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## WHEN IS A BCC-ALGEBRA EQUIVALENT TO AN MV-ALGEBRA?

**Abstract.** The aim of this paper is to characterize BCC-algebras which are term equivalent to MV-algebras. It turns out that they are just the bounded commutative BCC-algebras. Further, we characterize congruence kernels as deductive systems. The explicit description of a principal deductive system enables us to prove that every subdirectly irreducible bounded commutative BCC-algebra is a chain (with respect to the induced order).

### 1. Introduction

By a **BCC-algebra** we mean an algebra  $\mathcal{A} = (A; \rightarrow, 1)$  of type  $(2, 0)$  satisfying the following axioms

$$(BCC1) \quad (x \rightarrow y) \rightarrow ((z \rightarrow x) \rightarrow (z \rightarrow y)) = 1;$$

$$(BCC2) \quad x \rightarrow x = 1;$$

$$(BCC3) \quad x \rightarrow 1 = 1;$$

$$(BCC4) \quad 1 \rightarrow x = x;$$

$$(BCC5) \quad (x \rightarrow y = 1 \text{ and } y \rightarrow x = 1) \text{ implies } x = y.$$

These algebras were introduced by Y. Komori [9] in connection with the problem whether the class of all BCK-algebras forms a variety. The problem was solved in the negative.

Let us note that for a BCC-algebra  $\mathcal{A}$  the relation defined by

$$(*) \quad x \leq y \quad \text{if and only if} \quad x \rightarrow y = 1$$

is an order on  $A$  with greatest element 1, see e.g. [9]. Due to this fact, the

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identity (BCC1) can be read as

$$x \rightarrow y \leq (z \rightarrow x) \rightarrow (z \rightarrow y).$$

This equivalent formulation will be used in our paper.

We can prove the following

LEMMA 1. *Let  $\mathcal{A} = (A; \rightarrow, 1)$  be a BCC-algebra. Then*

- (i)  $x \leq y$  implies  $z \rightarrow x \leq z \rightarrow y$ ;
- (ii)  $x \leq y$  implies  $y \rightarrow z \leq x \rightarrow z$ ;
- (iii)  $y \leq x \rightarrow y$ .

**Proof.** Suppose  $x \leq y$ . Then  $x \rightarrow y = 1$  and, by (BCC1),

$$\begin{aligned} 1 &= (x \rightarrow y) \rightarrow ((z \rightarrow x) \rightarrow (z \rightarrow y)) = 1 \rightarrow ((z \rightarrow x) \rightarrow (z \rightarrow y)) = \\ &= (z \rightarrow x) \rightarrow (z \rightarrow y) \end{aligned}$$

thus  $z \rightarrow x \leq z \rightarrow y$  proving (i). Similarly,

$$\begin{aligned} 1 &= (y \rightarrow z) \rightarrow ((x \rightarrow y) \rightarrow (x \rightarrow z)) = (y \rightarrow z) \rightarrow (1 \rightarrow (x \rightarrow z)) = \\ &= (y \rightarrow z) \rightarrow (x \rightarrow z) \end{aligned}$$

whence  $y \rightarrow z \leq x \rightarrow z$  proving (ii).

Applying (ii) and the fact that  $x \leq 1$  for each  $x \in A$  we conclude  $y = 1 \rightarrow y \leq x \rightarrow y$ . ■

The concept of a **BCK-algebra** was introduced by K. Iséki and Y. Imai [7] as an algebra  $\mathcal{A} = (A; \rightarrow, 1)$  of type  $(2, 0)$  satisfying the following axioms

- (BCK1)  $(x \rightarrow y) \rightarrow ((y \rightarrow z) \rightarrow (x \rightarrow z)) = 1$ ;
- (BCK2)  $x \rightarrow ((x \rightarrow y) \rightarrow y) = 1$ ;
- (BCK3)  $x \rightarrow x = 1$ ;
- (BCK4)  $x \rightarrow 1 = 1$ ;
- (BCK5)  $(x \rightarrow y = 1 \text{ and } y \rightarrow x = 1) \text{ implies } x = y$ .

