stable set in  $\overline{G}$ , and so a clique in G. We can find that G has a Hamiltonian path. Thus,  $\operatorname{tpc}(G) = 3$ .

**Case 2**. One of  $n_1$ ,  $n_2$ ,  $n_3$  is equal to 1. Suppose  $n_2 = 1$ . Since  $\overline{G}$  is triangle-free, we know that  $N_{\overline{G}}^1$  and  $N_{\overline{G}}^3$  is a stable set in  $\overline{G}$ , and so a clique in G. Note that G has a Hamiltonian path, and so  $\operatorname{tpc}(G) = 3$ .

**Subcase 2.1.**  $n_1 = 1$ . Since  $\overline{G}$  is triangle-free, we obtain that  $N_{\overline{G}}^2$  is a clique in G. Define a total coloring of G as follows: assign color 3 to the vertex X, all edges between  $N_{\overline{G}}^2$  and  $N_{\overline{G}}^3$ , and all edges between  $N_{\overline{G}}^1$  and  $N_{\overline{G}}^3$ ; assign color 2 to the edges xu for  $u \in N_{\overline{G}}^2$ , and all vertices and edges in  $N_{\overline{G}}^3$ ; assign color 1 to the edges xu for  $u \in N_{\overline{G}}^3$ , and all vertices and edges in  $N_{\overline{G}}^1$ . We prove that there is a total proper path between any two distinct vertices u and v in G. Note that  $P = x_1x_3xx_2$  is a total proper path, where  $x_i \in N_{\overline{G}}^1$ . By means of the path P, we know that u and v are connected by some total proper path for any  $u \in N_{\overline{G}}^1$ . For any two vertices u,  $v \in N_{\overline{G}}^3$ , it is trivial if  $uv \in E(G)$ . If  $uv \notin E(G)$ , since u,  $v \in N_{\overline{G}}^3$ , there exist two vertices u',  $v' \in N_{\overline{G}}^3$  such that uu',  $vv' \in E(\overline{G})$ . Since  $\overline{G}$  is triangle-free, we can see that  $u' \neq v'$  and vu',  $uv' \in E(G)$ . Then P = uxu'v is a total proper path. Hence, G is total proper connected with the above coloring, and so tpc(G) = 3.

**Subcase 2.2.**  $n_3 = 1$ . Since  $\overline{G}$  is triangle-free, we know that  $N_G^1$  is a stable set in  $\overline{G}$ , and so a clique in G. Define a total coloring of G as follows: assign color 3 to the vertex X, and all edges between  $N_G^1$  and  $N_G^3$ ; assign color 2 to the edges X for X for X for X and all vertices and edges in X assign color 1 to the edges X for X for X and all vertices and edges in X and all vertices and edges in X assign color 1 to the edges X for X for X and all vertices X and all vertices and edges in X for X and all vertices X and all vertices and edges in X for X for X for X and X and X is a total proper path, where X is a total proper path between any two distinct vertices X and X are connected by some total proper path for any X is an end X for X for X is a clique in X and so X for X for X is a clique in X and so X for X for X is a clique in X for X

**Case 3.**  $n_1$ ,  $n_2$ ,  $n_3 \ge 2$ . Since  $\overline{G}$  is triangle-free, we have that  $N_{\overline{G}}^1$  is a stable set in  $\overline{G}$ , and so a clique in G. If any vertex in  $N_{\overline{G}}^3$  is adjacent to all vertices of  $N_{\overline{G}}^2$  in  $\overline{G}$ , then both  $N_{\overline{G}}^2$  and  $N_{\overline{G}}^3$  are stable sets in  $\overline{G}$ , and so cliques in G. Thus, G has a Hamiltonian path, and so  $\operatorname{tpc}(G) = 3$ .

Otherwise, we choose a vertex  $u \in N_{\overline{G}}^3$ , let  $V_u$  denote the set of neighbors of u in  $N_{\overline{G}}^2$  in G, we have  $V_u \neq \emptyset$ ,  $N_{\overline{G}}^2$ . Define a total coloring of G: assign color 2 to the vertex x, all vertices and edges in  $N_{\overline{G}}^1$ , and all edges between  $N_{\overline{G}}^2$ ,  $N_{\overline{G}}^3$ ; assign color 3 to the vertex u, all edges between  $N_{\overline{G}}^1$  and  $N_{\overline{G}}^3 \setminus \{u\}$ , and all edges between x and  $V_u$ ; assign color 2 to the remaining vertices and edges. Note that  $P = xvux_1$  is a total proper path, where  $v \in V_u$ ,  $x_1 \in N_{\overline{G}}^1$ . For any two vertices w,  $z \in N_{\overline{G}}^2$ ,  $P = wx_1uvxz$  is a total proper path, where  $x_1 \in N_{\overline{G}}^1$ ,  $v \in V_u$ . For any two vertices  $x_1 \in N_{\overline{G}}^2$ ,  $y \in V_u$ . For any two vertices  $x_1 \in N_{\overline{G}}^2$ ,  $y \in V_u$ . If  $|N_{\overline{G}}^2 \setminus V_y| = 1$ , then |P| = uxv| is a total proper path, where  $|u| \in V_y$ ,  $|V| \in N_{\overline{G}}^2 \setminus V_y| = 1$ , then |P| = uxv| is a total proper path, where  $|u| \in V_y$ ,  $|V| \in N_{\overline{G}}^2 \setminus V_y| = 1$ , then |P| = uxv| is a total proper path, where  $|V| \in V_y$ ,  $|V| \in N_{\overline{G}}^2 \setminus V_y| = 1$ , then  $|V| \in V_y$ ,  $|V| \in$ 

**Lemma 3.3.** For a connected graph G, if  $\overline{G}$  is triangle-free and  $\operatorname{diam}(\overline{G}) = 2$ , then  $\operatorname{tpc}(G) = 3$ .

**Proof.** Choose a vertex x with  $\operatorname{ecc}_{\overline{G}}(x) = \operatorname{diam}(\overline{G}) = 2$ . Since G is connected, we have  $n_1 \geq 2$ ,  $n_2 = 1$  or  $n_1, n_2 \geq 2$ , and there exist two vertices  $u \in N_{\overline{G}}^1, v \in N_{\overline{G}}^2$  such that  $uv \in E(G)$ . Assume  $n_1 \geq 2$  and  $n_2 = 1$ . Since  $\overline{G}$  is triangle-free, we know that  $N_{\overline{G}}^1$  is a stable set in  $\overline{G}$ , and so a clique in G. Note that G has a Hamiltonian path, and so  $\operatorname{tpc}(G) = 3$ .

Assume  $n_1$ ,  $n_2 \ge 2$ . Observe that  $N_{\overline{G}}^1$  is a stable set in  $\overline{G}$  since  $\overline{G}$  is triangle-free, and so a clique in G. We show a total coloring of G as follows: assign color 1 to the vertex X, the edge UV and all vertices in  $N_{\overline{G}}^1 \setminus U$ ; assign color 3 to the vertex V and all edges in  $N_{\overline{G}}^1$ ; assign color 2 to the remaining vertices and edges. If there exist some vertices  $W \in N_{\overline{G}}^2$  with  $d_{\overline{G}}(W) = n - 2$ , then W is adjacent to the remaining vertices except X in  $\overline{G}$ .

