Heuristics on pairing-friendly elliptic curves

John Boxall

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Abstract. We present a heuristic asymptotic formula as $x \to \infty$ for the number of isogeny classes of pairing-friendly elliptic curves over prime fields with fixed embedding degree $k \geq 3$, with fixed discriminant, with rho-value bounded by a fixed ρ_0 such that $1 < \rho_0 < 2$, and with prime subgroup order at most x.

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Introduction

Pairing-based cryptography protocols first became important with the work of Joux [18] and nowadays have numerous applications to the security of information transmission and other fields. Many of these protocols require the construction of elliptic curves over finite fields having very special properties. More precisely, let $q = p^f$ be a power of the prime p and let $k \geq 1$, $r \geq 1$ be integers. We need to be able to construct an elliptic curve E over the finite field $\mathbb{F}_q$ with q elements that satisfies the following:

(a) E has a point P of order r rational over $\mathbb{F}_q$.

(b) The group of points $E[r]$ of order r of E is isomorphic to $(\mathbb{Z}/r\mathbb{Z})^2$ and all the points of $E[r]$ are rational over the extension field $\mathbb{F}_{q^k}$ of degree k of $\mathbb{F}_q$.

In practical applications, if a security level of s bits is required, it is generally recommended that the integer r should have at least $2s$ bits (see for example [14, Table 1]). This is because the Pollard-rho algorithm is generally believed to be the best attack on the elliptic discrete logarithm problem. The subgroup of $E(\mathbb{F}_q)$ generated by P should be of small index in $E(\mathbb{F}_q)$. Since $\sharp(E(\mathbb{F}_q)) \in [(\sqrt{q}-1)^2, (\sqrt{q}+1)^2]$, so that $\sharp(E(\mathbb{F}_q)) \approx q$, a convenient measure of the suitability of the curve is the so-called rho-value, defined by $\rho = \frac{\log q}{\log r}$, which ideally should

be close to 1. On the other hand, the integer k needs to be sufficiently small to allow efficient arithmetic in $\mathbb{F}_{q^k}$, which in practice implies that k is at most about 50. These constraints on ρ and k imply very strong restrictions on the choice of elliptic curve, making suitable curves very rare [1, 15, 20, 24]. For this reason, a systematic search to obtain curves having parameters of cryptographic interest is completely out of the question.

Although there is considerable recent interest in protocols where the group order r is composite [5, 6, 13], we shall be concerned in this paper with the more familiar situation where r is a prime number, which is assumed to be the case from now on. Since known attacks on such protocols are based on the discrete logarithm in the subgroup of order r of the multiplicative group $\mathbb{F}_{q^k}^{\times}$, and this is believed to be the same difficulty as the discrete logarithm in $\mathbb{F}_{q^k}^{\times}$ itself, k cannot be too small. In what follows, therefore, we shall often suppose that $k \geq 3$.

Let E be an elliptic curve over $\mathbb{F}_q$ satisfying (a), where r is a prime different from p. Following what has become standard usage, the smallest integer k such that $q^k \equiv 1 \pmod{r}$ is called the embedding degree of (E, P) (or just of E if there is no possibility of confusion). Alternatively, the embedding degree is just the order of q in $(\mathbb{Z}/r\mathbb{Z})^{\times}$. An argument using the characteristic polynomial of the Frobenius endomorphism (see [1, Theorem 1]) shows that if E is an elliptic curve over $\mathbb{F}_q$ that satisfies (a) and if the embedding degree k of E is at least 2, then E also satisfies (b). Let $\Phi_k(w) \in \mathbb{Z}[w]$ denote the k-th cyclotomic polynomial. Then r divides $\Phi_k(q)$. On the other hand, if t denotes the trace of the Frobenius endomorphism of E over $\mathbb{F}_q$, then $\sharp(E(\mathbb{F}_q)) = q+1-t$ and so $q \equiv t-1 \pmod{r}$. It follows that r divides $\Phi_k(q)$ if and only if r divides $\Phi_k(t-1)$. Furthermore, we know from Hasse's bound that $|t| \leq 2\sqrt{q}$ and, if we suppose in addition that p does not divide t, then E is ordinary and there exists a unique square-free positive integer D and a unique integer $y > 0$ such that $t^2 + Dy^2 = 4q$. The endomorphism ring of E is then an order in the imaginary quadratic field $\mathbb{Q}(\sqrt{-D})$. Conversely, if t, D, y are integers, if $D > 0$ is square-free and if $t^2 + Dy^2 = 4q$ with $q = p^f$ a power of the prime p and p does not divide t, then a theorem of Deuring [11] implies that there exists an elliptic curve E over $\mathbb{F}_q$ such that $\sharp(E(\mathbb{F}_q)) = q + 1 - t$. If, further, r is a prime dividing both $q + 1 - t$ and $\Phi_k(t - 1)$, and if the rho-value $\frac{\log q}{\log r}$ is close to 1, then E is suitable for pairing-based cryptography. Since we only know how to construct the curve E corresponding to a choice of parameters (t, D, y) when D is fairly small ($D \leq 10^{15}$, say, see [12]), we shall suppose except in the last section that D is fixed.

The purpose of this note is to discuss the following heuristic asymptotic estimate.

Estimate 1 (Pairing-friendly curves estimate). *Let $k \geq 3$ be an integer, let $D \geq 1$ be a square-free integer and let $\rho_0 \in \mathbb{R}$ with $1 < \rho_0 < 2$. We suppose that*

(i) *$(k, D) \neq (3, 3)$, $(4, 1)$, $(6, 3)$;*

(ii) *if (k, D) is such that there exists a complete polynomial family (r_0, t_0, y_0) with generic rho-value equal to 1 (see remark (6) below and Section 3 for detailed definitions), then $\rho_0 > 1 + \frac{1}{\deg r_0}$.*

Let $e(k, D) = 2$ or 1 according as to whether $\sqrt{-D}$ belongs to the field generated over $\mathbb{Q}$ by the k-th roots of unity or not, let w_D be the number of roots of unity in the imaginary quadratic field $\mathbb{Q}(\sqrt{-D})$ and let h_D be the class number of $\mathbb{Q}(\sqrt{-D})$. Then the number of triples $(r, t, y) \in \mathbb{Z}^3$ with $2 \leq r \leq x$ a prime number dividing $\Phi_k(t - 1)$, $t^2 + Dy^2 = 4p$ with p prime, $y > 0$, r dividing $p + 1 - t$, and $p \leq r^{\rho_0}$ is asymptotically equivalent as $x \to \infty$ to

$$\frac{e(k, D)w_D}{2\rho_0 h_D} \int_2^x \frac{du}{u^{2-\rho_0}(\log u)^2}.$$

Several remarks are in order.

(1) If f is a function that is strictly positive for sufficiently large real x and if g is a second function defined for sufficiently large real x we say that g is asymptotically equivalent to f as $x \to \infty$ if $g(x) = f(x)(1 + o(1))$.

(2) Integrating by parts, we find

$$\int_2^x \frac{du}{u^{2-\rho_0}(\log u)^2} = \frac{1}{\rho_0 - 1} \frac{x^{\rho_0-1}}{(\log x)^2} \left(1 + O\left(\frac{1}{\log x}\right)\right),$$

where the constant implied by the O is independent of ρ_0. Thus, for fixed ρ_0, the number of triples is also asymptotically equivalent to

$$\frac{e(k, D)w_D}{2\rho_0(\rho_0 - 1)h_D} \frac{x^{\rho_0-1}}{(\log x)^2}. \tag{0.1}$$

However, in view of the term $\rho_0 - 1$ that appears in the denominator in this formula, the version with the integral seems preferable.