Moreover, a BCK-algebra  $\mathcal{A}$  is called **commutative** if it satisfies the so-called **commutative law** (see [1] for this notation)

$$(C) \quad (x \rightarrow y) \rightarrow y = (y \rightarrow x) \rightarrow x.$$

It is well-known (see [8], [10]) that if  $\mathcal{A}$  is a commutative BCK-algebra then it is a  $\vee$ -semilattice where  $x \vee y = (x \rightarrow y) \rightarrow y$ .

Analogously as for BCC-algebras, the relation  $\leq$  defined by  $(*)$  is an order on the support of a BCK-algebra  $\mathcal{A}$  and 1 is the greatest element.

The following three lemmas are known but their proofs were published in a different way in several hardly attainable papers. Thus we present our proofs for the reader's convenience.

LEMMA 2. Let  $\mathcal{A} = (A; \rightarrow, 1)$  be a BCK-algebra. Then

- (i)  $1 \rightarrow x = x$ ;
- (ii)  $x \leq y$  implies  $y \rightarrow z \leq x \rightarrow z$ ;
- (iii)  $x \leq y$  implies  $z \rightarrow x \leq z \rightarrow y$ ;
- (iv)  $y \leq x \rightarrow y$ .

Proof. (i) Using of (BCK3) and (BCK2), we get  $1 \rightarrow x = (x \rightarrow x) \rightarrow x \geq x$ . However,  $(1 \rightarrow x) \rightarrow x = ((x \rightarrow x) \rightarrow x) \rightarrow x \geq x \rightarrow x = 1$  thus  $1 = (1 \rightarrow x) \rightarrow x$  whence  $1 \rightarrow x \leq x$ . Together we have  $1 \rightarrow x = x$ .

(ii) Suppose  $x \leq y$ . Then  $x \rightarrow y = 1$  and, by (i) and (BCK1),

$$\begin{aligned} 1 &= (x \rightarrow y) \rightarrow ((y \rightarrow z) \rightarrow (x \rightarrow z)) = 1 \rightarrow ((y \rightarrow z) \rightarrow (x \rightarrow z)) = \\ &= (y \rightarrow z) \rightarrow (x \rightarrow z) \end{aligned}$$

giving  $y \rightarrow z \leq x \rightarrow z$ .

(iii) If  $x \leq y$  then by (i) and (BCK1) we derive

$$\begin{aligned} 1 &= (z \rightarrow x) \rightarrow ((x \rightarrow y) \rightarrow (z \rightarrow y)) = (z \rightarrow x) \rightarrow (1 \rightarrow (z \rightarrow y)) = \\ &= (z \rightarrow x) \rightarrow (z \rightarrow y) \end{aligned}$$

proving  $z \rightarrow x \leq z \rightarrow y$ .

(iv) Since  $x \leq 1$  by (BCK4), we apply (i) and (ii):  $y = 1 \rightarrow y \leq x \rightarrow y$ . ■

LEMMA 3. Every BCK-algebra satisfies the so-called **exchange identity**

(EI)  $x \rightarrow (y \rightarrow z) = y \rightarrow (x \rightarrow z)$ .

Proof. Substituting  $y$  by  $y \rightarrow z$  in (BCK1), we get

$$x \rightarrow (y \rightarrow z) \leq ((y \rightarrow z) \rightarrow z) \rightarrow (x \rightarrow z).$$

By (BCK2) we have  $y \leq (y \rightarrow z) \rightarrow z$ , thus, by Lemma 2 (ii),

$$((y \rightarrow z) \rightarrow z) \rightarrow (x \rightarrow z) \leq y \rightarrow (x \rightarrow z).$$

Together it yields  $x \rightarrow (y \rightarrow z) \leq y \rightarrow (x \rightarrow z)$ . Swapping  $x, y$ , we obtain the converse inequality. ■

LEMMA 4. Every BCK-algebra is a BCC-algebra. A BCC-algebra is a BCK-algebra if and only if it satisfies the exchange identity (EI).

