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#### Research Article

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# The optimal pebbling of spindle graphs

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**Abstract:** Given a distribution of pebbles on the vertices of a connected graph G, a *pebbling move* on G consists of taking two pebbles off one vertex and placing one on an adjacent vertex. The *optimal pebbling number* of G, denoted by  $\pi_{opt}(G)$ , is the smallest number m such that for some distribution of m pebbles on G, one pebble can be moved to any vertex of G by a sequence of pebbling moves. Let  $P_k$  be the path on K vertices. Snevily defined the N-K spindle graph as follows: take K copies of K and two extra vertices K and K, and then join the left endpoint (respectively, the right endpoint) of each K to K (respectively, K), the resulting graph is denoted by K0, and called the K1 spindle graph. In this paper, we determine the optimal pebbling number for spindle graphs.

Keywords: pebbling, optimal pebbling, spindle graph

MSC: 05C35

### 1. Introduction

*Graph pebbling* was first introduced into the literature by Chung (see [1]). Pebbling has developed its own subfield (see [2]). Let G be a simple graph with vertex set V(G) and edge set E(G). Let D be a distribution of pebbles on the vertices of G, or a *distribution* on G. For any vertex v of G, D(v) denotes the number of pebbles on v in D. For  $S \subseteq V(G)$ , we let  $D(S) = \sum_{v \in S} D(v)$  and  $|D| = \sum_{v \in V(G)} D(v)$ . A *pebbling move* consists of removing two pebbles from one vertex and then placing one pebble at an adjacent vertex. For  $v \in V(G)$ , v is *reachable* under distribution D if v has at least one pebble after some sequence of pebbling moves starting from D. A distribution D is *solvable* if all vertices of G are reachable under D. The *pebbling number* of G, denoted by  $\pi(G)$ , is the smallest number G0 solvable. The *optimal pebbling number* of G0, denoted by  $\pi_{opt}(G)$ 1, is the smallest number G2, denoted by G3 solvable. We say a distribution G4 is G5 pebbles and is solvable; that is, it is a solvable distribution of minimum size. We say a distribution G5 is G6 pebbles and is solvable on each vertex of degree 2. A vertex G6 is G6 under a distribution G6 if G7 if G8.

The optimal pebbling number of G was first introduced by Pachter, Snevily, and Voxman [3]. The optimal pebbling number has been determined for paths [3, 4], cycles [4], m-ary trees [5], caterpillars [6], and ladders [4]. Moews [7] used a probabilistic argument to show that the n-cube  $Q^n$  has  $\pi_{opt}(Q^n) = (4/3)^{n+O(\log n)}$ . In [8], Xue and Yerger investigated the optimal pebbling number of grids and found the optimal pebbling number for the 3 by n grid. For graphs of diameter two (respectively, three), Muntz et al. [9] characterized the classes of graphs having  $\pi_{opt}(G)$  equal to a value between 2 and 4 (respectively, between 3 and 8). The lower and upper bounds on the optimal pebbling number were further studied in [4]. Milans and Clark [10] showed that computing optimal pebbling number is NP-hard on arbitrary graphs. Interestingly, exact values for optimal

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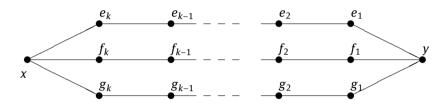
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pebbling number are known only for paths, cycles, caterpillars, *m*-ary trees, ladders, and the 3 by *n* grid. A survey of results of optimal pebbling number can be found in [2].

Let  $P_k$  be the path on k vertices. Snevily defined the n-k spindle graph as follows: take n copies of  $P_k$ and two extra vertices x and y, and then join the left endpoint (respectively, the right endpoint) of each  $P_k$ to x (respectively, y), the resulting graph is denoted by S(n, k), and called the n-k spindle graph. In fact, the spindle graph S(n, k) is the graph that the vertices x and y are connected by n internally-disjoint paths of length k + 1 (Figure 1 is the graph S(3, k)). Snevily and Foster [11] proposed the following Problem 1.1, which appears to be quite difficult.