(3) Several papers have appeared in the literature showing (either heuristically or unconditionally) that pairing-friendly elliptic curves are sparse (see for example [1], [15, §4.1], [20] and [24] and also Remark 4.1). However, to the best of our knowledge, this paper is the first to suggest a possible asymptotic formula.

(4) When $q = p$ is prime and $p \geq 5$, Hasse's bound implies that p divides t if and only if $t = 0$. Since only finitely many primes divide $\Phi_k(-1)$, we must have $t \neq 0$ for all but finitely many triples, so that Deuring's results [11] imply that all

but finitely many triples (r, t, y) actually give rise to elliptic curves. Furthermore, the elliptic curves thus obtained are ordinary. On the other hand, a well-known result of Tate [23] asserts that two elliptic curves E_1 and E_2 over $\mathbb{F}_q$ are isogenous if and only if $\sharp(E_1(\mathbb{F}_q)) = \sharp(E_2(\mathbb{F}_q))$, and it is clear that the embedding degree k and the rho-value $\frac{\log p}{\log r}$ are invariant under isogeny. Thus Estimate 1 is also a heuristic asymptotic estimate for the number of isogeny classes of elliptic curves with given k and D defined over prime fields $\mathbb{F}_p$ and possessing a subgroup of prime order $r \leq x$ such that $p \leq r^{\rho_0}$. For given D, the methods of [12] construct curves whose endomorphism ring is the maximal order of $\mathbb{Q}(\sqrt{-D})$. On the other hand, [26, Theorem 6.1] shows that every isogeny class of ordinary elliptic curves contains a curve whose endomorphism ring is the maximal order of $\mathbb{Q}(\sqrt{-D})$. Thus, if D is sufficiently small, the methods of [12] enable one to construct at least one member of an isogeny class corresponding to any triple (r, t, y).

(5) We have supposed that $t^2 + Dy^2 = 4p$ with p prime rather than a power of a prime. However, as is usually the case in analytic number-theoretical situations, we expect solutions with $t^2 + Dy^2 = 4p^f$ and $f > 1$ to be negligible in number as compared with those with $f = 1$, so they should not in general affect the asymptotic estimate.

(6) We know of only one pair (k, D) for which there is a complete polynomial family (r_0, t_0, y_0) with generic rho-value equal to 1. This is the pair $(12, 3)$, and the corresponding family is the well-known Barreto–Naehrig family [2]. In this case the degree $\deg r_0$ of the polynomial r_0 is 4. In general, as we shall explain in Section 3, the Bateman and Horn heuristic asymptotic formula [3] predicts that a complete polynomial family with generic rho-value equal to one will produce more triples than predicted by Estimate 1 when $\rho_0 < 1 + \frac{1}{\deg r_0}$. This will be a consequence of Theorem 3.1 below.

(7) On the other hand, the cases $(k, D) = (3, 3), (6, 3)$ and $(4, 1)$ have to be excluded for a trivial reason. These are exactly the values of (k, D) with $k \geq 3$ and $\mathbb{Q}(\sqrt{-D})$ is equal to the field generated over $\mathbb{Q}$ by the k-th roots of unity; one deduces easily that $t^2 + Dy^2$ cannot be of the form $4p$ with p a prime. See Remark 1.3 for further details. Recall however that this does *not* imply that there are no pairing-friendly curves when (k, D) takes one of these values, but only that such curves cannot be rational over prime fields. Indeed, when $(k, D) = (3, 3)$, there is a well-known construction of curves over fields of square cardinality (see [14, §3.3] and also Remark 4.1 below).

(8) We have excluded the cases $k = 1$ and $k = 2$.

When $k = 1$ and E has a point P of order r rational over $\mathbb{F}_q$, there are two possibilities:

(a) either all the points of $E[r]$ are rational over $\mathbb{F}_q$, in which case $r^2 \leq q + 1 + 2\sqrt{q}$ by the Weil bound, which implies that the rho-value is asymptotically at least 2, or

(b) the points of $E[r]$ that are not multiples of P become rational only after extension of scalars to $\mathbb{F}_{q^r}$, so that computations of any sort are completely infeasible.

When $k = 2$ and E has a point P of order r rational over $\mathbb{F}_q$, then r divides $q + 1 - t$ and also r divides $q + 1$, since $\Phi_2(w) = w + 1$. Hence r divides t and again there are two possibilities:

(a) if $t \neq 0$, then $r \leq |t| \leq 2\sqrt{q}$ and so the rho-value is asymptotically at least 2, or

(b) $t = 0$, in which case E is supersingular. Suppose for example that the prime r is such that $2r - 1$ is also prime and take $q = p = 2r - 1$. By Deuring's results [11], there exists a supersingular elliptic curve E over $\mathbb{F}_p$ with $\sharp(E(\mathbb{F}_p)) = p + 1 = 2r$. By the Bateman–Horn heuristics, there is a constant $C > 0$ such that the number of primes $r \leq x$ with $2r - 1$ prime is asymptotically equal to $C \int_2^x \frac{du}{(\log u)^2}$. For the corresponding elliptic curves, the rho-value approaches 1 as $r \to \infty$. Thus, when $k = 2$, we expect far more pairing-friendly elliptic curves with $r \leq x$ than predicted by Estimate 1.

Here is a brief outline of the paper. In Section 1, we briefly describe a heuristic argument which leads to Estimate 1 and in Section 2 we present numerical evidence for several values of $(k, D) \neq (12, 3)$. In Section 3, we review families of pairing friendly curves and in particular the Barreto–Naehrig complete family [2], and explain why Estimate 1 is expected to fail when (k, D) satisfies condition (ii) of Estimate 1 and, in particular, when $(k, D) = (12, 3)$. This involves the Bateman-Horn heuristic asymptotic estimate on polynomials with integer coefficients and its generalization by K. Conrad [9] to polynomials with rational coefficients that take integer values. Finally, in Section 4, we briefly discuss a variant of Estimate 1 where D is allowed to vary and compare this with the recent work of Urroz, Luca and Shparlinski [24] (see Remark 4.1).

We insist on the fact that Estimate 1 is only a heuristic assertion, not a theorem. Indeed, proofs of most of the steps that are used to derive it and described in Section 1 seem to be a long way off.

All calculations reported on in this paper where done using PARI/GP [4] running on the GMP kernel [17] and often using PARI's GP to C compiler gp2c.

1   A heuristic argument

As in the Introduction, we fix an integer $k \geq 1$ and a square-free integer $D \geq 1$. If r is a prime such that r does not divide kD, $r \equiv 1 \pmod{k}$ and $-D$ is a square $\pmod{r}$, the Cocks–Pinch method [8], as explained say in [14, Theorem 4.1], produces all parameters (r, t, y) corresponding to ordinary curves with embedding degree k and endomorphism ring an order in $\mathbb{Q}(\sqrt{-D})$ having a point of order r. This means that r divides $\Phi_k(t-1)$, $y > 0$ and $t^2 + Dy^2 = 4p$ with p prime, the corresponding curve having coefficients in $\mathbb{F}_p$. As is well known, the rho-value of the curve is usually around 2. The heuristic argument that follows will give a measure of the frequency with which it can be expected to give curves with smaller rho-values.

As before, we fix a real number ρ_0 with $1 < \rho_0 < 2$. We wish to estimate asymptotically as $x \to \infty$ the number of triples $(r, t, y) \in \mathbb{Z}^3$ as above with $r \leq x$ and $p \leq r^{\rho_0}$. Thus, the heuristic argument that follows is, in fact, an estimate of the expected number of curves with $r \leq x$ and $p \leq r^{\rho_0}$ that the Cocks–Pinch method produces.