Proof. To prove the first assertion, we need to verify only (BCC1). By Lemma 3, a BCK-algebra satisfies (EI) thus, using this and (BCK1), we compute

$$(x \rightarrow y) \rightarrow ((z \rightarrow x) \rightarrow (z \rightarrow y)) = (z \rightarrow x) \rightarrow ((x \rightarrow y) \rightarrow (z \rightarrow y)) = 1.$$

Conversely, let a BCC-algebra  $\mathcal{A}$  satisfy (EI). We need to verify only (BCK1) and (BCK2). By (BCC1) and (EI) we have

$$1 = (y \rightarrow z) \rightarrow ((x \rightarrow y) \rightarrow (x \rightarrow z)) = (x \rightarrow y) \rightarrow ((y \rightarrow z) \rightarrow (x \rightarrow z)),$$

which is (BCK1), and

$$1 = (x \rightarrow y) \rightarrow (x \rightarrow y) = x \rightarrow ((x \rightarrow y) \rightarrow y),$$

which is (BCK2). ■

We can apply the previous lemmas to state

**THEOREM 1.** *Every BCC-algebra satisfying (BCK2) is a BCK-algebra.*

**Proof.** By (BCK2) we have  $y \leq (y \rightarrow z) \rightarrow z$ . Applying (BCC1) we compute

$$y \leq (y \rightarrow z) \rightarrow z \leq (x \rightarrow (y \rightarrow z)) \rightarrow (x \rightarrow z)$$

thus, by Lemma 1 (ii) and (BCK2),

$$y \rightarrow (x \rightarrow z) \geq ((x \rightarrow (y \rightarrow z)) \rightarrow (x \rightarrow z)) \rightarrow (x \rightarrow z) \geq x \rightarrow (y \rightarrow z).$$

Interchanging the roles of  $x$  and  $y$ , we obtain the converse inequality proving the exchange identity (EI). By Lemma 4, we have shown that the given BCC-algebra is in fact a BCK-algebra. ■

## 2. Bounded BCC-algebras

We say that a BCC-algebra  $\mathcal{A}$  is **bounded** if it has a least element 0, i.e. if  $0 \leq x$  for each  $x \in A$ . Clearly, this property can be characterized by the identity

$$(Z) \quad 0 \rightarrow x = 1$$

and such an algebra will be denoted by  $\mathcal{A} = (A; \rightarrow, 1, 0)$  to indicate the existence of a new nullary operation explicitly.

A bounded BCC-algebra  $\mathcal{A}$  satisfies the **double negation law** if the identity

$$(DN) \quad (x \rightarrow 0) \rightarrow 0 = x$$

holds in  $\mathcal{A}$ .

For the sake of brevity, we will denote  $x \rightarrow 0$  by  $\neg x$  and call it the **negation** of  $x$ . Hence, (DN) can be read as

$$\neg\neg x = x.$$

The concept of an MV-algebra was introduced by C.C. Chang [4] as an axiomatization of the Łukasiewicz many-valued logic. We present the definition taken from the monograph [5]:

By an **MV-algebra** we mean an algebra  $\mathcal{M} = (M; \oplus, \neg, 0)$  of type  $(2, 1, 0)$  satisfying the following identities

- (MV1)  $x \oplus (y \oplus z) = (x \oplus y) \oplus z$ ;
- (MV2)  $x \oplus y = y \oplus x$ ;
- (MV3)  $x \oplus 0 = x$ ;
- (MV4)  $\neg\neg x = x$ ;
- (MV5)  $x \oplus \neg 0 = \neg 0$  ( $\neg 0$  is denoted by 1);
- (MV6)  $\neg(\neg x \oplus y) \oplus y = \neg(\neg y \oplus x) \oplus x$ .

The following result was proved by D. Mundici [10]:

**PROPOSITION.** *Let  $\mathcal{M} = (M; \oplus, \neg, 0)$  be an MV-algebra. Define  $x \rightarrow y = \neg x \oplus y$  and  $1 = \neg 0$ . Then  $\mathcal{A}(M) = (M; \rightarrow, 1, 0)$  is a bounded commutative BCK-algebra.*

*Let  $\mathcal{A} = (A; \rightarrow, 1, 0)$  be a bounded commutative BCK-algebra. Define  $x \oplus y = (x \rightarrow 0) \rightarrow y$  and  $\neg x = x \rightarrow 0$ . Then  $\mathcal{M}(A) = (A; \oplus, \neg, 0)$  is an MV-algebra.*

In the sequel, we are going to modify the Proposition for BCC-algebras. At first we prove