**Problem 1.1** [11] Find  $\pi(S(3, k))$ .

By  $S(1, k) = P_{k+2}$  and  $S(2, k) = C_{2k+2}$ , we have  $\pi(S(1, k)) = 2^{k+1}$  and  $\pi(S(2, k)) = 2^{k+1}$  (see [4]). Recently, Gao and Yin [12] determined  $\pi(S(3, k))$ , which solves Problem 1.1.



**Figure 1:** S(3, k).

The focus of this paper is to investigate the optimal pebbling number of S(n, k). By  $\pi_{opt}(C_n) = \pi_{opt}(P_n) = \pi_{opt}(P_n)$  $\lceil \frac{2n}{3} \rceil$  (see [4]), we have  $\pi_{opt}(S(1,k)) = \lceil \frac{2k+4}{3} \rceil$  and  $\pi_{opt}(S(2,k)) = \lceil \frac{4k+4}{3} \rceil$ . For  $n \ge 3$ , we further determine  $\pi_{opt}(S(n, k))$  in this paper. That is the following Theorem 1.1.

**Theorem 1.1** Let  $n \ge 3$  and  $p \ge 0$ , and denote  $\ell = \max\{t | 2^t \le n\}$ .

- (1) If  $k < 2\ell$ , then  $\pi_{opt}(S(n, k)) = 2^{\lfloor k/2 \rfloor} + 2^{\lceil k/2 \rceil}$ ;
- (2) If  $k = 2\ell + 3p$ , then  $\pi_{opt}(S(n, k)) = 2^{\ell+1} + 2np$ ;
- (3) If  $k = 2\ell + 3p + 1$ , then  $\pi_{opt}(S(n, k)) = 2^{\ell+1} + 2^{\ell} + 2np$ ;
- (4) If  $k = 2\ell + 3p + 2$  and  $2n \ge 2^{\ell+1} + 2^{\ell-1}$ , then  $\pi_{ont}(S(n, k)) = 2^{\ell+2} + 2np$ ;
- (5) If  $k = 2\ell + 3p + 2$  and  $2n < 2^{\ell+1} + 2^{\ell-1}$ , then  $\pi_{opt}(S(n, k)) = 2^{\ell-1} + 2^{\ell} + 2n(p+1)$ .

#### 2. Lemmas

In order to prove Theorem 1.1, we need the following lemmas.

Let *D* be a distribution on *G*, and let  $H \subseteq G$ . The *restriction* of *D* to *H* is a pebble distribution  $D_H$  which is defined as follows:  $D_H(u) = D(u)$  if  $u \in V(H)$  and  $D_H(u) = 0$  if  $u \notin V(H)$ . For convenience, we write  $A_1, A_2, \ldots, A_n$  for the *n* copies of  $P_k$  in S(n, k). For  $1 \le i \le n$  and  $1 \le s < t \le k$ , we write  $A_i = u_{i,1}u_{i,2} \ldots u_{i,k}$ and  $A_i[s, t] = u_{i,s}u_{i,s+1}...u_{i,t}$ .

**Lemma 2.1.** [4] If G is a connected n-vertex graph, with  $n \ge 3$ , then G has a smooth optimal distribution with all leaves unoccupied.

**Lemma 2.2.** [3, 4] Let  $P_k$  be the path on k vertices. Then  $\pi_{opt}(P_k) = \lceil 2k/3 \rceil$ . The following Lemma gives an upper bound on  $\pi_{opt}(S(n, k))$ .