Since diam( $\overline{G}$ ) = 2, there exists an edge  $w_1w_2 \in E(\overline{G})$  with  $w_1 \in N_{\overline{G}}^1$ ,  $w_2 \in N_{\overline{G}}^2$ . Thus,  $w_1, w_2$  is a triangle in  $\overline{G}$ , a contradiction. Hence,  $d_{\overline{G}}(w) < n-2$  for all  $w \in N_{\overline{G}}^2$ , and so  $d_G(w) \ge 2$ . For any  $z \in N_{\overline{G}}^1$ , we know that P = xvuz is a total proper path. For any  $y \in N_G^2 \setminus \{v\}$  and  $z \in N_G^1$ , if  $N_G(y) \cap N_G^1 \neq \emptyset$ , let  $w \in N_G(y) \cap N_G^1$ . Then *ywz* is a total proper path. Otherwise, let  $N_G(y) \cap N_{\overline{G}}^1 = \emptyset$ . We claim that y is adjacent to all the other vertices of  $N_{\overline{G}}^2$  in G. In fact, for any vertex  $w \in N_{\overline{G}}^2 \setminus \{y\}$ , there exists a vertex  $w' \in N_{\overline{G}}^1$  such that  $ww' \in E(\overline{G})$ . Since  $yw' \in E(\overline{G})$ , we know that  $yw \in E(G)$ . Then yvuz is a total proper path. Next we consider  $w, z \in N_{\overline{G}}^2$  such that  $wz \notin E(G)$ . Since  $\overline{G}$  is triangle-free, we have that G is claw-free, and at least one of w and z is adjacent to the v, without loss of generality, assume that  $wv \in E(G)$ . Since  $w, z \in N_{\mathbb{Z}}^2$ , there exist two vertices  $w', z' \in N_{\mathbb{Z}}^2$ such that  $ww', zz' \in E(\overline{G})$ , and  $w' \neq z'$ . Then  $zw', wz' \in E(G)$  and P = wvuw'z is a total proper path. Thus, *G* is total proper connected with the above coloring. Hence, tpc(G) = 3.

**Lemma 3.4.** Let G be a connected graph of order  $n \ge 3$ . If  $\overline{G}$  is disconnected and triangle-free, then tpc(G) = 3.

**Proof.** Suppose  $\overline{G}$  is triangle-free and contains two connected components one of which is trivial. Let  $\overline{G}_1$  and  $\overline{G}_2$  be the two components of  $\overline{G}$ , where  $V(\overline{G}_1) = \{u\}$ . Then u is adjacent to any other vertex in G. We will consider two cases according to the value of  $\delta$ , where  $\delta$  is the minimum degree of G. If  $\delta = 1$ , let  $d(v) = \delta$ . Since  $\overline{G}$  is triangle-free, we know that G is claw-free, and the subgraph  $G[V(G)\setminus\{v\}]$  is a complete graph. Thus, G has a Hamiltonian path, and so  $\operatorname{tpc}(G) = 3$ . If  $\delta \geq 2$ , let  $d(v) = \delta$ ,  $D = V(G) \setminus \{u, v\}$ , and  $V_v$  be the set of neighbors of v in G. Now we define a total coloring of G as follows: assign color 1 to the vertex v and all the edges between u and  $V_{\nu}$ ; assign color 3 to the vertex u and all the edges between  $\nu$  and  $V_{\nu}$ ; assign color 2 to the remaining vertices and edges. Since *G* is claw-free, we can find that the subgraph  $G[V(G)\setminus\{v\}\cup V_v]$  is a complete graph, and  $P = v_1 u v v_2$  is a total proper path, where  $v_1, v_2 \in V_v$ . For any  $w \in V_v, z \in D \setminus V_v$ , we obtain that P = wuz is a total proper path. Thus, G is total proper connected with the above coloring, and so tpc(G) = 3. Suppose  $\overline{G}$  contains at least three connected components or exactly two nontrivial components. Then we have tpc(G) = 3 by the similar proof of Theorem 1.1.

**Proof of Theorem 1.2.** If  $\overline{G}$  is connected, the result holds for the case diam( $\overline{G}$ )  $\geq 4$  by Lemma 3.1, the case  $\operatorname{diam}(\overline{G}) = 3$  by Lemma 3.2, and the case  $\operatorname{diam}(\overline{G}) = 2$  by Lemma 3.3. If  $\overline{G}$  is disconnected, the result holds by Lemma 3.4.

#### 4 Proof of Theorem 1.3

Suppose  $F \cong K_2$ . Note that *G* has a Hamiltonian path, and thus tpc(G) = 3. Next, we compute the total proper connection number of *G* by proving the following claim.

**Claim 1.** Let G be a graph obtained by adding two pendant vertices  $\{u_1, u_2\}$  to a vertex  $v_1$  of a complete graph  $K_t$ . Then tpc(G) = 3.

**Proof.** Since *G* is not a complete graph, we have  $tpc(G) \ge 3$ . Now we only need to prove  $tpc(G) \le 3$  by the following cases.

**Case 1.**  $t \equiv 0 \pmod{3}$ . Assign a total coloring c to G as follows: Let  $c(u_1v_1) = 1$ ,  $c(u_2v_1) = 3$ ;  $c(v_{3i+1}) = 1$  $c(v_{3i+2}v_{3i+3}) = 2$ ,  $c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = 1$ , and  $c(v_{3i+3}) = c(v_{3i+1}u_{3i+2}) = 3$ , where  $0 \le i \le \frac{t}{3} - 1$ . Observe that  $P_1 = u_1 v_1 v_2 \cdots v_{t-1} v_t$  and  $P_2 = u_2 v_1 v_t v_{t-1} \cdots v_3 v_2$  are two total proper paths.

**Case 2.**  $t \equiv 1 \pmod{3}$ . Assign a total coloring c to G as follows: Let  $c(u_1v_1) = c(v_{t-1}v_1) = 1$ ,  $c(v_t) = 2$ ,  $c(u_2v_1) = 1$  $c(v_i v_1) = 3$ . Let *i* be an integer with  $0 \le i \le \left| \frac{t}{3} \right| - 1$ ,  $c(v_{3i+1}) = c(v_{3i+2} v_{3i+3}) = 2$ ,  $c(v_{3i+2}) = c(v_{3i+3} v_{3i+4}) = 1$ , and  $c(v_{3i+3}) = c(v_{3i+1}u_{3i+2}) = 3$ . We can find that  $P_1 = u_1v_1v_2 \cdots v_{t-1}v_t$  and  $P_2 = u_2v_1v_{t-1} \cdots v_3v_2v_t$  are two total proper paths.

**Case 3.**  $t \equiv 2 \pmod{3}$ . Assign a total coloring c to G as follows: Let  $c(u_1v_1) = c(v_tv_1) = c(v_t) = c(v_{t-2}v_1) = 1$ ,  $c(v_{t-1}) = 2$ ,  $c(u_2v_1) = c(v_tv_{t-1}) = c(v_{t-1}v_2) = 3$ . Let i be an integer with  $0 \le i \le \left\lfloor \frac{t}{3} \right\rfloor - 1$ ,  $c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = 2$ ,  $c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = 1$ , and  $c(v_{3i+3}) = c(v_{3i+1}u_{3i+2}) = 3$ . We can easily verify that  $P_1 = u_1v_1v_2 \cdots v_{t-1}v_t$ ,  $P_2 = u_2v_1v_{t-2}v_{t-3} \cdots v_3v_2v_{t-1}$ , and  $P_3 = v_tv_1u_2$  are three total proper paths.

Thus, *G* is total proper connected with the above coloring, and so tpc(G) = 3. This completes the proof of Claim 1.

Suppose  $F \cong 2K_1$ . Assume that  $N_X(u_1) \cap N_X(u_2) = \emptyset$ . Then G has a Hamiltonian path, and so  $\operatorname{tpc}(G) = 3$ . Otherwise, if  $|N_X(u_1) \cap N_X(u_2)| = d_X(u_1) = d_X(u_2) = 1$ , then we know that  $\operatorname{tpc}(G) = 3$  from Claim 1. If  $|N_X(u_1) \cap N_X(u_2)| \ge 2$ , or  $|N_X(u_1) \cap N_X(u_2)| = 1$ , and  $\max\{d_X(u_1), d_X(u_2)\} \ge 2$ , then we can find that G has a Hamiltonian path. Thus,  $\operatorname{tpc}(G) = 3$ .

# 5 Proof of Theorem 1.4

Suppose  $F \cong K_3$  or  $P_3$ . Note that G has a Hamiltonian path, and so tpc(G) = 3. The following three claims will be used later.

**Claim 2.** Let *G* be a graph obtained by adding a pendant vertex  $u_3$  adjacent to vertex  $u_1$  or  $u_2$  of graph in Claim 1. Then tpc(G) = 3.