**LEMMA 5.** *Every bounded BCK-algebra satisfying the double negation law satisfies the contraposition law*

$$(CL) \quad x \rightarrow y = \neg y \rightarrow \neg x.$$

**Proof.**  $\neg y \rightarrow \neg x = (y \rightarrow 0) \rightarrow (x \rightarrow 0) = x \rightarrow ((y \rightarrow 0) \rightarrow 0) = x \rightarrow \neg\neg y = x \rightarrow y$ . ■

**REMARK.** Of course, the contraposition law entails the double negation law since

$$x = 1 \rightarrow x = \neg x \rightarrow \neg 1 = \neg x \rightarrow 0 = \neg\neg x.$$

It is well-known that commutative BCK-algebras form a variety which can be axiomatized by the following identities

- (i)  $(x \rightarrow y) \rightarrow y = (y \rightarrow x) \rightarrow x$ ;
- (ii)  $x \rightarrow (y \rightarrow z) = y \rightarrow (x \rightarrow z)$ ;
- (iii)  $x \rightarrow x = 1$ ;
- (iv)  $1 \rightarrow x = x$ .

Just as in case of BCK-algebras, we say that a BCC-algebra  $(A; \rightarrow, 1)$  is **commutative** if it satisfies the identity (C).

It easily follows that these algebras are precisely the commutative BCK-algebras:

**THEOREM 2.** *Every commutative BCC-algebra  $(A; \rightarrow, 1)$  is a commutative BCK-algebra.*

**Proof.** By Lemma 1 (iii) and (C) we have

$$x \leq (y \rightarrow x) \rightarrow x = (x \rightarrow y) \rightarrow y.$$

Thus  $(A; \rightarrow, 1)$  satisfies (BCK2), which by Theorem 1 entails that it is a BCK-algebra. ■

**COROLLARY.** *Let  $(A; \rightarrow, 1, 0)$  be a bounded commutative BCC-algebra. Define*

$$x \oplus y = (x \rightarrow 0) \rightarrow y \quad \text{and} \quad \neg x = x \rightarrow 0.$$

*Then  $(A; \oplus, \neg, 0)$  is an MV-algebra.*

**Proof.** This is an immediate consequence of the previous Theorem 2 and Mundici's proposition. ■

### 3. Congruence kernels

Let  $\mathcal{A} = (A; \rightarrow, 1)$  be a BCC-algebra and  $\Theta \in \text{Con}\mathcal{A}$ . The congruence class  $[1]_\Theta$  is called the **kernel** of  $\Theta$ .

Congruence kernels of BCC-algebras are in a close relationship with congruences and hence it is important to have their characterization.

Let  $\mathcal{A} = (A; \rightarrow, 1)$  be a BCC-algebra. A subset  $D \subseteq A$  is called a **deductive system** of  $\mathcal{A}$  if

- (a)  $1 \in D$ ;
- (b) if  $a \in D$  and  $a \rightarrow b \in D$  then also  $b \in D$ .

Of course, the condition (b) is in fact the deduction rule **Modus Ponens** which justifies the name “deductive system”.

However, the class of BCC-algebras is not closed under homomorphic images thus there is not a one-to-one correspondence between congruences and their kernels. This correspondence exists only for the so-called relative congruences, i.e. such  $\Theta \in \text{Con}\mathcal{A}$  that  $\mathcal{A}/\Theta$  is a BCC-algebra again.

Further, note that every deductive system of a BCC-algebra  $\mathcal{A}$  is an order filter (i.e. an upset) with respect to the induced order.

**THEOREM 3.** *Let  $\mathcal{A} = (A; \rightarrow, 1, 0)$  be a bounded BCC-algebra satisfying the contraposition law (CL). Then  $D \subseteq A$  is a congruence kernel if and only if  $D$  is a deductive system of  $\mathcal{A}$ .*

**Proof.** Obviously, if  $D = [1]_\Theta$  for some  $\Theta \in \text{Con}\mathcal{A}$ ,  $a \in D$  and  $a \rightarrow b \in D$  then  $\langle a, 1 \rangle \in \Theta$ ,  $\langle a \rightarrow b, 1 \rangle \in \Theta$  thus also  $\langle a \rightarrow b, b \rangle = \langle a \rightarrow b, 1 \rightarrow b \rangle \in \Theta$  and, due to transitivity of  $\Theta$ ,  $\langle b, 1 \rangle \in \Theta$  proving  $b \in D$ . Thus  $D$  is a deductive system.