We first recall the following well-known lemmas, the first of which can be extracted from [25, Chapter 2, §2]:

Lemma 1.1. *Let $k \geq 1$ be an integer and let r be a prime number not dividing k. The following statements are equivalent.*

(i) *The cyclotomic polynomial $\Phi_k(w)$ has a root $\pmod{r}$.*

(ii) *$\Phi_k(w)$ splits into distinct linear factors $\pmod{r}$.*

(iii) *$r \equiv 1 \pmod{k}$.*

(iv) *r splits completely in the cyclotomic field $\mathbb{Q}(\zeta_k)$ generated over $\mathbb{Q}$ by a primitive k-th root of unity ζ_k.*

Lemma 1.2. *Let $K = \mathbb{Q}(\sqrt{-D})$ be an imaginary quadratic field. Let h_D be the class number of K and w_D the number of roots of unity of K. Then, as $X \to \infty$, the number of pairs of integers (t, y) with $t^2 + Dy^2 = 4p$, p a prime and $p \leq X$ is asymptotically equivalent to*

$$\frac{w_D X}{h_D \log X}.$$

Proof. Let t, y and p be as in the statement of the lemma and put

$$\pi = \frac{t + y\sqrt{-D}}{2}.$$

Then π is a root of $x^2 - tx + p$, so that π is an algebraic integer of K. Write $N(\alpha)$ for the norm down to $\mathbb{Q}$ of an element α of K. Then $N(\pi) = p$ so that the condition that p be prime is equivalent to the condition that π generate a principal prime ideal of K. By the prime ideal theorem in K (see for example [22, Chapter 7, §2]), the number of principal prime ideals $\mathfrak{p}$ of K of prime norm p bounded by X is equivalent to $\frac{X}{h_D \log X}$ as $X \to \infty$. On the other hand, every non-zero principal ideal of K has w_D generators all having the same norm. The lemma follows. $\quad\square$

We are now ready to explain our heuristic argument.

Step 1. Let $r \geq 2$ be any integer. Taking $X = r^{\rho_0}$ in Lemma 1.2, we find that the expected number of pairs (t, y) with $t^2 + Dy^2 = 4p$, p prime and $p \leq r^{r_0}$ is asymptotically equivalent to $\frac{w_D r^{\rho_0}}{h_D \rho_0 \log r}$.

Step 2. Let again $r \geq 2$ be any integer. By Lemma 1.1, the probability that r is prime and splits completely in $\mathbb{Q}(\zeta_k)$ is equal to the probability that r is prime and that $r \equiv 1 \pmod{k}$. Since there are $\phi(k)$ residue classes $\pmod{k}$ consisting of integers prime to k, the prime number theorem generalized to arithmetic progressions implies that the probability that r is prime and splits completely in $\mathbb{Q}(\zeta_k)$ is equal to $\frac{1}{\phi(k) \log r}$.

If, furthermore, t is the first member of a pair (t, y) as in Step 1, we assume that t is a sufficiently random integer for the probability that $\Phi_k(t - 1) \equiv 0 \pmod{r}$ to be equal to $\frac{\phi(k)}{r}$. Since $\Phi_1(w) = w + 1$ and $\Phi_2(w) = w - 1$, this is reasonable only when $k \geq 3$. Thus, when $k \geq 3$ and $r \geq 2$ are fixed integers, the expected number of triples (r, t, y) assuming r to be a prime dividing $\Phi_k(t - 1)$ and (t, y) as in Step 1 should be

$$\frac{1}{\phi(k) \log r} \frac{\phi(k)}{r} \frac{w_D r^{\rho_0}}{h_D \rho_0 \log r} = \frac{w_D r^{\rho_0}}{h_D \rho_0 r (\log r)^2}.$$

Step 3. We now estimate the probability that r divides $p + 1 - t$, given that r is prime. Now $p + 1 - t = N(\pi - 1)$, so that r divides $p + 1 - t$ if and only if there exists a prime ideal $\mathfrak{r}$ lying above r and dividing $\pi - 1$. Since $\rho_0 < 2$, this implies that r splits in $\mathbb{Q}(\sqrt{-D})$ as a product $\mathfrak{r}\bar{\mathfrak{r}}$ of two prime ideals of degree one. The probability that a random algebraic integer π satisfies $\pi \equiv 1 \pmod{\mathfrak{r}}$ is $\frac{1}{r}$ and the generalization to $\mathbb{Q}(\sqrt{-D})$ of Dirichlet's theorem on primes in arithmetic progressions implies that this remains true if π generates a prime ideal. Since there are two prime ideals $\mathfrak{r}$ and $\bar{\mathfrak{r}}$ dividing r, the probability that r divides $p + 1 - t$ given that it splits in $\mathbb{Q}(\sqrt{-D})$ is $\frac{2}{r}$.

On the other hand, the probability that r splits as a product of two degree one primes in $\mathbb{Q}(\sqrt{-D})$ is 1 if $\sqrt{-D} \in \mathbb{Q}(\zeta_k)$, and $\frac{1}{2}$ if not. This is equal to $\frac{e(k,D)}{2}$.

Step 4. Putting Steps 1–3 together, we conclude that if $r \geq 2$ is a fixed integer, the expected number of triples (r, t, y) with $t^2 + Dy^2 = 4p$, p a prime such that

$p \leq r^{\rho_0}$, and r prime and dividing both $\Phi_k(t-1)$ and $p+1-t$ is asymptotically equivalent to

$$\frac{e(k,D)w_D r_0^\rho}{\rho_0 h_D r^2 (\log r)^2} = \frac{e(k,D)w_D}{\rho_0 h_D r^{2-\rho_0}(\log r)^2}.$$

Summing over all integers r such that $2 \leq r \leq x$ and taking into account that

$$\sum_{2 \leq r \leq x} \frac{1}{r^{2-\rho_0}(\log r)^2} \sim \int_2^x \frac{du}{u^{2-\rho_0}(\log u)^2},$$

we obtain an estimate that differs by a factor of 2 from that in Estimate 1, the difference being due to the fact that we assumed in Estimate 1 that $y > 0$ whereas in the preceding argument the sign of y is arbitrary.

Remark 1.3. Step 3 assumes that π is an essentially random element of the set of algebraic integers of $\mathbb{Q}(\sqrt{-D})$ such that $\pi - 1$ belongs to one of the prime ideals dividing r. In particular, the probability that it generates a prime ideal should be that predicted by the prime ideal theorem. This is not true when $(k,D) = (3,3)$, $(6,3)$ or $(4,1)$, in other words in those cases where $\mathbb{Q}(\zeta_k) = \mathbb{Q}(\sqrt{-D})$. Suppose for example that $(k,D) = (3,3)$. The condition $r|\Phi_3(t-1)$ then implies that $4r$ divides $4t^2 - 4t + 4$. On the other hand, since $4r$ divides $(t-2)^2 + 3y^2 = t^2 - 4t + 4 + 3y^2$, we find by subtraction that $4r$ divides $3(t^2 - y^2)$. When $r \geq 5$, this implies that $t \equiv \pm y \pmod{4r}$. Since $|t| \leq 2r$ and $|y| \leq 2r$, this implies that $t = \pm y$ when r is sufficiently large and so $t^2 + 3y^2$ cannot be of the form $4p$ with p a prime. A similar argument works when $(k,D) = (6,3)$ or $(4,1)$. Thus the heuristic argument fails these cases.