**Proof.** Without loss of generality, assume that  $u_3$  is adjacent to  $u_1$ . Let  $c(u_1) = \{1, 2, 3\} \setminus \{c(u_1v_1), c(v_1)\}$ ,  $c(u_1u_3) = c(v_1)$ , and the remaining vertices and edges are assigned the same color as Claim 1. We can verify that G is total proper connected with the above coloring. Then tpc(G) = 3. This completes the proof of Claim 2.

**Claim 3.** Let *G* be a graph obtained by adding three vertices  $\{u_1, u_2, u_3\}$  to a complete graph  $K_t$  such that  $d(u_1) = d(u_3) = 1$ ,  $d(u_2) = 2$ ,  $N(u_1) \cap N(u_3) = \emptyset$ , and  $|N(u_2) \cap N(u_i)| = 1$ , where  $1 \le i \le 2$ . Then tpc(G) = 3.

**Proof.** Without loss of generality, assume that  $N(u_2) \cap N(u_1) = v_1$ ,  $N(u_2) \cap N(u_3) = v_i$ . Let  $c(u_3v_i) = c(v_iv_{i+1})$ ,  $c(v_1u_2) = c(v_iv_{i-1})$ , and the remaining vertices and edges are assigned the same color as Claim 1. Suppose  $t \equiv 0 \pmod{3}$ . Note that  $P_1 = u_1v_1v_2 \cdots v_{t-1}v_t$ ,  $P_2 = u_2v_1v_tv_{t-1} \cdots v_3v_2$ ,  $P_3 = u_3v_iu_2$ ,  $P_4 = u_1v_1v_2 \cdots v_iu_3$ , and  $P_5 = u_3v_iv_{i-1} \cdots v_1v_iv_{t-1} \cdots v_{i+1}$  are five total proper paths. Suppose  $t \equiv 1 \pmod{3}$ . Note that  $P_1 = u_1v_1v_2 \cdots v_{t-1}v_t$ ,  $P_2 = u_2v_1v_{t-1} \cdots v_3v_2v_t$ ,  $P_3 = u_3v_iu_2$ ,  $P_4 = u_1v_1v_2 \cdots v_iu_3$  and  $P_5 = u_3v_iv_{i-1} \cdots v_2v_tv_{t-1} \cdots v_{i+1}$  are five total proper paths. Suppose  $t \equiv 2 \pmod{3}$ . We can find that  $P_1 = u_1v_1v_2 \cdots v_{t-1}v_t$ ,  $P_2 = u_2v_1v_{t-2} \cdots v_3v_2v_{t-1}$ ,  $P_3 = u_3v_iu_2$ ,  $P_4 = u_1v_1v_2 \cdots v_iu_3$ ,  $P_5 = u_3v_iv_{i-1} \cdots v_2v_tv_{t-1}v_{t-2} \cdots v_{i+1}$ , and  $P_6 = u_3v_iv_{i-1}v_t$  are six total proper paths. Therefore, G is total proper connected with the above coloring, and so tpc(G) = 3. This completes the proof of Claim 3.

**Claim 4.** Let *G* be a graph obtained by adding a vertex *u* to the graph in Claim 1 such that d(u) = 2 and  $v_1$  is adjacent to *u*. Then tpc(G) = 3.

Suppose  $F \cong K_2 + K_1$ . Let  $V(K_2) = \{u_1, u_2\}$  and  $V(K_1) = \{u_3\}$ . Since diam(G) = 2, we have  $N_X(u_1) \cap N_X(u_2) \cap V(u_2) \cap V(u_3)$  $N_X(u_3) = \{v\}$ , and so tpc(G) = 3 by Claim 2. Suppose  $F \cong 3K_1$ . Assume  $N_X(u_1) \cap N_X(u_2) \cap N_X(u_3) = \emptyset$ . Then tpc(G) = 3 by Claim 3. Assume  $N_X(u_1) \cap N_X(u_2) \cap N_X(u_3) \neq \emptyset$ . If  $d_X(u_1) = d_X(u_2) = d_X(u_3) = 1$ , then  $tpc(G) \ge 4$  by [26, Proposition 2]. Define a total coloring of G as follows:  $c(u_3v) = 4$  with  $v \in N_X(u_3)$ , and the remaining vertices and edges are assigned the same color as Claim 1. We check that any two vertices have a total proper path, and so tpc(G) = 4. Otherwise, we have  $d_X(u_1) + d_X(u_2) + d_X(u_3) \ge 4$ . Without loss of generality, let  $d_X(u_1) \ge 2$ , and  $u \in X \setminus \{v\}$  where  $v \in N_X(u_1) \cap N_X(u_2) \cap N_X(u_3)$ . Thus,  $\operatorname{tpc}(G) = 3$  by Claim 4.

# 6 Proof of Theorem 1.5

**Case 1.** diam(G) = 3. We prove Case 1 by analyzing the structure of F.

**Subcase 1.1.**  $F \cong K_3$  or  $P_3$ . Note that G has a Hamiltonian path, and so tpc(G) = 3.

**Subcase 1.2.**  $F \cong K_2 + K_1$ . Denote  $V(K_2) = \{u_1, u_2\}$  and  $V(K_1) = \{u_3\}$ . Suppose  $N_X(u_1) \cap N_X(u_3) \neq \emptyset$  or  $N_X(u_2) \cap N_X(u_3) \neq \emptyset$ . Without of loss generality, we may assume that  $N_X(u_1) \cap N_X(u_3) \neq \emptyset$ . Then tpc(G) =3 by Claim 2. Suppose  $N_X(u_1) \cap N_X(u_3) = \emptyset$  and  $N_X(u_2) \cap N_X(u_3) = \emptyset$ . Since diam(G) = 3, we have  $N_X(u_1) \neq \emptyset$ and  $N_X(u_3) \neq \emptyset$ . Then *G* has a Hamiltonian path, and so tpc(*G*) = 3.

**Subcase 1.3.**  $F \cong 3K_1$ . Let  $V(F) = \{u_1, u_2, u_3\}$ . Since diam(G) = 3, we have  $N_X(u_1) \cap N_X(u_2) \cap N_X(u_3) = \emptyset$ . Suppose there exists two vertices  $u_i$ ,  $u_i \in V(F)$  satisfy  $N_X(u_i) \cap N_X(u_i) \neq \emptyset$ . Without loss of generality, let  $u_1$ and  $u_2$  satisfy  $N_X(u_1) \cap N_X(u_2) \neq \emptyset$  and  $v_1 \in N_X(u_1) \cap N_X(u_2)$ . Assume  $d_X(u_1) = d_X(u_2) = 1$  and  $v_i \in N(u_3)$ . Since G is not complete, we have  $tpc(G) \ge 3$ . To the contrary, suppose there exists a total coloring c of G using three colors such that G is total proper connected. Since any two vertices of G are connected by a total proper path, we have  $c(u_1v_1) \neq c(v_1) \neq c(u_2v_1)$ . Without loss of generality, let  $c(u_1v_1) = 1$ ,  $c(v_1) = 2$  and  $c(u_2v_1) = 3$ . Consider the total proper path P between  $u_1$  and  $u_2$ , then the color of vertices and edges in P follows the sequence 1, 2, 3, ..., 1, 2, 3, .... Thus, the value of  $(c(v_i), c(v_iu))$  is (1, 2), (2, 3), or (3, 1). Consider the total proper path Q between  $u_2$  and u, then the color of vertices and edges in Q follows the sequence 3, 2, 1, ..., 3, 2, 1, .... But the value of  $(c(v_i), c(v_iu))$  is (3, 2), (2, 1), or (1, 3), a contradiction. Assign a total coloring c to G as follows:  $c(u_1v_1) = 1$ ,  $c(v_1) = c(v_1u) = 2$ ,  $c(u_2v_1) = 3$ , assign 4 to the remaining edges, and assign 1 to the remaining vertices. We can check that *G* is total proper connected with the above coloring, and so tpc(G) = 4. Assume  $d_X(u_1) + d_X(u_2) \ge 3$ , without loss of generality, let  $d_X(u_1) \ge 2$ . If  $d_X(u_3) = 1$  and  $N_X(u_1) \cap N_X(u_3) \neq \emptyset$ , then we have tpc(G) = 3 by Claim 3; if  $N_X(u_1) \cap N_X(u_3) = \emptyset$ , or  $N_X(u_1) \cap N_X(u_3) \neq \emptyset$  and  $d_X(u_3) \ge 2$ , then G has a Hamiltonian path, and so tpc(G) = 3.