**Lemma 2.3.** Let  $n \ge 3$  and  $p \ge 0$ , and denote  $\ell = \max\{t | 2^t \le n\}$ . (1) If  $k < 2\ell$ , then  $\pi_{ont}(S(n, k)) \le 2^{\lfloor k/2 \rfloor} + 2^{\lceil k/2 \rceil}$ ;

- (2) If  $k = 2\ell + 3p$ , then  $\pi_{ont}(S(n, k)) \le 2^{\ell+1} + 2np$ ;
- (3) If  $k = 2\ell + 3p + 1$ , then  $\pi_{ont}(S(n, k)) \le 2^{\ell+1} + 2^{\ell} + 2np$ ;
- (4) If  $k = 2\ell + 3p + 2$  and  $2n \ge 2^{\ell+1} + 2^{\ell-1}$ , then  $\pi_{ont}(S(n, k)) \le 2^{\ell+2} + 2np$ ;
- (5) If  $k = 2\ell + 3p + 2$  and  $2n < 2^{\ell+1} + 2^{\ell-1}$ , then  $\pi_{opt}(S(n, k)) \le 2^{\ell-1} + 2^{\ell} + 2n(p+1)$ .

**Proof.** Clearly,  $2^{\ell} \le n < 2^{\ell+1}$  and  $\ell \ge 1$ .

- (1) Assume that  $k < 2\ell$ . Let D be a distribution such that  $D(x) = 2^{\lfloor k/2 \rfloor}$ ,  $D(y) = 2^{\lceil k/2 \rceil}$ , and D(v) = 0 for each  $v \in V(S(n,k)) \setminus \{x,y\}$ . Then D is solvable and  $\pi_{opt}(S(n,k)) \le 2^{\lfloor k/2 \rfloor} + 2^{\lceil k/2 \rceil}$ .
- (2) Assume that  $k = 2\ell + 3p$ . Let D be a distribution such that  $D(x) = D(y) = 2^{\ell}$ ,  $D(u_{i,j}) = 2$  for all  $i \in \{1, 2, ..., n\}$ ,  $j \in \{\ell + 2, ..., \ell + 3(p-1) + 2\}$ , and D(v) = 0 for each  $v \in V(S(n, k)) \setminus \{x, y, u_{1, \ell+2}, ..., u_{1, (\ell+3p-1)}, ..., u_{n, (\ell+2)}, ..., u_{n, (\ell+3p-1)}\}$ . It is clear to see that D is solvable. Then  $\pi_{opt}(S(n, k)) \le 2^{\ell+1} + 2np$ .
- (3) Assume that  $k = 2\ell + 3p + 1$ . Let D be a distribution such that  $D(x) = 2^{\ell}$ ,  $D(y) = 2^{\ell+1}$ ,  $D(u_{i,j}) = 2$  for all  $i \in \{1, 2, ..., n\}$ ,  $j \in \{\ell + 2, ..., \ell + 3(p-1) + 2\}$ , and D(v) = 0 for each  $v \in V(S(n, k)) \setminus \{x, y, u_{1,(\ell+2)}, ..., u_{1,(\ell+3p-1)}, ..., u_{n,(\ell+3p-1)}\}$ . Then D is solvable and  $\pi_{opt}(S(n, k)) \le 2^{\ell+1} + 2^{\ell} + 2nv$ .
- (4) Assume that  $k = 2\ell + 3p + 2$  and  $2n \ge 2^{\ell+1} + 2^{\ell-1}$ . Let D be a distribution such that  $D(x) = D(y) = 2^{\ell+1}$ ,  $D(u_{i,j}) = 2$  for all  $i \in \{1, 2, ..., n\}$ ,  $j \in \{\ell + 3, ..., \ell + 3p\}$ , and D(v) = 0 for each  $v \in V(S(n, k)) \setminus \{x, y, u_{1,(\ell+3)}, ..., u_{1,(\ell+3p)}, ..., u_{n,(\ell+3p)}\}$ . Then D is solvable and  $\pi_{opt}(S(n, k)) \le 2^{\ell+2} + 2np$ .
- (5) Assume that  $k = 2\ell + 3p + 2$  and  $2n < 2^{\ell+1} + 2^{\ell-1}$ . Let D be a distribution such that  $D(x) = 2^{\ell}$ ,  $D(y) = 2^{\ell-1}$ ,  $D(u_{i,j}) = 2$  for all  $i \in \{1, 2, \dots, n\}$ ,  $j \in \{\ell + 2, \dots, \ell + 3p + 2\}$ , and D(v) = 0 for each  $v \in V(S(n,k)) \setminus \{x,y,u_{1,(\ell+2)},\dots,u_{1,(\ell+3p+2)},\dots,u_{n,(\ell+2p+2)}\}$ . Then D is solvable and  $\pi_{opt}(S(n,k)) \le 2^{\ell-1} + 2^{\ell} + 2n(p+1)$ .  $\square$