2   Numerical evidence

In order to test Estimate 1 numerically, we wrote a programme in PARI/GP [4] to search for all triples (r,t,y) with r in some interval $[a,b]$, k, D and ρ_0 being given. Thus, for each prime $r \equiv 1 \pmod{k}$ belonging to $[a,b]$ such that $-D$ is a square $\pmod{r}$, the programme finds all the roots of $\Phi_k(t-1) \equiv 0 \pmod{r}$ and searches for those for which $|t| \leq 2r^{\frac{\rho_0}{2}}$. Given such a t, it tests whether there exists $y > 0$ such that $t^2 + Dy^2 = 4p$ with p prime and $p \leq r^{\rho_0}$. It outputs the vector of all sextuples (r,t,y,h,p,ρ) satisfying these conditions, with r, t, y and p as before, h the cofactor defined by $p + 1 - t = rh$, and $\rho = \frac{\log p}{\log r}$ the actual rho-value.

For a given r, there are two possible strategies for finding t. The first is to factor $\Phi_k(x) \pmod{r}$ using a standard factorization algorithm for univariate polynomials over finite fields. The second is to first choose at random a primitive root g

(mod r), so that if $s = g^{\frac{r-1}{k}}$ (mod r), then s is a primitive k-th root of unity in the field with r elements. The possible values of t are then $s^\ell + 1$ (mod r) as ℓ ranges over the integers between 1 and k that are prime to k. This is justified by the fact that the roots of Φ_k are precisely the primitive k-th roots of unity. In the range where the systematic search for all triples (r, t, y) is feasible, the second method turned out to be the faster although it is clear that for large values of r the first method is preferable since $k \leq 50$ and the exponentiation to the power $\frac{r-1}{k}$ becomes costly.

In view of the discussion in Section 1, our programme is basically an implementation of the Cocks–Pinch method that selects only those curves with $\rho \leq \rho_0$. However, as all primes $r \equiv 1$ (mod k) need to be tested, this cannot be expected in reasonable time to find curves in an interval $[a, b]$ where a and b are of a sufficiently large size for the curves to be of cryptographic interest (unless the value ρ is taken to be close to 2). In practice, it was found that for given k and D the vector of all sextuples (r, t, y, h, p, ρ) could be calculated in between 15 and 75 seconds when $b - a = 10^8$ and b was smaller than about 10^{15}. Under these conditions, the time taken was roughly proportional to $1/\phi(k)$. Also, in view of the irregularity that one expects when k and D vary and r is very small, it was decided to restrict attention to $r \geq 10^6$.

In what follows we present, for different values of k, D, ρ_0, a and b, the number $N = N(k, D, \rho_0, a, b)$ of triples (r, t, y) as in Estimate 1 with $a \leq r \leq b$ and, for comparison, the value of the corresponding integral

$$I = I(k, D, \rho_0, a, b) = \frac{e(k, D)w_D}{2\rho_0 h_D} \int_a^b \frac{du}{u^{2-\rho_0}(\log u)^2}. \tag{2.1}$$

We define I_0 by $I_0(k, D, \rho_0, a, b) = e(k, D)^{-1} I(k, D, \rho_0, a, b)$: note that I_0 depends only on D and ρ_0 but not on k.

Table 1 gives the values of $N(k, D, 1.7, 10^6, 85\,698\,768)$ for all k such that $3 \leq k \leq 30$ and all squarefree D with $D \leq 15$ as well as $D = 19, 23, 43$ and 47. This choice of D includes all imaginary quadratic fields of class number one except $\mathbb{Q}(\sqrt{-163})$ and, for each integer h less than or equal to 5, at least one field whose class number is equal to h. The second line of the table recalls the class number h_D of $\mathbb{Q}(\sqrt{-D})$. The third line gives the value of $I_0 = e(k, D)^{-1} I(k, D, 1.7, 10^6, 85\,698\,768)$. The values of I_0 are the reason for the choice of 85 698 768 as upper limit. In fact, when D is such that $w_D = 2$ and $h_D = 1$, then

$$I_0 = \frac{1}{1.7} \int_{10^6}^{85\,698\,768} \frac{du}{u^{0.3}(\log u)^2} \simeq 1000.00$$

so that the predicted value of $N(k, D, 1.7, 10^6, 85\,698\,768)$ is 1000 in these cases. The main part of the table contains the values of $N(k, D, 1.7, 10^6, 85\,698\,768)$, the entries corresponding to values of (k, D) with $e(k, D) = 2$ are marked with an asterisk; Estimate 1 predicts that they should be close to $2I_0$ and therefore roughly twice as large as the other entries in the same column. The last line of Table 1 gives the average value of each column as k varies from 3 to 30, the cases where $e(k, D) = 2$ being counted with weight $\frac{1}{2}$ and the excluded values $(k, D) = (3, 3)$, $(4, 1)$ and $(6, 3)$ omitted. Estimate 1 predicts that each of these averages be close to the corresponding value of I_0.

Table 2 gives the values of $N(k, D, 1.5, 10^6, 2 \times 10^8)$ for the same values of (k, D) as Table 1. When D is such that $w_D = 2$ and $h_D = 1$, we now have

$$I_0 = \frac{1}{1.5} \int_{10^6}^{2 \cdot 10^8} \frac{du}{u^{0.5}(\log u)^2} \simeq 58.17.$$

Although all the entries in Tables 1 and 2 (with the exception of those for $(k, D) = (3, 3)$, $(4, 1)$ and $(6, 3)$) are of the order of magnitude predicted by Estimate 1, there is considerable variation in the actual values, especially in Table 2. This is perhaps not unexpected, as similar variation occurs when one computes the number of values for which polynomials simultaneously take prime values and compares the result to the Bateman–Horn heuristics. In fact, if $\pi(x)$ denotes as usual the number of primes less than or equal to the real positive x, no explicit formula analogous to Riemann's formula for $\pi(x) - \int_2^x \frac{du}{\log u}$ seems to be known in the Bateman–Horn context (see for example [19] for a discussion of the case of prime pairs). So, presumably it would also be a hard problem to find one in the context of Estimate 1.

In order to obtain numerical data for larger values of x and examine what happens when ρ_0 varies, it is necessary to restrict the values of k and D. The case $(k, D) = (12, 3)$ will be discussed in the next section. Table 3 presents data for the three cases $(k, D) = (28, 1)$, $(27, 11)$ and $(8, 23)$. In each case, the table gives the values of $N(\rho_0) = N(k, D, \rho_0, a, b)$ and $I(\rho_0) = I(k, D, \rho_0, a, b)$ for $\rho_0 \in \{1.1, 1.2, 1.3, 1.4, 1.5\}$ and for each of the three intervals $(a, b) = (10^6, 10^8)$, $(10^8, 10^{10})$ and $(10^{12} - 10^{10}, 10^{12} + 10^{10})$. These results emphasize just how rare triples with rho-values close to one are. For example, if one wanted to construct a table like Table 1 with $I_0 = 1000$ but taking $\rho_0 = 1.2$ instead of 1.7, Estimate 1 suggests that one would need to test all r up to about 7.9×10^{29}, which is obviously completely out of the question.