Now, we may suppose  $N_X(u_1) \cap N_X(u_2) = \emptyset$ ,  $N_X(u_1) \cap N_X(u_3) = \emptyset$ , and  $N_X(u_2) \cap N_X(u_3) = \emptyset$ . Since G is not complete, we have  $tpc(G) \ge 3$ . To the contrary, assume that  $v_i \in N(u_i)$ , and there exists a total coloring c of G using three colors such that G is total proper connected. Since any two vertices of G are connected by a total proper path, we have  $c(u_1v_1) \neq c(v_1)$ . Without loss of generality, let  $c(u_1v_1) = 1$  and  $c(v_1) = 2$ . Consider the total proper path P between  $u_1$  and  $u_2$ , then the color of vertices and edges in P follows the sequence 1, 2, 3, ..., 1, 2, 3, .... Thus, the value of  $(c(v_2), c(v_2u_2))$  is (1, 2), (2, 3), or (3, 1). Consider the total proper path Q between  $u_1$  and  $u_3$ , then the color of vertices and edges in Q follows the sequence 1, 2, 3, ..., 1, 2, 3, .... Hence, the value of  $(c(v_3), c(v_3u_3))$  is (1, 2), (2, 3), or (3, 1). Consider the total proper path W between  $u_2$ and  $u_3$ . If  $c(v_2) = 1$  and  $c(v_2u_2) = 2$ , then the color of vertices and edges in W follows the sequence 2, 1, 3, ..., 2, 1, 3, .... Note that the value of  $(c(v_3), c(v_3u_3))$  is (2, 1), (1, 3), or (3, 2), a contradiction. If  $c(v_2) = 2$ and  $c(v_2u_2) = 3$ , then the color of vertices and edges in W follows the sequence 3, 2, 1, ..., 3, 2, 1, .... Note that the value of  $(c(v_3), c(v_3u_3))$  is (2, 1), (1, 3), or (3, 2), a contradiction. If  $c(v_2) = 3$  and  $c(v_2u_2) = 1$ , then the color of vertices and edges in W follows the sequence  $1, 3, 2, \dots, 1, 3, 2, \dots$ . Note that the value of  $(c(v_3), c(v_3u_3))$  is (2, 1), (1, 3), or (3, 2), a contradiction. Assign a total coloring c to G as follows:  $c(u_1v_1) = c(v_2) = 1$ ,  $c(v_1) = c(u_2v_2) = c(u_3v_3) = 2$ , assign 4 to the remaining vertices, and assign 3 to the remaining edges. We can verify that G is total proper connected with the above coloring, and so tpc(G) = 4. **Case 2.** diam $(G) \ge 4$ . Thus,  $F \cong P_3$  or  $F \cong K_2 + K_1$ . Assume  $F \cong P_3$ . Obviously, G has a Hamiltonian path, and so  $\operatorname{tpc}(G) = 3$ . Assume  $F \cong K_2 + K_1$ . Denote  $V(K_2) = \{u_1, u_2\}$  and  $V(K_1) = \{u_3\}$ , without loss of generality, we have  $d_X(u_2) = 0$ ,  $d_X(u_1) \ge 1$ ,  $d_X(u_3) \ge 1$  satisfying  $N_X(u_1) \cap N_X(u_2) = \emptyset$ . Hence, we can find that G has a Hamiltonian path, and so tpc(G) = 3.

# 7 Proof of Theorem 1.6

The proof of Theorem 1.6 follows from the next two lemmas. First, we shall determine the total proper *k*-connection numbers of the circular ladders.

**Lemma 7.1.** Let n be an integer with  $n \ge 3$ . Then  $tpc(CL_{2n}) = tpc_2(CL_{2n}) = 3$ ,  $tpc_3(CL_{2n}) = 4$ .

**Proof.** Let *n* be an integer with  $n \ge 3$ . Since  $CL_{2n}$  contains a Hamiltonian path that is not complete, we have  $tpc(CL_{2n}) = 3$ . Since  $tpc_2(CL_{2n}) \ge tpc(CL_{2n}) = 3$ , we only need to prove  $tpc_2(CL_{2n}) \le 3$ .

**Case 1.**  $n \equiv 0 \pmod{3}$ . Let n = 3t. Assign a total coloring c to  $CL_{2n}$  as follows: Let i be an integer with  $0 \le n$  $i \le t - 1, c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) = c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = 1, c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(v_{3i+3}) = c(v_{3i+1}v_{3i+2}) = 2,$ and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = 3$ ;  $c(u_iv_i), c(u_i), c(v_i) \in \{1, 2, 3\}$  with  $c(u_iv_i) \neq c(u_i) \neq c(v_i)$ for  $1 \le i \le n$ . Let x and y be any two distinct vertices of  $CL_{2n}$ . By symmetry, we may assume that  $x = u_1$ . If  $y = u_i$  for  $2 \le i \le n$ , then  $xu_2u_3 \cdots u_{i-1}y$  and  $xu_nu_{n-1} \cdots u_{i+1}y$  are two total proper paths connecting x and y. If  $y = v_1$ , then xy and  $xu_nu_{n-1}\cdots u_2y$  are two total proper paths connecting x and y. If  $y = v_i$ , then  $xv_1v_nv_{n-1}\cdots v_{i+1}y$  and  $xu_nu_{n-1}\cdots u_iy$  are two total proper paths connecting x and y. Thus,  $CL_{2n}$  is total proper 2-connected with the above coloring.

**Case 2.**  $n \equiv 1 \pmod{3}$ . Let n = 3t + 1. Define a total coloring c of  $CL_{2n}$  as follows: Let i be an integer with  $0 \le i \le t - 1, \ c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) = c(v_{3i+3}) = c(v_{3i+1}v_{3i+2}) = 1, \ c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = c(v_{3i+$ = 3, and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = 2$ ;  $c(u_n) = c(v_1v_n) = 1$ ,  $c(u_nv_n) = 2$ ,  $c(v_n) = c(u_1u_n) = 2$ 3, and  $c(u_iv_i) = 3$  for  $1 \le i \le n - 1$ . Let x and y be any two distinct vertices of  $CL_{2n}$ . We may assume that  $x = u_i$ for  $1 \le i \le n-1$ . If  $y = u_i$  for  $i \le j \le n-1$ , then  $xu_{i+1}u_{i+2}\cdots u_{j-1}y$  and  $xu_{i-1}u_{i-2}\cdots u_1v_1v_2\cdots v_{n-1}u_{n-1}u_{n-2}\cdots u_{i+1}y$ are two total proper paths connecting x and y. If  $y = v_i$  for  $1 \le j \le n - 1$ , then  $xu_{i-1}u_{i-2} \cdots u_1v_1v_2 \cdots v_{i-1}y$  and  $xu_{i+1}\cdots u_{n-1}v_{n-1}v_{n-2}\cdots v_{i+1}y$  are two total proper paths connecting x and y. If  $y=u_n$ , then  $xu_{i+1}\cdots u_{n-1}y$ and  $xu_{i-1}u_{i-2}\cdots u_1y$  are two total proper paths connecting x and y. If  $y=v_n$ , then  $xu_{i+1}\cdots u_{n-1}u_ny$  and  $xu_{i-1}u_{i-2}\cdots u_1v_1v_2\cdots v_{n-1}y$  are two total proper paths connecting x and y. Assume that  $x=u_n$ . If  $y=v_i$  for  $1 \le j \le n-1$ , then  $xu_1u_2 \cdots u_{n-1}v_{n-1}v_{n-2} \cdots v_{j+1}y$  and  $xu_nv_nv_1v_2 \cdots v_{j-1}y$  are two total proper paths connecting x and y. If  $y = v_n$ , then xy and  $xu_{n-1}u_{n-2}\cdots u_1v_1y$  are two total proper paths connecting x and y. Thus,  $CL_{2n}$  is total proper 2-connected with the above coloring.