Let

$$g(n,k) = \begin{cases} 2^{\lfloor k/2 \rfloor} + 2^{\lceil k/2 \rceil} & \text{if } k < 2\ell, \\ 2^{\ell+1} + 2np & \text{if } k = 2\ell + 3p, \\ 2^{\ell+1} + 2^{\ell} + 2np & \text{if } k = 2\ell + 3p + 1, \\ \min\{2^{\ell+2} + 2np, 2^{\ell-1} + 2^{\ell} + 2n(p+1)\} & \text{if } k = 2\ell + 3p + 2. \end{cases}$$

Then  $\pi_{opt}(S(n, k)) \leq g(n, k)$ .

**Lemma 2.4.** Assume that  $n \ge 3$ ,  $k \ge 2$ , and  $\alpha_1 \ge 0$ ,  $\alpha_2 \ge 0$ ,  $\alpha_1 + \alpha_2 \le k$ . Let D be a smooth solvable distribution on S(n, k). If there are at most  $2^{\alpha_1} + \beta_1$  pebbles can be put on x by a sequence of pebbling moves starting from D, and there are at most  $2^{\alpha_2} + \beta_2$  pebbles can be put on y by a sequence of pebbling moves, then  $|D| \ge 2^{\alpha_1} + 2^{\alpha_2} + n\lceil 2(k - \alpha_1 - \alpha_2)/3 \rceil$ , where  $0 \le \beta_i \le 2^{\alpha_i} - 1$  for  $i \in \{1, 2\}$ .

**Proof.** Clearly, for  $i \in \{1, 2, ..., n\}$ , we can move at most one pebble to x from  $A_i$  as D is a smooth distribution. Similarly, we can move at most one pebble to y from  $A_i$ . Consider a smooth solvable distribution  $D_0$  with  $D_0(x) = 2^{\alpha_1}$ ,  $D_0(y) = 2^{\alpha_2}$ .

If  $|D_0| \le 2^{\alpha_1} + 2^{\alpha_2} + n[2(k - \alpha_1 - \alpha_2)/3] - 1$ , then there are at most

$$|D_0| - D_0(x) - D_0(y) \le n\lceil 2(k - \alpha_1 - \alpha_2)/3 \rceil - 1$$

pebbles on  $A_1[1 + \alpha_1, k - \alpha_2], \ldots, A_n[1 + \alpha_1, k - \alpha_2]$ . For  $i \in \{2, \ldots, n\}$ , by Lemma 2.2, we need at least  $\lceil 2(k - \alpha_1 - \alpha_2)/3 \rceil$  pebbles that ensure each vertex in  $A_i[1 + \alpha_1, k - \alpha_2]$  to be reachable. Thus, there is some vertex in  $A_i[1 + \alpha_1, k - \alpha_2]$  which is not reachable. Hence,  $|D_0| \ge 2^{\alpha_1} + 2^{\alpha_2} + n\lceil 2(k - \alpha_1 - \alpha_2)/3 \rceil$ . Note that each pebbling move reduces the size of D. Therefore,  $|D| \ge |D_0| \ge 2^{\alpha_1} + 2^{\alpha_2} + n\lceil 2(k - \alpha_1 - \alpha_2)/3 \rceil$ .  $\square$ 