D	1	2	3	5	6	7	10	11	13	14	15	19	23	43	47
h_D	1	1	1	2	2	1	2	1	2	4	2	1	3	1	5
I_0	2000	1000	3000	500	500	1000	500	1000	500	250	500	1000	333.3	1000	200
$k=3$	2087	1053	0†	534	512	1012	514	1049	512	246	529	1049	362	991	195
4	0†	998	3132	568	568	1033	515	1066	510	282	507	1085	328	992	220
5	2193	1001	3219	513	544	963	552	1079	510	271	507	1004	345	1066	194
6	2118	1008	0†	535	517	1049	497	1032	521	261	509	1088	323	1044	209
7	2107	1024	3112	533	517	2098*	512	1047	530	270	533	1061	346	1036	208
8	4226*	2117*	3115	505	520	1018	510	1039	507	249	515	1056	338	1062	174
9	2120	1014	6139*	484	503	1041	507	984	512	228	549	1077	329	1060	191
10	2167	1039	3171	492	536	995	509	1038	539	267	523	990	347	1029	195
11	2064	1033	3121	518	489	1009	447	2084*	524	264	537	1035	345	1069	205
12	4239*	1048	6368*	519	547	1009	518	1055	502	259	519	1030	334	1078	205
13	1970	1065	3061	544	504	988	476	1059	521	229	526	1076	333	1028	192
14	2095	1102	3243	560	546	2001*	540	1023	532	278	533	1048	364	999	225
15	2030	981	6221*	526	516	1130	525	982	502	289	975*	1058	347	1077	191
16	4183*	2058*	3007	528	536	1071	502	998	511	260	491	1001	361	1071	205
17	2073	1008	3194	517	506	1023	509	1015	482	254	470	1096	374	1020	206
18	2139	1017	6215*	534	512	1013	537	1021	558	273	520	1016	334	1001	207
19	2073	1031	3115	529	564	1049	497	1048	566	229	518	2127*	356	1025	205

D	1	2	3	5	6	7	10	11	13	14	15	19	23	43	47
20	4063*	1071	3111	1073*	517	1039	502	1096	481	234	491	1028	325	1101	196
21	2035	1068	6304*	526	509	2016*	500	995	568	293	503	1060	371	1019	199
22	2145	996	3048	557	512	1042	533	2138*	519	239	545	1059	345	988	216
23	2113	1012	3185	530	521	1043	476	1071	492	271	527	1059	682*	1064	219
24	4161*	2110*	6247*	510	1055*	1003	543	996	529	260	525	1031	333	1113	214
25	1971	1102	3082	499	504	1031	481	1038	540	248	523	996	374	997	227
26	2065	1055	3230	493	525	1058	542	1042	530	257	541	1083	336	1071	196
27	2148	1049	6327*	483	521	1035	516	1062	503	270	541	976	323	1053	179
28	4189*	1038	3119	547	514	2047*	513	1042	506	268	480	1006	367	1054	197
29	2153	979	3017	581	509	1072	551	1040	522	263	500	1030	334	1086	201
30	2153	1041	6198*	494	535	1029	519	1030	534	271	996*	1068	361	955	211
Avg	2094.4	1034.8	3126.6	524.8	522.6	1029.9	513.3	1037.8	520.4	260.1	516.0	1043.9	345.6	1041.0	202.9

Table 1. Values of $N(k, D, 1.7, 10^6, 85\,698\,768)$ for $3 \leq k \leq 30$ and various D. An asterisk denotes entries corresponding to values of (k, D) with $e(k, D) = 2$. A dagger indicates the three exceptional cases $(k, D) = (3, 3)$, $(4, 1)$, $(6, 3)$ when no triples (r, t, y) can exist. See Section 2 for detailed explanations.

D	1	2	3	5	6	7	10	11	13	14	15	19	23	43	47
h_D	1	1	1	2	2	1	2	1	2	4	2	1	3	1	5
I_0	116.3	58.17	174.5	29.09	29.09	58.17	29.09	58.17	29.09	14.54	29.09	58.17	19.39	58.17	11.63
$k=3$	132	69	0†	29	34	57	35	54	29	14	27	59	17	54	12
4	0†	63	198	20	31	65	31	65	27	17	37	64	22	59	10
5	123	49	211	31	26	55	24	53	30	18	26	45	21	73	12
6	132	58	0†	36	41	61	22	61	32	10	29	63	14	56	13
7	111	59	190	34	32	119*	29	67	32	21	27	75	15	63	6
8	235*	131*	181	30	26	56	27	47	34	16	30	64	9	61	9
9	132	60	367*	31	27	52	32	63	34	22	32	80	18	52	6
10	118	55	205	28	33	69	39	59	38	13	37	46	15	66	10
11	111	64	197	31	38	58	26	119*	29	17	28	58	15	59	13
12	255*	42	419*	22	21	62	30	67	25	27	28	61	15	59	16
13	125	66	164	21	27	37	26	61	43	20	32	51	28	58	9
14	122	74	168	29	35	133*	29	45	31	13	32	55	14	69	16
15	119	59	381*	32	30	64	28	57	30	19	57*	58	16	61	9
16	244*	130*	193	30	32	58	33	53	28	9	27	71	18	77	15
17	133	62	194	32	33	60	22	55	30	10	36	78	16	66	11
18	133	59	316*	34	36	65	32	62	33	18	23	63	15	71	11
19	111	64	176	36	27	53	31	46	38	18	32	127*	24	63	15

D	1	2	3	5	6	7	10	11	13	14	15	19	23	43	47
20	249*	60	176	64*	31	73	27	57	28	12	30	63	21	61	9
21	113	66	378*	26	25	114*	26	51	33	18	30	60	25	57	12
22	123	62	184	25	34	55	30	127*	36	19	29	68	17	54	15
23	103	61	192	30	44	53	38	71	32	24	17	60	44*	71	13
24	207*	129*	343*	28	48*	64	25	69	26	14	40	60	15	51	15
25	96	65	186	40	26	60	33	79	34	12	28	67	20	57	10
26	144	57	173	33	35	66	36	65	31	14	32	45	18	59	11
27	135	51	354*	44	40	59	27	76	21	17	17	62	27	56	10
28	266*	66	220	25	30	123*	31	66	31	19	34	65	23	71	11
29	113	69	170	34	23	69	29	60	26	21	25	69	23	43	12
30	109	67	388*	24	37	47	26	55	29	13	69*	47	25	55	12
Avg	121.0	61.50	186.6	30.25	31.36	59.38	29.43	60.25	31.07	16.61	29.57	61.45	18.86	60.79	11.54

Table 2. Values of $N(k, D, 1.5, 10^6, 2 \times 10^8)$ for $3 \leq k \leq 30$ and various D. An asterisk denotes entries corresponding to values of (k, D) with $e(k, D) = 2$. A dagger indicates the three exceptional cases $(k, D) = (3, 3), (4, 1), (6, 3)$ when no triples (r, t, y) can exist. See Section 2 for detailed explanations.

Data for $k = 28$, $D = 1$, $\rho_0 \in \{1.1, 1.2, 1.3, 1.4, 1.5\}$.

Interval	$10^6 \le r \le 10^8$		$10^8 \le r \le 10^{10}$		$10^{12} - 10^{10} \le r \le 10^{12} + 10^{10}$	
ρ_0	$I(\rho_0)$	$N(\rho_0)$	$I(\rho_0)$	$N(\rho_0)$	$I(\rho_0)$	$N(\rho_0)$
1.1	0.325	0	0.311	1	0.002	0
1.2	1.502	3	2.286	6	0.022	0
1.3	7.104	8	17.22	24	0.321	0
1.4	34.39	37	132.71	135	4.723	5
1.5	170.07	188	1044.7	1128	69.86	73

Data for $k = 27$, $D = 11$, $\rho_0 \in \{1.1, 1.2, 1.3, 1.4, 1.5\}$.