**Case 3.**  $n \equiv 2 \pmod{3}$ . Let n = 3t + 2. Define a total coloring c of  $CL_{2n}$  as follows: Let i be an integer with  $0 \leq i \leq t-2, c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) = c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = 1, c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = c(v_{3i+3}v_{3$ = 3, and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = 2$ ;  $c(u_n) = c(u_{n-4}) = c(v_{n-2}) = c(v_n) = 1$ ,  $c(u_{n-1}) = c(u_{n-1}) = c(u_{n-1})$  $c(v_{n-1}) = c(u_{n-3}) = c(v_{n-3}) = 3$ ,  $c(u_{n-2}) = c(v_{n-4}) = 2$ ,  $c(u_{n-4}u_{n-3}) = c(u_{n-1}u_n) = c(v_{n-1}v_n) = c(v_{n-3}v_{n-2}) = c(v_{n-3}v_{n-2})$  $c(v_{n-2}v_{n-1}) = 2$ ,  $c(u_{n-3}u_{n-2}) = c(u_{n-2}u_{n-1}) = c(v_{n-4}v_{n-3}) = c(u_2v_2) = 1$ ,  $c(u_1u_1) = c(u_1u_n) = c(v_1v_n) = 3$ ;  $c(u_{4j}v_{4j}) = c(u_{4j}v_{4j}) = c(u_{4j}v_{$ 1,  $c(u_{4j-1}v_{4j-1}) = 2$ ,  $c(u_{4j-1}v_{4j+1}) = 3$ , where  $1 \le j \le t-1$ . Note that  $u_1u_2 \cdots u_{n-2}v_{n-2}v_{n-2}v_{n-3} \cdots v_1u_1$  is a total proper cycle, any two distinct vertices of the cycle have two disjoint total proper paths. Now, we may assume that  $x = u_i$  for  $1 \le i \le n - 2$ . If  $y = u_{n-1}$ , then  $xu_{i+1} \cdots u_{n-2}v_{n-2}v_{n-1}y$  and  $xu_{i-1} \cdots u_1u_ny$  are two total proper paths connecting x and y. If  $y = u_n$ , then  $xu_{i+1} \cdots u_{n-2}v_{n-2}v_{n-1} \cdots v_1v_ny$  and  $xu_{i-1} \cdots u_1v_1v_ny$  are two total proper paths connecting x and y. If  $y = v_{n-1}$ , then  $xu_{i-1} \cdots u_3v_3v_2v_1v_ny$  and  $xu_{i+1} \cdots u_{n-2}v_{n-2}y$  are two total proper paths connecting x and y. If  $y = v_n$ , then  $xu_{i-1} \cdots u_1u_ny$  and  $xu_{i+1} \cdots u_{n-2}v_{n-2}v_{n-1}y$  are two total proper paths connecting x and y. Assume that  $x = v_i$  for  $1 \le i \le n - 2$ . If  $y = v_{n-1}$ , then  $xv_{i+1} \cdots v_{n-2}u_{n-2}u_{n-1} \cdots u_1u_nu_{n-1}y$  and  $xv_{i-1}\cdots v_1v_ny$  are two total proper paths connecting x and y. If  $y=v_n$ , then  $xv_{i+1}\cdots v_{n-2}u_{n-2}u_{n-1}u_ny$  and  $xv_{i-1} \cdots v_1y$  are two total proper paths connecting x and y. Thus,  $CL_{2n}$  is total proper 2-connected with the above coloring.

To the contrary, suppose there exists a total proper 3-connected coloring c of  $CL_{2n}$  using three colors. Considering  $u_1$  and  $v_2$ ,  $u_1u_2v_2$ ,  $u_1v_1v_2$ , and  $u_1u_nv_nv_{n-1}\cdots v_3v_2$  must be three total proper paths connecting  $u_1$  and  $v_2$ . Then  $c(u_1u_2) \neq c(u_2v_2) \neq c(u_2)$ . Considering  $u_1$  and  $u_3$ ,  $u_1u_2u_3$ ,  $u_1u_2u_3$ ,  $u_1u_2u_3$ , and  $u_1v_1v_2v_3u_3$  must be three total proper paths connecting  $u_1$  and  $u_3$ . Hence,  $c(u_1u_2) \neq c(u_2u_3) \neq c(u_2)$ , and so  $c(u_2v_2) = c(u_2u_3)$ . But then, there is no set of three disjoint total proper paths connecting  $u_3$  and  $v_2$ , a contradiction. Hence,  $tpc_3(CL_{2n}) \ge 4$ . Now we only need to prove  $tpc_3(CL_{2n}) \le 4$ .

**Case 1.**  $n \equiv 0 \pmod{3}$ . Let n = 3t. Assign a total coloring c to  $CL_{2n}$  as follows: Let i be an integer with  $0 \le c$  $i \le t-1$ ,  $c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) = c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = 1$ ,  $c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(v_{3i+3}) = c(v_{3i+3}v_{3i+2}) = 2$ ,

and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = 3$ ;  $c(u_iv_i) = 4$  for  $1 \le i \le n$ . Let x and y be any two distinct vertices of  $CL_{2n}$ . By symmetry, we may assume that  $x=u_1$ . If  $y=u_i$  for  $2 \le i \le n$ , then  $xu_2u_3\cdots u_{i-1}y_1xu_nu_{n-1}\cdots u_{i+1}y_1$ , and  $xv_1v_2\cdots v_iy_1$  are three total proper paths connecting x and y. If  $y = v_1$ , then xy,  $xu_2v_2y$  and  $xu_nv_ny$  are three total proper paths connecting x and y. If  $y = v_i$  for  $2 \le i \le n$ , then  $xu_1u_2 \cdots u_iy_1xv_1v_2 \cdots v_{i-1}y_i$ , and  $xu_nv_nv_{n-1}\cdots u_iy_i$  are three total proper paths connecting x and y. Hence,  $CL_{2n}$  is total proper 3-connected with the above coloring.

**Case 2.**  $n \equiv 1 \pmod{3}$ . Let n = 3t + 1. Define a total coloring c of  $CL_{2n}$  as follows:  $c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) =$  $c(v_{3i+3}) = c(v_{3i+1}v_{3i+2}) = 1$ ,  $c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = 3$ , and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(u_{3i+1}u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(u_{3i+3}u_{3i+$  $c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = 2$ , where  $0 \le i \le t-1$ ;  $c(u_n) = c(v_1v_n) = 1$ ,  $c(u_nv_n) = 2$ ,  $c(v_n) = c(u_1u_n) = 4$ , and  $c(u_iv_i) = 1$ = 4 for  $2 \le j \le n - 1$ . Let x and y be any two distinct vertices of  $CL_{2n}$ . By symmetry, we may assume that  $x = u_1$ . If  $y = u_i$  for  $2 \le i \le n$ , then  $xu_2u_3 \cdots u_{i-1}y$ ,  $xu_nu_{n-1} \cdots u_{i+1}y$ , and  $xv_1v_2 \cdots v_iy$  are three total proper paths connecting x and y. If  $y = v_1$ , then xy,  $xu_2v_2y$  and  $xu_nv_ny$  are three total proper paths connecting x and y. If  $y = v_i$  for  $2 \le i \le n$ , then  $xu_1u_2 \cdots u_iy$ ,  $xv_1v_2 \cdots v_{i-1}y$ , and  $xu_nv_nv_{n-1} \cdots u_iy$  are three total proper paths connecting x and y. Hence,  $CL_{2n}$  is total proper 3-connected with the above coloring.