**Lemma 2.5.** *Let a be a nonnegative integer,* a = 3p + s*,*  $s \in \{0, 1, 2\}$ *. Then*  $2^{\lceil a/2 \rceil} + 2^{\lfloor a/2 \rfloor} \ge 4p + s + 2$ *.* 

**Proof.** Firstly, assume that a = 2b. Then p = (2b - s)/3, and  $2^{\lceil a/2 \rceil} + 2^{\lfloor a/2 \rfloor} = 2^{b+1}$ . Let  $f(t) = 2^{t+1} - 2 - s - (8t - 4s)/3$ . If a = 0, then b = p = s = 0, and  $f(0) \ge 0$ . If a = 2, then b = 1, p = 0, s = 2, and  $f(1) \ge 0$ . We have that  $f'(t) = 2^{t+1} \ln 2 - 8/3 > 0$  for  $t \ge 1$ . Hence,  $f(t) > f(1) \ge 0$ .

Now, we assume that a = 2b + 1. Then p = (2b + 1 - s)/3, and  $2^{\lceil a/2 \rceil} + 2^{\lfloor a/2 \rfloor} = 3 \times 2^b$ . If a = 1, then b = p = 0, s = 1, and  $2^{\lceil a/2 \rceil} + 2^{\lfloor a/2 \rfloor} \ge 4p + s + 2$ . Let  $f(t) = 3 \times 2^t - 2 - s - (8t + 4 - 4s)/3$ . If a = 3, then b = p = 1, s = 0, and  $f(1) \ge 0$ . We get  $f'(t) = 3 \times 2^t \ln 2 - 8/3 > 0$  for  $t \ge 1$ . Thus,  $f(t) > f(1) \ge 0$ . Therefore,  $2^{\lceil a/2 \rceil} + 2^{\lfloor a/2 \rfloor} \ge 4p + s + 2$ .  $\square$ 

#### 3. Proof of Theorem 1.1

**Proof of Theorem 1.1.** Clearly, S(n, 1) is isomorphic to  $K_{2,n}$ , and  $\pi_{opt}(K_{2,n}) = 3$ . We assume  $k \ge 2$ . By Lemma 2.3, we have that  $\pi_{opt}(S(n, k)) \le g(n, k)$ . We now show that  $\pi_{opt}(S(n, k)) \ge g(n, k)$ . By Lemma 2.1, we can assume that D is a smooth optimal distribution. Now, we assume that there are at most  $2^{\alpha_1} + \beta_1$  pebbles can be put on X by a sequence of pebbling moves starting from X0, and there are at most  $2^{\alpha_2} + \beta_2$  pebbles can be put on X1 by a sequence of pebbling moves, where X2 by X3 considering moves, where X4 by X5 considering moves, where X6 considering moves assume X6 considering moves, where X8 considering moves are X9 considering moves. We consider the following two cases.

**Case 1.**  $\alpha_1 + \alpha_2 > k$ .

Then  $|D| = D(x) + D(y) + \sum_{i=1}^{n} D(A_i) \ge 2^{\alpha_1} + \beta_1 - m_1 + 2^{\alpha_2} + \beta_2 - m_2 + m_1 + m_2 \ge 2^{\alpha_1} + 2^{\alpha_2}$ . Hence,  $|D| \ge 2^{\alpha_1} + 2^{\alpha_2} \ge 2^{\lceil (\alpha_1 + \alpha_2)/2 \rceil} + 2^{\lfloor (\alpha_1 + \alpha_2)/2 \rfloor} > 2^{\lceil k/2 \rceil} + 2^{\lfloor k/2 \rfloor}$  since  $2^a + 2^b \ge 2^{a-1} + 2^{b+1}$  for all integers satisfying a > b.