To prove the converse, we only need to show that the relation defined by

$$\langle x, y \rangle \in \Theta_D \text{ if and only if } x \rightarrow y, y \rightarrow x \in D$$

is a congruence on  $\mathcal{A}$  whenever  $D$  is a deductive system. It is clear that then  $[1]_{\Theta_D} = D$ .

Suppose that  $D$  is a deductive system of  $\mathcal{A}$ . By definition,  $\Theta_D$  is a reflexive and symmetric binary relation on  $A$ . Assume  $\langle x, y \rangle \in \Theta_D$  and  $\langle y, z \rangle \in \Theta_D$ . Then  $x \rightarrow y, y \rightarrow x, y \rightarrow z, z \rightarrow y \in D$  and, by (BCC1),

$$(y \rightarrow x) \rightarrow ((z \rightarrow y) \rightarrow (z \rightarrow x)) = 1 \in D.$$

Since  $y \rightarrow x \in D$  and  $z \rightarrow y \in D$ , we conclude  $z \rightarrow x \in D$ . Analogously, it can be shown  $x \rightarrow z \in D$ , thus  $\langle x, z \rangle \in \Theta_D$  proving that  $\Theta_D$  is an equivalence of  $A$ . It remains to show that  $\Theta_D$  has the Substitution Property with respect to  $\rightarrow$ .

Since  $\mathcal{A}$  satisfies the contraposition law, we have

$$(P) \quad \langle x, y \rangle \in \Theta_D \quad \text{iff} \quad \langle \neg x, \neg y \rangle \in \Theta_D.$$

Suppose now  $\langle x, y \rangle \in \Theta_D$ , i.e.  $x \rightarrow y, y \rightarrow x \in D$ . By (BCC1) we have  $(x \rightarrow y) \rightarrow ((z \rightarrow x) \rightarrow (z \rightarrow y)) = 1 \in D$  thus, due to Modus Ponens, also

$$(z \rightarrow x) \rightarrow (z \rightarrow y) \in D.$$

Analogously, we can show  $(z \rightarrow y) \rightarrow (z \rightarrow x) \in D$ , thus  $\langle z \rightarrow x, z \rightarrow y \rangle \in \Theta_D$ . Using (P), we obtain  $\langle \neg z \rightarrow \neg x, \neg z \rightarrow \neg y \rangle \in \Theta_D$ , thus also  $\langle x \rightarrow z, y \rightarrow z \rangle \in \Theta_D$ . Due to transitivity of  $\Theta_D$ , we have shown the Substitution Property of  $\Theta_D$  thus  $\Theta_D$  is a congruence on  $\mathcal{A}$ . ■

Of course, the set of all congruence kernels of a BCC-algebra  $\mathcal{A}$  forms a complete lattice with respect to set inclusion. Hence, for a given subset  $X \subseteq A$  there exists the least congruence kernel containing  $X$ , which will be denoted by  $F(X)$ . If  $X = \{a\}$  is a singleton,  $F(X)$  will be denoted briefly by  $F(a)$ .

In what follows we are going to characterize  $F(a)$  explicitly:

**LEMMA 6.** *Let  $\mathcal{A} = (A; \rightarrow, 1, 0)$  be a bounded BCC-algebra satisfying (CL) and  $a \in A$ . Define  $F_0^a = \{x \in A; a \leq x\}$  and  $F_i^a = \{x \in A; \alpha \rightarrow x = \beta \text{ for some } \alpha, \beta \in F_{i-1}^a\}$  for  $i = 1, 2, \dots$ . Then  $F(a) = \bigcup \{F_i^a; i = 0, 1, \dots\}$ .*

**Proof.** Certainly,  $1 \in F_0^a$ . If  $1 \in F_i^a$  then also  $1 \in F_{i+1}^a$  since  $1 \rightarrow 1 = 1$ . Thus  $1 \in F_i^a$  for all  $i = 0, 1, 2, \dots$ . In particular,  $1 \in F = \bigcup \{F_i^a; i = 0, 1, 2, \dots\}$ .