Interval	$10^6 \le r \le 10^8$		$10^8 \le r \le 10^{10}$		$10^{12} - 10^{10} \le r \le 10^{12} + 10^{10}$	
ρ_0	$I(\rho_0)$	$N(\rho_0)$	$I(\rho_0)$	$N(\rho_0)$	$I(\rho_0)$	$N(\rho_0)$
1.1	0.081	0	0.078	0	0.00038	0
1.2	0.375	0	0.57	2	0.0055	0
1.3	1.78	1	4.31	5	0.080	0
1.4	8.60	9	33.18	30	1.18	1
1.5	42.52	57	261.17	271	17.46	22

Data for $k = 8$, $D = 23$, $\rho_0 \in \{1.1, 1.2, 1.3, 1.4, 1.5\}$.

Interval	$10^6 \le r \le 10^8$		$10^8 \le r \le 10^{10}$		$10^{12} - 10^{10} \le r \le 10^{12} + 10^{10}$	
ρ_0	$I(\rho_0)$	$N(\rho_0)$	$I(\rho_0)$	$N(\rho_0)$	$I(\rho_0)$	$N(\rho_0)$
1.1	0.027	0	0.026	0	0.00013	0
1.2	0.125	0	0.191	0	0.00183	0
1.3	0.592	0	1.435	1	0.0267	0
1.4	2.866	1	11.06	16	0.394	0
1.5	14.47	7	87.06	76	5.821	6

Table 3. Results for $(k, D) = (28, 1)$, $(27, 11)$ and $(8, 23)$. Here, $I(\rho_0)$ is an abbreviation for $I(k, D, \rho_0, a, b)$ as defined in (2.1) with $(a, b) = (10^6, 10^8)$, $(10^8, 10^{10})$ or $(10^{12} - 10^{10}, 10^{12} + 10^{10})$. Similarly, $N(\rho_0)$ is an abbreviation for $N(k, D, \rho_0, a, b)$, the number of triples (r, t, y) as in Estimate 1 with $a \le r \le b$.

3 The Barreto–Naehrig family and the case $k = 12$, $D = 3$

The various known methods of constructing pairing-friendly elliptic curves are reviewed in [14]. Since Estimate 1 is primarily concerned with ordinary elliptic curves over prime fields and assumes that $k \geq 3$, we limit our attention to those methods which apply in these situations. We want to understand asymptotically as $x \to \infty$ the number of triples (r, t, y) with $r \leq x$ that belong to such families and have rho-value at most ρ_0 and compare this with Estimate 1. Clearly we can only compare constructions where k and D are fixed.

Apart from the Cocks–Pinch method, which constructs all parameters corresponding to ordinary curves and on which our heuristic estimate is based, the other well-known constructions with k and D fixed are the polynomial families. These fall into two kinds: (a) sparse families, of which the most familiar example is MNT families [21]; (b) complete families, of which the general construction is due to Brezing and Weng [7]. We refer to [14, §5, §6] for a detailed review of the two kinds of families.

The idea behind both constructions is to find polynomials $r_0(w)$, $t_0(w)$ and $p_0(w) \in \mathbb{Q}[w]$ such that $r_0(w)$ divides both $\Phi_k(t_0(w) - 1)$ and $p_0(w) + 1 - t_0(w)$. One then seeks values w_0 of w for which $r_0(w_0)$, $t_0(w_0)$ and $p_0(w_0)$ are all integers with $r_0(w_0)$ prime (or a prime multiplied by a very small factor) and $p_0(w_0)$ is prime (or a prime power). The values of the integral parameters r, t and p are then respectively $r_0(w_0)$, $t_0(w_0)$ and $p_0(w_0)$ with $r_0(w_0)$ and $p_0(w_0)$ prime. By definition, the generic rho value of the family is $\frac{\deg p_0}{\deg r_0}$. As w_0 tends to infinity, the rho-value of the elliptic curve corresponding w_0 approaches the generic rho-value.

However, the two constructions differ in the way they treat the parameter y. Define the polynomial $h_0(w)$ by $p_0(w) + 1 - t_0(w) = r_0(w)h_0(w)$. If $r = r_0(w_0)$, $t = t_0(w_0)$, $p = p_0(w_0)$ and $h = h_0(w_0)$, then the corresponding y parameter satisfies

$$Dy^2 = 4p - t^2 = 4hr - (t - 2)^2.$$

In the case of sparse families, the general idea is to choose r_0, t_0 and p_0 in such a way that $4p_0(w) - t_0(w)^2$ is of degree two. When this is the case, the affine curve with (w, y)-equation $Dy^2 = 4p_0(w) - t_0(w)^2$ is of genus 0. If this curve is to have infinitely many integral points, its real locus must be either a parabola or a hyperbola. In all the cases of which we are aware, the real locus is a hyperbola. Thus, an affine change of coordinates transforms this into a generalized Pell equation $Z^2 - aY^2 = b$, with $a > 0$ is not a square. The integral solutions of this are of the form $Z + \sqrt{a}Y = \alpha\varepsilon^n$, where α runs through a finite set of elements of the real quadratic field $\mathbb{Q}(\sqrt{a})$, ε is a fundamental unit of $\mathbb{Q}(\sqrt{a})$, and $n \in \mathbb{Z}$. From this we deduce that the number of values of $r \leq x$ that can

arise from a sparse family is $O((\log x)^2)$. On the other hand, Estimate 1 predicts that there are at least $\gg \frac{x^{\rho_0-1}}{(\log x)^2}$ choices of the parameters (r, t, y, p) with $r \leq x$ and $p \leq r^{\rho_0}$. Thus, sparse families can only contribute a negligible proportion of pairing friendly-curves with given k and D.

In the case of complete families, the basic strategy was described in full generality by Brezing and Weng [7]. In addition to r_0, t_0, h_0 and p_0, we also require a polynomial y_0 such that $t_0(w)^2 + Dy_0(w)^2 = 4p_0(w)$, so that the y parameter is the corresponding value $y_0(w_0)$. Now, the polynomials r_0, t_0, y_0, h_0, p_0 simultaneously take integral values at integers w_0 varying over a finite set of congruence classes modulo some fixed integer. Furthermore, if r_0 and p_0 are to give rise to triples (r, t, y) corresponding to elliptic curves, they must simultaneously take prime values.

Before going further, we recall the Bateman–Horn heuristics [3] in the case of two polynomials f and g with integral coefficients. We assume that f and g are distinct and irreducible. For any prime p let N_p denote the number of solutions of the congruence $f(x)g(x) \equiv 0 \pmod{p}$ and suppose that $N_p < p$ for all p. Then let C be given by the conditionally convergent infinite product

$$C = \prod_{p \geq 2 \text{ prime}} \left(1 - \frac{N_p}{p}\right)\left(1 - \frac{1}{p}\right)^{-2}. \tag{3.1}$$

Then the number of integers w_0 with $2 \leq w_0 \leq X$ such that $f(w_0)$ and $g(w_0)$ are simultaneously prime is asymptotically equivalent to

$$\frac{C}{\deg r_0 \deg p_0} \int_2^X \frac{du}{(\log u)^2} \tag{3.2}$$

as $X \to \infty$. In particular, since $C > 0$, there are infinitely many integers w_0 such that $f(w_0)$ and $g(w_0)$ are simultaneously prime.

We need to adapt this statement to polynomials whose coefficients are rational. Let $f, g \in \mathbb{Q}[w]$ and let $n \geq 1$ be a common denominator of the coefficients of f and g. Then there are integers m_i with $0 \leq m_i < n$ such that $f(nw_0 + m_i) \in \mathbb{Z}$ and $g(nw_0 + m_i) \in \mathbb{Z}$ for all i and for all $w_0 \in \mathbb{Z}$.