Assume that  $k = 2\ell + 3p + s$ ,  $s \in \{0, 1, 2\}$ . By Lemma 2.5, we have that

$$\begin{split} |D| &> 2^{\lceil k/2 \rceil} + 2^{\lfloor k/2 \rfloor} \\ &= 2^{\lceil (2\ell+3p+s)/2 \rceil} + 2^{\lfloor (2\ell+3p+s)/2 \rfloor} \\ &= 2^{\ell} (2^{\lceil (3p+s)/2 \rceil} + 2^{\lfloor (3p+s)/2 \rfloor}) \\ &\geq 2^{\ell} (4p+s+2). \end{split}$$

If s=0, then  $|D|>(4p+2)2^{\ell}>2^{\ell+1}+2np$ . If s=1, then  $|D|>4p\times2^{\ell}+3\times2^{\ell}>2^{\ell+1}+2^{\ell}+2np$ . If s=2, then  $|D|>4p\times2^{\ell}+4\times2^{\ell}>2^{\ell+2}+2np$   $\geq \min\{2^{\ell+2}+2np,2^{\ell-1}+2^{\ell}+2n(p+1)\}$ . Thus,  $|D|=\pi_{opt}(S(n,k))>g(n,k)$ . This is impossible as it would be greater than the known upper bound.

**Case 2.**  $\alpha_1 + \alpha_2 \le k$ .

Denote  $\omega = \alpha_1 + \alpha_2$ . By Lemma 2.4, we get

$$|D| \ge 2^{\alpha_1} + 2^{\alpha_2} + n\lceil 2(k - \omega)/3 \rceil \ge 2^{\lceil \omega/2 \rceil} + 2^{\lfloor \omega/2 \rfloor} + n\lceil 2(k - \omega)/3 \rceil \tag{3.1}$$

as  $2^a + 2^b \ge 2^{a-1} + 2^{b+1}$ . Now we consider the following three subcases.

**Subcase 2.1.**  $k < 2\ell$ .

If  $\omega = k$ , by (3.1), then  $|D| \ge 2^{\lceil k/2 \rceil} + 2^{\lfloor k/2 \rfloor}$ . If  $\omega = k - 1$ , by (3.1), then

$$|D| \ge 2^{\lceil \omega/2 \rceil} + 2^{\lfloor \omega/2 \rfloor} + n\lceil 2(k-\omega)/3 \rceil$$

$$\ge 2^{\lceil (k-1)/2 \rceil} + 2^{\lfloor (k-1)/2 \rfloor} + 2^{\ell}$$

$$\ge 2^{\lceil (k-1)/2 \rceil} + 2^{\lfloor (k-1)/2 \rfloor} + 2^{\lceil k/2 \rceil}$$

$$\ge 2^{\lceil k/2 \rceil} + 2^{\lfloor k/2 \rfloor}.$$

If  $\omega \leq k-2$ , then

$$|D| \ge 2^{\lceil \omega/2 \rceil} + 2^{\lfloor \omega/2 \rfloor} + n\lceil 2(k-\omega)/3 \rceil \ge 2^{\lceil \omega/2 \rceil} + 2^{\lfloor \omega/2 \rfloor} + 2 \times 2^{\ell} \ge 2^{\lceil k/2 \rceil} + 2^{\lfloor k/2 \rfloor}.$$

**Subcase 2.2.**  $k = 2\ell + 3p + s$  and  $0 \le \omega \le 2\ell + s$ , where  $s \in \{0, 1, 2\}$ .

Let  $\omega = 2\ell + s - j$ ,  $0 \le j \le 2\ell + s$ . By (3.1), we get

$$|D| \ge 2^{\lceil \omega/2 \rceil} + 2^{\lfloor \omega/2 \rfloor} + n \lceil (6p + 2j)/3 \rceil \ge 2^{\lceil (2\ell + s - j)/2 \rceil} + 2^{\lfloor (2\ell + s - j)/2 \rfloor} + 2np + n \lceil 2j/3 \rceil. \tag{3.2}$$