Furthermore,  $F_i^a \subseteq F_{i+1}^a$ . Indeed, if  $x \in F_i^a$  then  $1 \rightarrow x = x$ , thus  $x \in F_{i+1}^a$ . Now we assume  $x, x \rightarrow y \in F$ . By the previous observation, there is an integer  $j$  such that  $x, x \rightarrow y \in F_j^a$ . By definition, this means  $y \in F_{j+1}^a$  and hence  $y \in F$ . We have shown that  $F$  is a deductive system.

It is obvious that  $F(a) \subseteq F$  since  $F$  is a deductive system containing  $a$ . We show the converse inclusion by induction. Trivially,  $F_0^a \subseteq F(a)$ . Assume

$F_i^a \subseteq F(a)$  and let  $x \in F_{i+1}^a$ . Then there are  $\alpha, \beta \in F_i^a \subseteq F(a)$  such that  $\alpha \rightarrow x = \beta$ . This yields  $x \in F(a)$  and hence  $F(a) = F$ . ■

Moreover, if  $\mathcal{A}$  is a bounded commutative BCC-algebra, we can prove the following

LEMMA 7. *Let  $\mathcal{A} = (A; \rightarrow, 1, 0)$  be a bounded commutative BCC-algebra and  $a, b \in A$ ,  $a \parallel b$  and  $a \vee b = 1$ . Then for each  $x \in F(a)$  and  $y \in F(b)$  we have  $x \vee y = 1$ .*

Proof. (i) At first we prove that  $x \vee y = 1$  for each  $x \in F(a)$  and  $y \in F_0^b$ . This is clearly equivalent with  $x \vee b = 1$ . Since  $x \in F(a)$ , there is an index  $i$  such that  $x \in F_i^a$ . Evidently,  $x \vee b = 1$  for  $i = 0$ . Suppose now  $z \vee b = 1$  for all  $z \in F_i^a$  and let  $x \in F_{i+1}^a$ . Then there are  $\alpha, \beta \in F_i^a$  with  $\alpha \rightarrow x = \beta$ . Since  $\alpha \vee b = \beta \vee b = 1$  by the induction hypothesis, we obtain

$$\begin{aligned} x \rightarrow b &\leq (\alpha \rightarrow x) \rightarrow (\alpha \rightarrow b) = \beta \rightarrow ((\alpha \vee b) \rightarrow b) = \beta \rightarrow (1 \rightarrow b) = \\ &= \beta \rightarrow b = (\beta \vee b) \rightarrow b = 1 \rightarrow b = b. \end{aligned}$$

Since  $b \leq x \rightarrow b$ , we have  $x \rightarrow b = b$ , thus  $x \vee b = (x \rightarrow b) \rightarrow b = b \rightarrow b = 1$ .

(ii) Suppose now generally  $x \in F(a)$ ,  $y \in F(b)$ . Then  $y \in F_j^b$  for some index  $j$ . If  $j = 0$  then  $x \vee y = 1$  by (i). Assume  $x \vee z = 1$  for all  $z \in F_i^b$  and take  $y \in F_{i+1}^b$ . Then  $\alpha \rightarrow y = \beta$  for  $\alpha, \beta \in F_i^b$  thus, by the induction hypothesis,  $x \vee \alpha = x \vee \beta = 1$ . Hence

$$\begin{aligned} y \rightarrow x &\leq (\alpha \rightarrow y) \rightarrow (\alpha \rightarrow x) = \beta \rightarrow (\alpha \rightarrow x) = \beta \rightarrow ((\alpha \vee x) \rightarrow x) = \\ &= \beta \rightarrow (1 \rightarrow x) = \beta \rightarrow x = (\beta \vee x) \rightarrow x = 1 \rightarrow x = x. \end{aligned}$$

Since  $x \leq y \rightarrow x$ , we conclude  $y \rightarrow x = x$  thus

$$y \vee x = (y \rightarrow x) \rightarrow x = x \rightarrow x = 1. \quad \blacksquare$$

We are able to prove our main result.