**Case 3.**  $n \equiv 2 \pmod{3}$ . Let n = 3t + 2. Define a total coloring c of  $CL_{2n}$  as follows:  $c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) =$  $c(v_{3i+4}) = c(v_{3i+2}v_{3i+3}) = 1$ ,  $c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(v_{3i+3}) = c(v_{3i+1}v_{3i+2}) = 3$ , and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(u_{3i+2}u_{3i+3}) = c(u_{3i+3}u_{3i+4}) = c(u_{3i+3}u_{3i+$  $c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = 2$ , where  $0 \le i \le t-1$ ;  $c(u_n) = c(v_{n-1}v_n) = 3$ ,  $c(v_n) = c(u_{n-1}u_n) = c(u_1v_1) = 2$ ,  $c(u_{n-1}v_n) = 2$  $=c(v_1v_n)=c(u_nu_1)=1$ ,  $c(v_1)=4$ , and  $c(u_iv_j)=4$  for  $2 \le j \le n$ . Let x and y be any two distinct vertices of  $CL_{2n}$ . By symmetry, we may assume that  $x = u_1$ . If  $y = u_i$  for  $2 \le i \le n$ , then  $xu_2u_3 \cdots u_{i-1}y$ ,  $xu_nu_{n-1} \cdots u_{i+1}y$ , and  $xv_1v_2 \cdots v_jy$  are three total proper paths connecting x and y. If  $y = v_1$ , then xy,  $xu_2v_2y$ , and  $xu_nv_ny$  are three total proper paths connecting x and y. If  $y = v_i$  for  $2 \le i \le n$ , then  $xu_1u_2 \cdots u_iy$ ,  $xv_1v_2 \cdots v_{i-1}y$  and  $xu_nv_nv_{n-1}\cdots u_iy$  are three total proper paths connecting x and y. Hence,  $CL_{2n}$  is total proper 3-connected with the above coloring.

Next, we shall determine the total proper *k*-connection numbers of the Möbius ladders.

**Lemma 7.2.** Let n be an integer with  $n \ge 3$ . Then  $\operatorname{tpc}(M_{2n}) = \operatorname{tpc}_2(M_{2n}) = 3$ ,  $\operatorname{tpc}_3(M_{2n}) = 4$ .

**Proof.** Since  $M_{2n}$  contains a Hamiltonian path and is not complete, we have  $tpc(M_{2n}) = 3$ . Since  $tpc_2(M_{2n}) \ge 1$  $tpc(M_{2n}) = 3$ , we only need to prove  $tpc_2(M_{2n}) \le 3$ .

**Case 1.**  $n \equiv 0 \pmod{3}$ . Define a total coloring  $c = 0 \pmod{3}$  as follows: Let  $i = 0 \pmod{3}$ .  $c(u_i) = c(v_{n-i+1}) = 1$ ,  $c(u_{i+1}) = c(v_{n-i}) = 3$ , and  $c(u_{i+2}) = c(v_{n-i-1}) = 2$ ;  $c(u_iu_{i+1})$ ,  $c(u_i)$ ,  $c(u_{i+1}) \in \{1, 2, 3\}$  with  $c(u_iu_{i+1}) \neq c(u_i) \neq c(u_{i+1})$  for  $1 \leq i \leq n-1$ ;  $c(v_iv_{i+1}), c(v_i), c(v_{i+1}) \in \{1, 2, 3\}$  with  $c(v_iv_{i+1}) \neq c(v_i) \neq c(v_{i+1})$  for  $1 \le i \le n - 1$ ;  $c(u_1v_1) = 3$ ,  $c(u_nv_n) = 2$ , and  $c(u_jv_{n-j+1}) = 3$  for  $1 \le j \le n$ . Let x and y be any two distinct vertices of  $M_{2n}$ . By symmetry, we may assume that  $x = u_1$ . If  $y = u_i$  for  $2 \le i \le n$ , then  $xu_2u_3 \cdots u_{i-1}y$  and  $xv_1v_2$  $\cdots v_n u_n u_{n-1} \cdots u_{i+1} y$  are two total proper paths connecting x and y. If  $y = v_1$ , then xy and  $x u_2 u_3 y$  are two total proper paths connecting x and y. If  $y = v_i$  for  $2 \le i \le n$ , then  $xu_1u_2 \cdots u_nv_nv_{n-1} \cdots v_{i+1}y$  and  $xv_1v_2 \cdots v_{i-1}y$ are two total proper paths connecting x and y. Thus,  $M_{2n}$  is total proper 2-connected with the above coloring.

**Case 2.**  $n \equiv 1 \pmod{3}$ . Let n = 3t + 1. Define a total coloring c of  $M_{2n}$  as follows: Let i be an integer with  $0 \le i \le t - 1, c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) = c(v_{3i+3}) = c(v_{3i+3}u_{3i+2}) = 1, c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = c(v_{3i+3}u_{3i+3}) = c(v_{3i+3}u_$ = 3, and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = 2$ ; if j = 3i for  $1 \le i \le t$ , then  $c(u_jv_{n-j+1}) = 3$ ; if  $j \ne 3i$ for  $1 \le i \le t$ , then  $c(u_i v_{n-j+1})$ ,  $c(u_i)$ ,  $c(v_{n-j+1}) \in \{1, 2, 3\}$  with  $c(u_i v_{n-j+1}) \ne c(u_i) \ne c(v_{n-j+1})$ ;  $c(u_n) = 1$ ,  $c(u_n v_n) = 1$  $c(u_1v_1) = 2$ ,  $c(v_n) = 3$ . Let x and y be any two distinct vertices of  $M_{2n}$ . We may assume that  $x = u_i$  for  $1 \le i \le n$ . If  $y = u_i$  for  $i \le j \le n$ , then  $xu_{i+1}u_{i+2}\cdots u_{j-1}y$  and  $xu_{i-1}u_{i-2}\cdots u_1v_nu_nu_{n-1}\cdots u_{j+1}y$  are two total proper paths connecting x and y. If  $y = v_j$  for  $1 \le j \le n$ , then  $xu_{i-1}u_{i-2}\cdots u_1v_nv_{n-1}\cdots v_{j+1}y$  and  $xu_{i+1}\cdots u_{n-j+1}y$  are two total proper paths connecting x and y. Thus,  $M_{2n}$  is total proper 2-connected with the above coloring.

**Case 3**.  $n \equiv 2 \pmod{3}$ . Let n = 3t + 2. Define a total coloring c of  $M_{2n}$  as follows: Let i be an integer with  $0 \le i \le t - 1, \ c(u_{3i+1}) = c(u_{3i+2}u_{3i+3}) = c(v_{3i+3}) = c(v_{3i+1}v_{3i+2}) = c(u_{3i+2}v_{n-3i-1}) = 1, \ c(u_{3i+2}) = c(u_{3i+3}u_{3i+4}) = 1$  $c(v_{3i+1}) = c(v_{3i+2}v_{3i+3}) = c(u_{3i+3}v_{n-3i-2}) = 3$ , and  $c(u_{3i+3}) = c(u_{3i+1}u_{3i+2}) = c(v_{3i+2}) = c(v_{3i+3}v_{3i+4}) = c(u_{3i+4}v_{n-3i-3})$ = 2;  $c(u_{n-1}) = c(u_n v_1) = c(v_{n-1} v_n) = 1$ ,  $c(u_n) = c(v_{n-1}) = c(u_1 v_n) = 3$ ,  $c(v_n) = c(u_{n-1} u_n) = 2$ . Let x and y be any two distinct vertices of  $M_{2n}$ . We may assume that  $x = u_i$  for  $1 \le i \le n$ . If  $y = u_i$  for  $i \le j \le n$ , then

 $xu_{i+1}u_{i+2}\cdots u_{j-1}y$  and  $xu_{i-1}u_{i-2}\cdots u_1v_nu_nu_{n-1}\cdots u_{j+1}y$  are two total proper paths connecting x and y. If  $y=v_i$ for  $1 \le j \le n$ , then  $xu_{i-1}u_{i-2} \cdots u_1v_nv_{n-1} \cdots v_{j+1}y$  and  $xu_{i+1} \cdots u_{n-j+1}y$  are two total proper paths connecting x and y. Thus,  $M_{2n}$  is total proper 2-connected with the above coloring.