If s = 0 and  $0 \le j \le 1$ , by (3.2), then  $|D| \ge 2^{\ell+1} + 2np$ . If s = 0 and  $j \ge 2$ , then

$$|D| \geq 2^{\lceil (2\ell-j)/2 \rceil} + 2^{\lfloor (2\ell-j)/2 \rfloor} + 2np + 2^{\ell} \lceil 2j/3 \rceil > 2^{\ell+1} + 2np.$$

If s = 1 and  $j \le 3$ , then

$$|D| \ge 2^{\lceil (2\ell+1-j)/2 \rceil} + 2^{\lfloor (2\ell+1-j)/2 \rfloor} + 2np + 2^{\ell} \lceil 2j/3 \rceil \ge 2^{\ell+1} + 2^{\ell} + 2np.$$

If s = 1 and  $j \ge 4$ , then

$$|D| \ge 2^{\lceil (2\ell+1-j)/2 \rceil} + 2^{\lfloor (2\ell+1-j)/2 \rfloor} + 2np + 2^{\ell} \lceil 2j/3 \rceil > 2^{\ell+1} + 2^{\ell} + 2np.$$

Assume that s=2. If j=0, then  $|D|\geq 2^{\ell+2}+2np\geq \min\{2^{\ell+2}+2np,2^{\ell-1}+2^{\ell}+2n(p+1)\}$ . If j=1, then  $|D|\geq 2^{\ell+1}+2^{\ell}+2np+n\geq 2^{\ell+2}+2np\geq \min\{2^{\ell+2}+2np,2^{\ell-1}+2^{\ell}+2n(p+1)\}$ . If  $2\leq j\leq 3$ , then  $|D|\geq 2^{\ell}+2^{\ell-1}+2np+2n\geq \min\{2^{\ell+2}+2np,2^{\ell-1}+2^{\ell}+2n(p+1)\}$ . If j=4, then  $|D|\geq 2^{\ell-1}+2^{\ell-1}+2np+3n\geq 2^{\ell+2}+2np\geq \min\{2^{\ell+2}+2np,2^{\ell-1}+2^{\ell}+2n(p+1)\}$ . If  $j\geq 5$ , then  $|D|\geq 2np+4n\geq 2^{\ell+2}+2np\geq \min\{2^{\ell+2}+2np,2^{\ell-1}+2^{\ell}+2n(p+1)\}$ .

**Subcase 2.3.**  $k = 2\ell + 3p + s$  and  $2\ell + s < \omega \le 2\ell + 3p + s$ , where  $s \in \{0, 1, 2\}$ .

Note that  $p \ge 1$ . Let  $\omega = 2\ell + 3p + s - j$ ,  $0 \le j \le 3p - 1$ . By (3.1), we have

$$|D| \ge 2^{\ell} \left( 2^{\lceil (3p+s-j)/2 \rceil} + 2^{\lfloor (3p+s-j)/2 \rfloor} \right) + n\lceil 2j/3 \rceil. \tag{3.3}$$

Assume that 3p + s - j = 1. Then  $|D| \ge 3 \times 2^{\ell} + n\lceil 2j/3 \rceil = 2^{\ell+1} + 2^{\ell} + 2np + n\lceil (2s - 2)/3 \rceil$ . If s = 0, then  $|D| > 2^{\ell+1} + 2np$ . If s = 1, then  $|D| \ge 2^{\ell+1} + 2^{\ell} + 2np$ . If s = 2, then  $|D| \ge 2^{\ell+1} + 2^{\ell} + 2np + n > 2^{\ell+2} + 2np$ .