Then, for each i, we can apply the generalization by K. Conrad (see [9, §2]) of the Bateman–Horn heuristics to the pair of polynomials $w \mapsto f(nw + m_i)$ and $w \mapsto f(nw + m_i)$. This implies that (3.2) still holds, although the value of C will no longer be given by (3.1) in general, but can be computed using [9, Conjecture 5]. Since in what follows we only need the actual value of C in the case of polynomials with integer coefficients, we do not discuss this in detail.

Returning to our discussion of complete families, it follows that there exists a constant $C' > 0$ such that the number of triples (r, t, y) with $r \leq x$ coming from

the family is asymptotically equivalent to

$$\frac{C'}{\deg r_0 \deg p_0} \int_2^{(x/c_{r_0})^{1/\deg r_0}} \frac{du}{(\log u)^2} \sim \frac{C' \deg r_0}{c_{r_0}^{1/\deg r_0} \deg p_0} \frac{x^{1/\deg r_0}}{(\log x)^2}, \qquad (3.3)$$

where c_{r_0} is the leading coefficient of r_0 and $\deg r_0$ is the degree of r_0, and the asymptotic equivalence of the two displayed formulae is seen by integrating by parts. (Note that in general C' will not be equal to C, since both positive and negative values of w_0 may yield triples (r, t, y).)

As $w_0 \to \infty$, the rho-value of the triple $(r_0(w_0), t_0(w_0), y_0(w_0))$ approaches $\frac{\deg p_0}{\deg r_0}$. Comparing (0.1) and (3.3), we deduce that if $\frac{1}{\deg r_0} > \rho_0 - 1$, then the Bateman–Horn heuristics implies the complete family parametrized by (r_0, t_0, y_0) asymptotically contains more choices of triples than predicted by Estimate 1. It also follows that this family can contain infinitely many triples with rho-value $\leq \rho_0$ only if $\frac{\deg p_0}{\deg r_0} \leq \rho_0$. It is clear that $\deg p_0 \geq \deg r_0$ so, since $\deg p_0$ and $\deg r_0$ are integers, the conditions

$$\frac{\deg p_0}{\deg r_0} \leq \rho_0 \quad \text{and} \quad \frac{1}{\deg r_0} > \rho_0 - 1$$

are satisfied only if $\deg p_0 = \deg r_0$. We deduce (i) of the following theorem.

Theorem 3.1. *We keep the notation that has just been introduced and assume the Bateman–Horn heuristics together with their generalization by K. Conrad.*

(i) *Suppose that $\rho_0 < 1 + \frac{1}{\deg r_0}$. Then the complete family (r_0, t_0, y_0) asymptotically contains more choices of parameters than predicted by Estimate 1. Furthermore, one has $\deg p_0 = \deg r_0$.*

(ii) *On the other hand, if $\rho_0 > 1 + \frac{1}{\deg r_0}$ then the family does not contain sufficiently many triples to contradict Estimate 1.*

Point (ii) is proved in a similar way to (i), again comparing of (0.1) and (3.3).

On the other hand, what happens when $\rho_0 = 1 + \frac{1}{\deg r_0}$ depends on the relative values of the constants appearing in (0.1) and the right hand side of (3.3).

Table 8.2 of [14] summarizes, for all k up to 50, the construction of the family with the smallest rho-value and the corresponding value of D. When $k \geq 4$, the families listed are all complete families, and all have $\deg p_0 > \deg r_0$ *except when* $k = 12$, in which case the corresponding value of D is 3. When $k = 3$, the family is also a complete family and $D = 3$ and also satisfies $\deg p_0 = \deg r_0$, except that $p_0(w) = (3w - 1)^2$ cannot represent primes (see [14, §3.3]).

The case $k = 12$ and $D = 3$ is thus expected to provide a genuine counterexample to Estimate 1. The corresponding family is the well-known Barreto–Naehrig family [2], where

$$r_0(w) = 36w^4 + 36w^3 + 18w^2 + 6w + 1, \quad t_0(w) = 6w^2 + 1, \quad h_0(w) = 1,$$
$$y_0(w) = 6w^2 + 4w + 1, \quad p_0(w) = 36w^4 + 36w^3 + 24w^2 + 6w + 1.$$

Since the degree of r_0 is 4, we expect the family to provide more curves than Estimate 1 when $\rho_0 < 1.25$.

This can be tested numerically using similar calculations to those presented in Section 1. To see the contribution of the Barreto–Naehrig family, we need to calculate the constant C appearing in the Bateman–Horn heuristics for it. For any prime p, let $N_{r_0,p}$ denote the number of solutions of $r_0(w) \equiv 0 \pmod{p}$ and define $N_{p_0,p}$ analogously. Write N_p for the number of solutions of $r_0(w)p_0(w) \equiv 0 \pmod{p}$. Then $N_2 = N_3 = 0$. Furthermore, when $p \geq 5$, we deduce from the relation $p_0(w) = r_0(w) + 6w^2$ that r_0 and p_0 cannot have a common root $\pmod{p}$, so that $N_p = N_{r_0,p} + N_{p_0,p}$. Since r_0 and p_0 have integral coefficients, the Bateman–Horn constant is given by (3.1).

As written, the product (3.1) is conditionally convergent and therefore unsuitable for numerical computation. Instead, we apply the formula given by the theorem of Davenport and Schinzel [10]. This gives

$$C = \frac{\gamma}{\rho(K_{r_0})\rho(K_{p_0})} \prod_{p \geq 5} \left(1 - \frac{N_p}{p}\right)\left(1 - \frac{1}{p}\right)^{-N_p} \prod_{p \geq 5} \left(1 - \frac{1}{p^2}\right)^{-N_p^{(2)}} \left(1 - \frac{1}{p^4}\right)^{-N_p^{(4)}},$$

where $N_p^{(2)}$ and $N_p^{(4)}$ denote respectively the number of irreducible factors of $r_0(x)p_0(x) \pmod{p}$ of degree 2 and of degree 4, $\rho(K_{r_0})$ and $\rho(K_{p_0})$ the residue at 1 of the zeta function of the number fields K_{r_0} and K_{p_0} generated over $\mathbb{Q}$ respectively by a root of r_0 and a root of p_0 and

$$\gamma = \left(1 - \frac{1}{2^2}\right)^{-2}\left(1 - \frac{1}{3^2}\right)^{-1}\left(1 - \frac{1}{3}\right)^{-1} = 3.$$

The two infinite products in the Davenport–Schinzel formula for C are now absolutely convergent. When $p \geq 5$, Table 4 gives the value of $N_p^{(j)}$ when $j = 2$ and $j = 4$.

Using these formulae and taking the product over all p with $5 \leq p \leq 10^6$, we find that the first product appearing in the formula for C equals $0.88576\ldots$ and the second equals $1.26250\ldots$. On the other hand, $\rho(K_{r_0}) = 0.36105\ldots$ and $\rho(K_{p_0}) = 0.52642\ldots$. It follows that $C \simeq 17.651$. Since neither of the

$p \bmod 12$	$p_0(w) \bmod p$	N_p	$N_p^{(2)}$	$N_p^{(4)}$
1	4 roots	8	0	0
1	0 roots	4	2	0
5		0	2	1
7		2	3	0
11		0	4	0

Table 4

polynomials r_0 and p_0 are even functions, the values of $r_0(w_0)$ and $p_0(w_0)$ at negative integers w_0 will, with finitely many exceptions, be different to those at positive integers. Hence $C' = 2C$ so that $\frac{C'}{16} \simeq 2.206$ and, if the Bateman–Horn heuristics are correct, we can expect the number of triples (r, t, y) arising from the Barreto–Naehrig family with $x' \leq r \leq x$ should be approximately equal to

$$J_{\mathrm{BN}}(x', x) = 2.206 \int_{x'^{1/4}/\sqrt{6}}^{x^{1/4}/\sqrt{6}} \frac{du}{(\log u)^2}.$$

Table 5 gives the values of $N(12, 3, \rho_0, 10^6, 10^8)$ together with $N(12, 3, \rho_0, 10^8, 10^{10})$ for $\rho_0 \in \{1.1, 1.2, 1.3, 1.4, 1.5\}$ and compares them with the corresponding expected value of $I(12, 3, \rho_0, a, b)$.