To the contrary, suppose there exists a total proper 3-connected coloring c of  $M_{2n}$  using three colors. By considering the pair  $\{u_2, v_n\}$ ,  $u_2u_3 \cdots u_nv_n$ ,  $u_2u_1v_n$ , and  $u_2v_{n-1}v_n$  must be three total proper paths connecting  $u_2$  and  $v_n$ . Then  $c(u_2v_{n-1}) \neq c(v_{n-1}v_n) \neq c(v_{n-1})$ . By considering the pair  $\{u_2, v_{n-2}\}, u_2u_1v_nv_{n-1}v_{n-2}, u_2u_3v_{n-2}, u_2u_3v_{n-2}\}$ and  $u_2v_{n-1}v_n$  must be three total proper paths connecting  $u_2$  and  $v_{n-2}$ . Thus,  $c(u_2v_{n-1}) \neq c(v_{n-1}v_{n-2}) \neq c(v_{n-1})$ , and hence  $c(v_{n-1}v_{n-2}) = c(v_{n-1}v_n)$ . But then, there is no set of three disjoint total proper paths connecting  $v_{n-2}$  and  $v_n$ , a contradiction. Thus,  $\operatorname{tpc}_3(M_{2n}) \geq 4$ . Now we only need to prove  $\operatorname{tpc}_3(M_{2n}) \leq 4$ .

**Case 1.**  $n \equiv 0 \pmod{2}$ . Let n = 2t. Assign a total coloring c of  $M_{2n}$  as follows:  $c(u_{2i+1}) = c(v_{n-2i}) = c(v_{n-2$  $c(u_{2i+2}v_{n-2i-1}) = 1, c(u_{2i+2}) = c(u_{n-2i-1}) = c(u_{2i+1}v_{n-2i}) = 3, c(v_{n-2i}v_{n-2i-1}) = c(u_{2i+1}u_{2i+2}) = 2, \text{ where } 0 \le i \le t-1;$  $c(u_{2i+2}u_{2i+3}) = c(v_{n-2i-1}v_{n-2i-2}) = 4$  for  $0 \le i \le t-2$ ;  $c(u_nv_n) = c(u_1v_1) = 4$ . Let x and y be any two distinct vertices of  $M_{2n}$ . By symmetry, we may assume that  $x = u_1$ . If  $y = u_i$  for  $2 \le i \le n$ , then  $xu_2 \cdots u_{i-1}y$ ,  $xv_nu_nu_{n-1}\cdots u_{i+1}y$ , and  $xv_1v_2\cdots v_{n-i+1}u_i$  are three total proper paths connecting x and y. If  $y=v_1$ , then xy,  $xu_2u_3\cdots u_ny$ , and  $xv_nv_{n-1}\cdots v_2y$  are three total proper paths connecting x and y. If  $y=v_i$  for  $2\leq i\leq n$ , then  $xv_1v_2 \cdots v_{i-1}y$ ,  $xv_nv_{n-1} \cdots v_{i+1}y$  and  $xu_2u_3 \cdots u_{n-i+1}y$  are three total proper paths connecting x and y. Hence,  $M_{2n}$  is total proper 3-connected with the above coloring.

**Case 2.**  $n \equiv 1 \pmod{2}$ . For n = 3, we assign a total coloring c to  $M_6$  as follows:  $c(u_1) = c(v_3) = 1$  $c(u_2u_3) = c(v_1v_2) = 1$ ,  $c(u_2) = c(v_2) = c(u_3v_3) = c(u_1v_1) = 2$ ,  $c(u_3) = c(v_1) = c(u_1u_2) = c(v_2v_3) = 3$ ,  $c(u_1v_{n-i+1}) = 4$ for  $1 \le i \le 3$ . We can verify that  $M_6$  is total proper 3-connected with the above total coloring, and so  $tpc_3(M_6) = 4$ .

For n = 5, define a total coloring c of  $M_{10}$  as follows:  $c(u_1) = c(v_5) = 1$ ,  $c(u_2) = c(v_4) = 2$ ,  $c(u_3) = c(v_3) = 2$  $c(u_5) = c(v_1) = 3, c(u_4) = c(v_2) = 4, c(u_2u_3) = c(u_4u_5) = c(v_1v_2) = c(v_3v_4) = 1, c(u_3u_4) = c(u_5v_5) = c(u_1v_1) = c(v_2v_3) = c(v_3v_4) = 1, c(u_3v_4) = 1$ 2,  $c(u_4v_2) = c(u_1u_2) = c(v_4v_5) = 3$ ,  $c(u_iv_{n-i+1}) = 4$  for i = 1, 2, 3, 5. We can check that  $M_{10}$  is total proper 3-connected with the above total coloring, and so  $tpc_3(M_{10}) = 4$ .

**Subcase 2.1.** Let  $n \equiv 1 \pmod{4}$  for  $n \ge 7$ . Let n = 2t + 1. Assign a total coloring c to  $M_{2n}$  as follows:  $c(u_{4i+1}) = c(v_{n-4i}) = 1$ ,  $c(u_{4i+2}) = c(v_{n-4i-1}) = 2$ ,  $c(u_{4i+3}) = c(v_{n-4i-2}) = 3$ , and  $c(u_{4i+4}) = c(v_{n-4i-3}) = 4$ , where  $0 \le i \le \frac{t-2}{2}; c(u_{4i}u_{4i+1}) = c(v_{n-4i+1}v_{n-4i}) = c(u_{4i+2}v_{n-4i-1}) = 3, c(u_{4i+1}u_{4i+2}) = c(v_{n-4i}v_{n-4i-1}) = c(u_{4i+3}v_{n-4i-2}) = 4,$  $c(u_{4i+2}u_{4i+3}) = c(v_{n-4i-1}v_{n-4i-2}) = c(u_{4i}v_{n-4i+1}) = 1$ , and  $c(v_{n-4i-2}v_{n-4i-3}) = c(u_{4i+3}u_{4i+4}) = c(u_{4i+1}v_{n-4i}) = 2$ , where  $1 \le i \le \frac{t-2}{2}$ ;  $c(u_n) = c(v_1) = c(u_1u_2) = c(v_nv_{n-1}) = 3$ ,  $c(u_2u_3) = c(v_{n-1}v_{n-2}) = c(v_1v_2) = c(u_{n-1}u_n) = c(u_{n-1}v_2)$ = 1,  $c(u_3u_4) = c(v_{n-2}v_{n-3}) = c(u_1v_1) = c(u_nv_n) = 2$ , and  $c(u_1v_n) = c(u_2v_{n-1}) = c(u_3v_{n-2}) = 4$ . Let x and y be any two distinct vertices of  $M_{2n}$ . By symmetry, we may assume that  $x = u_1$ . If  $y = u_i$  for  $2 \le i \le n$ , then  $xu_2u_3\cdots u_{i-1}y$ ,  $xv_nu_nu_{n-1}\cdots u_{i+1}y$ , and  $xv_1v_2\cdots v_{n-i+1}u_i$  are three total proper paths connecting x and y. If  $y = v_1$ , then xy,  $xu_2u_3 \cdots u_ny$ , and  $xv_nv_{n-1} \cdots v_2y$  are three total proper paths connecting x and y. If  $y = v_1$ for  $2 \le i \le n$ , then  $xv_1v_2 \cdots v_{i-1}y$ ,  $xv_nv_{n-1} \cdots v_{i+1}y$  and  $xu_2u_3 \cdots u_{n-i+1}y$  are three total proper paths connecting xand y. Hence,  $M_{2n}$  is total proper 3-connected with the above coloring.