THEOREM 4. *Let  $\mathcal{A} = (A; \rightarrow, 1, 0)$  be a bounded commutative BCC-algebra. If  $\mathcal{A}$  is subdirectly irreducible then  $\mathcal{A}$  is a chain with respect to the induced order.*

Proof. Suppose that  $\mathcal{A}$  is not a chain, i.e. there exist  $a', b' \in A$  such that  $a' \parallel b'$ . Since  $\mathcal{A}$  is commutative, it is a  $\vee$ -semilattice where  $x \vee y = (x \rightarrow y) \rightarrow y$ . The commutative law (C) implies

$$\neg \neg x = (x \rightarrow 0) \rightarrow 0 = (0 \rightarrow x) \rightarrow x = 1 \rightarrow x = x$$

thus  $\mathcal{A}$  satisfies also the double negation law and, due to Lemma 2 (ii), the mapping  $x \mapsto \neg x$  is an antitone involution. Hence,  $\mathcal{A}$  is in fact a lattice where  $x \wedge y = \neg(\neg x \vee \neg y)$ .



The same is in fact true for every section  $[p, 1]$ . The mapping assigning to  $x \in [p, 1]$  the element  $x \rightarrow p$  is an antitone involution on  $[p, 1]$  and hence

$$x \wedge y = ((x \rightarrow p) \vee (y \rightarrow p)) \rightarrow p$$

for any  $x, y \in [p, 1]$ .

$$\begin{aligned} \text{Therefore, } a \rightarrow (a \wedge b) &= a \rightarrow (((a \rightarrow 0) \vee (b \rightarrow 0)) \rightarrow 0) = \\ &= ((a \rightarrow 0) \vee (b \rightarrow 0)) \rightarrow (a \rightarrow 0) = \\ &= (((b \rightarrow 0) \rightarrow (a \rightarrow 0)) \rightarrow (a \rightarrow 0)) \rightarrow (a \rightarrow 0) = \\ &= (b \rightarrow 0) \rightarrow (a \rightarrow 0) = a \rightarrow b \text{ and hence} \end{aligned}$$

$$(a \rightarrow b) \vee (b \rightarrow a) = (a \rightarrow (a \wedge b)) \vee (b \rightarrow (a \wedge b)) = (a \wedge b) \rightarrow (a \wedge b) = 1.$$

Thus for  $a = a' \rightarrow b' \neq 1$  and  $b = b' \rightarrow a' \neq 1$  we obtain

$$\begin{aligned} a \vee b &= (a' \rightarrow b') \vee (b' \rightarrow a') = (a' \rightarrow (a' \wedge b')) \vee (b' \rightarrow (a' \wedge b')) = \\ &= (a' \wedge b') \rightarrow (a' \wedge b') = 1. \end{aligned}$$

By Lemma 7 we have  $x \vee y = 1$  for each  $x \in F(a)$  and  $y \in F(b)$ . This yields  $F(a) \cap F(b) = \{1\}$ . Since  $F(a) \neq \{1\} \neq F(b)$  and  $F(a)$  or  $F(b)$  uniquely determines the congruence  $\Theta(a, 1)$  or  $\Theta(b, 1)$  on  $\mathcal{A}$ , respectively, we have  $\Theta(a, 1) \neq \omega \neq \Theta(b, 1)$  but  $\Theta(a, 1) \cap \Theta(b, 1) = \omega$  thus  $\mathcal{A}$  is not subdirectly irreducible. ■

This result together with Theorem 2 and the Proposition yields that if an MV-algebra is subdirectly irreducible then it is a chain with respect to the induced order. This result is known, see e.g. [5], however, our new proof is much more simple.

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## A COMMON GENERALIZATION OF ORTHOLATTICES AND BOOLEAN QUASIRINGS

**Abstract.** In [2] a common generalization of Boolean algebras and Boolean rings was introduced. In a similar way we introduce a common generalization of ortholattices and Boolean quasirings.

In [2] a common generalization of Boolean algebras and Boolean rings was considered under the name N-algebra. In [1] the natural one-to-one correspondence between Boolean algebras and Boolean rings was generalized from Boolean algebras to ortholattices. The ring-like structures corresponding to ortholattices this way were called Boolean quasirings. Hence it is natural to ask for a common generalization of ortholattices and Boolean quasirings.