$\rho_0 =$	1.1	1.2	1.3	1.4	1.5
$N(10^6, 10^8)$	3	8	21	57	305
$I(10^6, 10^8)$	0.49	2.25	10.66	51.58	255.11
$N(10^8, 10^{10})$	6	10	44	221	1655
$I(10^8, 10^{10})$	0.47	3.43	25.83	199.07	1567.0

Table 5

The column $\rho_0 = 1.1$ of Table 5 shows three triples with $10^6 \leq r \leq 10^8$ and six with $10^8 \leq r \leq 10^{10}$. All these nine triples (r, t, y) are in fact members of the Barreto–Naehrig family: they correspond to the values of the polynomials $r_0(x)$ etc. at $x = -107, -55, -52, -41, -15, 20, 78, 82, 123$. This should be compared with the expected contributions from the Barreto–Naehrig family which are respectively $J_{\mathrm{BN}}(10^6, 10^8) = 6.05$ and $J_{\mathrm{BN}}(10^8, 10^{10}) = 10.26$. Although there is some discrepancy between the computed and expected values, this is perhaps not surprising in view of our previous remarks concerning the variation between computed and expected values in the Bateman–Horn heuristics.

4 What happens when D varies

Let again D denote a square-free positive integer. As before, we denote the discriminant of the imaginary quadratic field $\mathbb{Q}(\sqrt{-D})$ by d_D, thus $d_D = -D$ if $D \equiv 3 \pmod 4$ and $d_D = -4D$ if $D \equiv 1, 2 \pmod 4$. If z is small with respect to x, Estimate 1 suggests that the number of triples (r, t, y) as above with $r \leq x$, $p \leq r^{\rho_0}$ and $|d_D| \leq z$ should be equivalent to

$$\sum_{|d_D| \leq z} \frac{e(k, D) w_D}{2\rho_0 h_D} \int_2^x \frac{du}{u^{2-\rho_0}(\log u)^2}.$$

Here we shall not try to give a precise meaning to the condition that z be small with respect to x, which would require a discussion of the error term in Estimate 1 which would take us too far afield. We content ourselves with a heuristic asymptotic estimate for the sum

$$\sum_{|d_D| \leq z} \frac{e(k, D) w_D}{2\rho_0 h_D}$$

as $z \to \infty$. It is well known that $\sqrt{-D} \in \mathbb{Q}(\zeta_k)$ if and only if d_D divides k. Furthermore, $w_D = 2$ except when $D = 1$ or $D = 3$. Therefore

$$\sum_{|d_D| \leq z} \frac{e(k, D) w_D}{2\rho_0 h_D} = \frac{1}{\rho_0} \sum_{|d_D| \leq z} \frac{1}{h_D} + O(1),$$

where the constant implied by the $O(1)$ depends only on k. Estimates for the sum $\sum_{|d_D| \leq z} h_D^\alpha$ for various positive values of α, and in particular $\alpha = 1$, have been studied since the time of Gauss (see for example [16] and the references cited therein). However, we have been unable to find any reference to the case $\alpha = -1$ which is of interest here. On the other hand, heuristic considerations involving the prime ideal theorem and the residue of zeta functions at $s = 1$ for imaginary quadratic fields suggest that

$$\sum_{|d_D| \leq z} \frac{1}{h_D} \sim \frac{6}{\pi} \sqrt{z}, \quad z \to \infty$$

and this seems to be confirmed by numerical calculation. This suggest the following heuristic:

Estimate 2 (Variable D estimate). *Let $k \geq 3$ and ρ_0 such that $1 < \rho_0 < 2$ be fixed. If z is small with respect to x, then, as $x \to \infty$ the number $\mathcal{N}(k, z, \rho_0, x)$ of triples (r, t, y) as in Estimate 1 with $|d_D| \leq z$ is equivalent to*

$$\frac{6}{\rho_0 \pi} \sqrt{z} \int_2^x \frac{du}{u^{2-\rho_0}(\log u)^2}.$$

In particular, if we can take $z = x^\alpha$ for some small positive α then, integrating by parts, we find that the number of triples (r, t, y) with $r \leq x$ and $|d_D| \leq x^\alpha$ should be equivalent to

$$\frac{6}{\rho_0(\rho_0 - 1)\pi} \frac{x^{\frac{\alpha}{2}+\rho_0-1}}{(\log x)^2}.$$

At present it is not quite clear how large we can take α for this estimate to be reasonable. This depends in particular on the size of the error term in Estimate 1, a problem which certainly deserves study but we prefer to leave this for future work. One reason for this is that, to the best of our knowledge, no detailed discussion of the error term in the Bateman–Horn heuristics has appeared in the literature up till now.

Remark 4.1. In [24], Urroz, Luca and Shparlinski prove a result which implies an unconditional upper bound on $\mathcal{N}(k, z, \rho_0, x)$. In fact, their Theorem 1 implies that

$$\mathcal{N}(k, z, \rho_0, x) \ll \phi(k)\big(x^{\rho_0-1} + x^{\frac{\rho_0}{2}}\big)z^{\frac{1}{2}}\frac{\log x}{\log\log x} \ll \phi(k)x^{\frac{\rho_0}{2}}z^{\frac{1}{2}}\frac{\log x}{\log\log x},$$

where the constants implied by the $\ll$ are absolute. This follows from the hypothesis that $1 < \rho_0 < 2$, the variable x of [24] corresponds to our x^{ρ_0}, the y of [24] to our x, and the z of [24] is contained between $\frac{1}{4}z$ and z when z is used in our sense. For constant z, this is much weaker than Estimate 1, but when $(k, D) = (3, 3)$ there exists the complete family

$$r_0(w) = 9w^2 - 3w + 1, \quad t_0(w) = -3w + 1, \quad y_0(w) = 3w - 1,$$
$$h_0(w) = 1, \quad q_0(w) = (3w - 1)^2,$$

together with a similar family with $r_0(w) = 9w^2 - 9w + 3$ (see [14, §3.3]). The Bateman–Horn heuristics therefore imply that

$$\mathcal{N}(3, z, \rho_0, x) \gg \frac{x^{\frac{1}{2}}}{(\log x)^2}$$

for any $z \geq 3$ and any ρ_0. A similar argument using the Barreto–Naehrig family suggests that

$$\mathcal{N}(12, z, \rho_0, x) \gg \frac{x^{\frac{1}{4}}}{(\log x)^2}$$

for any $z \geq 3$ and any ρ_0. Thus, the Urroz–Luca–Shparlinski upper bound for a given k is strongly related to the existence of complete families with rho-value 1 for at least one value of D.

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Author information

John Boxall, Laboratoire de Mathématiques Nicolas Oresme, CNRS – UMR 6139,
Université de Caen Basse-Normandie, Boulevard Maréchal Juin, BP 5186,
14032 Caen cedex, France.
E-mail: `john.boxall@unicaen.fr`