**Subcase 2.2.** Let  $n \equiv 3 \pmod{4}$  for  $n \ge 7$ . Let n = 2t + 1. Assign a total coloring c to  $M_{2n}$  as follows: Let ibe an integer with  $0 \le i \le \frac{t-3}{2}$ ,  $c(u_{4i+1}) = c(v_{n-4i}) = 1$ ,  $c(u_{4i+2}) = c(v_{n-4i-1}) = 2$ ,  $c(u_{4i+3}) = c(v_{n-4i-2}) = 3$ , and  $c(u_{4i+4}) = c(v_{n-4i-3}) = 4$ . Let *i* be an integer with  $1 \le i \le \frac{t-1}{2}$ ,  $c(u_{4i}u_{4i+1}) = c(v_{n-4i-1}v_{n-4i}) = c(u_{4i+2}v_{n-4i-1}) = 3$ ,  $c(u_{4i+1}u_{4i+2}) = c(v_{n-4i}v_{n-4i-1}) = c(u_{4i+3}v_{n-4i-2}) = 4,$  $c(u_{4i+2}u_{4i+3}) = c(u_{4i}v_{n-4i+1}) = c(v_{n-4i-1}v_{n-4i-2}) = 1,$ and  $c(u_{4i+1}v_{n-4i}) = c(u_{4i-1}u_{4i}) = c(v_{n-4i+2}v_{n-4i+1}) = 2$ ;  $c(u_n) = c(v_1) = c(u_1u_2) = c(v_nv_{n-1}) = 3$ ,  $c(u_{n-1}) = c(v_2) = c(v_nv_{n-1}) = 3$  $c(u_1v_1) = c(u_nv_n) = 2$ , and  $c(u_{n-2}) = c(v_3) = c(u_2u_3) = c(v_{n-1}v_{n-2}) = 1$ . By the similar proof of the above subcase, we can verify  $M_{2n}$  is total proper 3-connected with the above coloring.

## 8 Proof of Theorem 1.7

Note that  $K_3 \square K_2 = CL_6$ ,  $K_{3,3} = M_6$ ,  $Q_3 = CL_8$  and  $M_8$ . By means of Theorem 1.6, we can obtain their total proper k-connection numbers. Now, we only need to consider  $K_4$ ,  $F_1$ ,  $F_2$ ,  $F_3$ .

**Lemma 8.1.**  $tpc(K_4) = 1$ ,  $tpc_2(K_4) = 3$ ,  $tpc_3(K_4) = 4$ .

**Proof.** By [26], we know that  $tpc(K_4) = 1$ . Suppose  $tpc_2(K_4) = 2$ , then there is no set of two disjoint total proper paths connecting  $u_1$  and  $u_2$ , a contradiction. Thus,  $tpc_2(K_4) \ge 3$ . Let  $V(K_4) = \{u_1, u_2, u_3, u_4\}$ , we assign a total coloring c to  $K_4$  as follows:  $c(u_1) = c(u_2u_3) = c(u_3u_4) = c(u_2u_4) = 1$ ,  $c(u_2) = c(u_4) = c(u_1u_3) = 2$ ,  $c(u_3) = c(u_1u_2) = c(u_1u_4) = 3$ . We can verify that the  $K_4$  is total proper 2-connected, so  $tpc_2(K_4) = 3$ .

Now, we suppose there exists a total proper 3-connected coloring c of  $K_4$  using three colors. Considering  $u_1$  and  $u_2$ ,  $u_1u_2$ ,  $u_1u_4u_2$ , and  $u_1u_3u_2$  must be the three total proper paths connecting  $u_1$  and  $u_2$ . Then  $c(u_1u_4) \neq c(u_2u_4) \neq c(u_4)$ . Considering  $u_1$  and  $u_3$ ,  $u_1u_3$ ,  $u_1u_2u_3$ , and  $u_1u_4u_3$  must be the three total proper paths connecting  $u_1$  and  $u_3$ . Thus,  $c(u_1u_4) \neq c(u_3u_4) \neq c(u_4)$ , and so  $c(u_2u_4) = c(u_4u_3)$ . But then, there is no set of three disjoint total proper paths connecting  $u_2$  and  $u_3$ , a contradiction. Hence,  $\operatorname{tpc}_3(K_4) \geq 4$ . Define a total coloring c of c0 of c1 as follows: c1 of c2 and c3 of c3 of c4 as follows: c4. We can easily check that c4 is total proper 3-connected, and so  $\operatorname{tpc}_3(K_4) = 4$ .

**Lemma 8.2.** 
$$tpc(F_1) = tpc_2(F_1) = 3$$
,  $tpc_3(F_1) = 4$ .

**Proof.** Since  $F_1$  has a Hamiltonian path that is not complete, we know that  $tpc(F_1) = 3$ . It is easy to verify that  $F_1$  is total proper 2-connected depicted in Figure 4(a), and so  $tpc_2(F_1) = 3$ . Now, we suppose there exists a total proper 3-connected coloring c of  $F_1$  using three colors. By considering the pair  $\{u_2, u_8\}$ ,  $u_2u_8$ ,  $u_2u_1u_8$ , and  $u_2u_3 \cdots u_8$  must be the three total proper paths connecting  $u_2$  and  $u_8$ . Then  $c(u_1u_2) \neq c(u_1u_8) \neq c(u_1)$ . By considering the pair  $\{u_5, u_8\}$ ,  $u_5u_6u_7u_8$ ,  $u_5u_1u_8$ , and  $u_5u_4u_3u_2u_8$  must be the three total proper paths connecting  $u_5$  and  $u_8$ . Then  $c(u_1u_5) \neq c(u_1u_8) \neq c(u_1)$ , and hence  $c(u_1u_2) = c(u_1u_5)$ . But then, there is no set of three disjoint total proper paths connecting  $u_2$  and  $u_5$ , a contradiction. Hence,  $tpc_3(F_1) \geq 4$ . By Figure 4(b), we know that  $F_1$  is total proper 3-connected, and so  $tpc_3(F_1) = 4$ .

**Lemma 8.3.**  $tpc(F_2) = tpc_2(F_2) = 3$ .

**Proof.** Since  $F_1$  has a Hamiltonian path that is not complete, we know that  $tpc(F_2) = 3$ . We can check that the coloring shown in Figure 4(c) is total proper 2-connected. Thus,  $tpc_2(F_2) = 3$ .

**Lemma 8.4.**  $tpc(F_3) = tpc_2(F_3) = 3$ ,  $tpc_3(F_3) = 4$ .

**Proof.** Since  $F_3$  has a Hamiltonian path that is not complete, we know that  $tpc(F_3) = 3$ . It is easy to check that the coloring shown in Figure 4(d) is total proper 2-connected using three colors. Thus,  $tpc_2(F_3) = 3$ . Now, we suppose there exists a total proper 3-connected coloring c of  $F_3$  using three colors. Considering  $u_2$  and  $u_8$ ,  $u_2u_8$ ,  $u_2u_1u_8$  and  $u_2u_3 \cdots u_8$  must be the three total proper paths connecting  $u_2$  and  $u_8$ . Then  $c(u_1u_2) \neq c(u_1u_8) \neq c(u_1)$ . Considering  $u_5$  and  $u_8$ ,  $u_5u_6u_7u_8$ ,  $u_5u_1u_8$ , and  $u_5u_4u_3u_2u_8$  must be the three

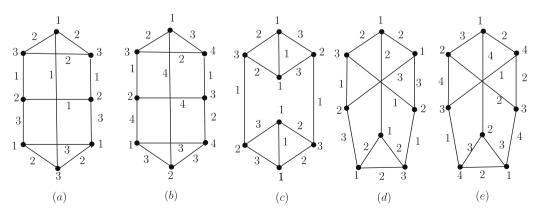


Figure 4: The total proper k-connected coloring of F<sub>1</sub>, F<sub>2</sub> and F<sub>3</sub>.

total proper paths connecting  $u_5$  and  $u_8$ . Thus,  $c(u_1u_5) \neq c(u_1u_8) \neq c(u_1)$ , and so  $c(u_1u_2) = c(u_1u_5)$ . But then, there is no set of three disjoint total proper paths connecting  $u_2$  and  $u_5$ , a contradiction. Hence,  $tpc_3(F_3) \ge 4$ . By Figure 4(e), we know that  $F_3$  is total proper 3-connected, and so  $tpc_3(F_3) = 4$ .

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