Assume that 3p + s - j = 2. Then  $|D| \ge 2 \times 2^{\ell+1} + n\lceil 2j/3 \rceil = 2^{\ell+2} + 2np + n\lceil (2s - 4)/3 \rceil$ . If s = 0, then  $|D| > 2^{\ell+1} + 2np$ . If s = 1, then  $|D| > 2^{\ell+1} + 2np$ . If s = 2, then  $|D| \ge 2^{\ell+2} + 2np$ .

Assume that 3p + s - j = 3. Then  $|D| \ge 6 \times 2^{\ell} + n\lceil 2j/3 \rceil = 2^{\ell+2} + 2^{\ell+1} + 2np + n\lceil (2s - 6)/3 \rceil$ . If s = 0, then  $|D| > 2^{\ell+1} + 2np$ . If s = 1, then  $|D| \ge 2^{\ell+1} + 2^{\ell} + 2np$ . If s = 2, then  $|D| > 2^{\ell+2} + 2np$ .

Assume that 3p + s - j = 4. Then  $|D| \ge 2^{\ell+3} + 2np + n\lceil (2s - 8)/3 \rceil$ . If s = 0, then  $|D| > 2^{\ell+1} + 2np$ . If s = 1, then  $|D| > 2^{\ell+1} + 2^{\ell} + 2np$ . If s = 2, then  $|D| > 2^{\ell+2} + 2np$ .

Assume that 3p + s - j = 2c and  $c \ge 3$ . By (3.3), and

$$p + s + \lceil 2j/3 \rceil - j \ge (j + 1 - s)/3 + s + \lceil 2j/3 \rceil - j = \lceil 2j/3 \rceil - 2j/3 + (1 + 2s)/3 > 0$$

we have that

$$\begin{split} |D| &\geq 2^{\ell} (2^{\lceil (3p+s-j)/2 \rceil} + 2^{\lfloor (3p+s-j)/2 \rfloor}) + n \lceil 2j/3 \rceil \\ &= 2^{\ell+1} 2^{c} + n \lceil 2j/3 \rceil \\ &\geq 2^{\ell+1} (2+2c) + n \lceil 2j/3 \rceil \\ &= 2^{\ell+2} + (3p+s-j)2^{\ell+1} + n \lceil 2j/3 \rceil \\ &> 2^{\ell+2} + (3p+s-j)n + n \lceil 2j/3 \rceil \\ &\geq 2^{\ell+2} + 2np + n(p+s+\lceil 2j/3 \rceil - j) \\ &> 2^{\ell+2} + 2np. \end{split}$$

Now, we assume that 3p + s - j = 2c + 1 and  $c \ge 2$ . By (3.3), and

$$\frac{p+3s+1-3j}{4} + \lceil 2j/3 \rceil \geq \frac{j+1-s}{12} + \frac{3s+1-3j}{4} + \lceil 2j/3 \rceil = \lceil 2j/3 \rceil - 2j/3 + \frac{1+2s}{3} > 0,$$

we have that

$$\begin{split} |D| & \geq 2^{\ell} (2^{\lceil (3p+s-j)/2 \rceil} + 2^{\lfloor (3p+s-j)/2 \rfloor}) + n\lceil 2j/3 \rceil \\ & = 3 \times 2^{\ell} \times 2^{c} + n\lceil 2j/3 \rceil \\ & \geq 3 \times 2^{\ell} (2+c) + n\lceil 2j/3 \rceil \\ & = 2^{\ell+2} + 2^{\ell+1} + 3c \times 2^{\ell} + n\lceil 2j/3 \rceil \\ & > 2^{\ell+2} + n(1 + 3c/2 + \lceil 2j/3 \rceil) \\ & > 2^{\ell+2} + n(1 + \frac{9p+3s-3j-3}{4} + \lceil 2j/3 \rceil) \\ & > 2^{\ell+2} + 2nn. \end{split}$$

Therefore,  $\pi_{ont}(S(n,k)) \ge g(n,k)$ . The proof of Theorem 1.1 is completed.  $\square$ 

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