### 1. Ortholattices and Boolean quasirings

We start with the definition of an ortholattice:

**DEFINITION 1.1.** An *ortholattice* is an algebra  $(L, \vee, \wedge, ', 0, 1)$  of type  $(2, 2, 1, 0, 0)$  such that  $(L, \vee, \wedge, 0, 1)$  is a bounded lattice and

$$(x')' = x, (x \vee y)' = x' \wedge y', (x \wedge y)' = x' \vee y', x \vee x' = 1 \text{ and } x \wedge x' = 0$$

for all  $x, y \in L$ .

Next we define Boolean quasirings.

**DEFINITION 1.2 ([1]).** A *Boolean quasiring* is an algebra  $(R, +, \cdot, ', 0, 1)$  of type  $(2, 2, 1, 0, 0)$  satisfying

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$$\begin{aligned}
x + y &= y + x, \\
x + x &= 0, \\
x + 0 &= x, \\
(xy)z &= x(yz), \\
xy &= yx, \\
xx &= x, \\
x0 &= 0, \\
x1 &= x, \\
(xy + 1)(x + 1) + 1 &= x, \\
((x + 1)(y + 1) + 1)(xy + 1) &= x + y \text{ and} \\
x' &= x + 1.
\end{aligned}$$

REMARK 1.3. The definition given here is a slight modification of the original one given in [1] since the operation  $' : x \mapsto x + 1$  was added to the family of fundamental operations.

Now we can state the the correspondence between the two algebras introduced above:

THEOREM 1.4 ([1]). *Let  $\mathcal{L} = (L, \vee, \wedge, ', 0, 1)$  be an ortholattice. Define*

$$x + y := (x \vee y) \wedge (x \wedge y)' \text{ and } xy := x \wedge y$$

*for all  $x, y \in L$ . Then  $\mathbf{R}(\mathcal{L}) := (L, +, \cdot, ', 0, 1)$  is a Boolean quasiring. Conversely, let  $\mathcal{R} = (R, +, \cdot, ', 0, 1)$  be a Boolean quasiring. Define*

$$x \vee y := (x + 1)(y + 1) + 1 \text{ and } x \wedge y := xy$$

*for all  $x, y \in R$ . Then  $\mathbf{L}(\mathcal{R}) := (R, \vee, \wedge, ', 0, 1)$  is an ortholattice. Moreover,  $\mathbf{L}(\mathbf{R}(\mathcal{L})) = \mathcal{L}$  and  $\mathbf{R}(\mathbf{L}(\mathcal{R})) = \mathcal{R}$  for every ortholattice  $\mathcal{L}$  and every Boolean quasiring  $\mathcal{R}$ .*

## 2. QN-algebras

In this section we present a common generalization of ortholattices and Boolean quasirings. Since the common generalization of Boolean algebras and Boolean rings introduced in [2] was called an *N-algebra* we call our algebras *Quasi-N-algebras* or *QN-algebras*.

DEFINITION 2.1. A *QN-algebra* is an algebra  $(R, +, \cdot, ', 0, 1)$  of type  $(2, 2, 1, 0, 0)$  satisfying

$$\begin{aligned}
(xy)z &= x(yz), \\
xy &= yx, \\
x0 &= 0,
\end{aligned}$$

$$\begin{aligned}
x1 &= x, \\
(x')' &= x, \\
xx' &= 0, \\
(x'y')'x &= x, \\
((xy)'x')' &= x \text{ and} \\
x + y &= ((1 + 1)'xy)'(x'y')'.
\end{aligned}$$

We first prove that every QN-algebra induces an ortholattice.

LEMMA 2.2. *Let  $(R, +, \cdot, ', 0, 1)$  be a QN-algebra. Define  $x \vee y := (x'y')'$  for all  $x, y \in R$ . Then  $(R, \vee, \cdot, ', 0, 1)$  is an ortholattice.*

Proof. Because of

$$\begin{aligned}
(xy)z &= x(yz), \\
xy &= yx, \\
(x \vee y) \vee z &= (((x'y')')'z')' = ((x'y')z')' = (x'(y'z'))' = (x'((y'z')')')' \\
&= x \vee (y \vee z), \\
x \vee y &= (x'y')' = (y'x')' = y \vee x, \\
(x \vee y)x &= (x'y')'x = x, \\
(xy) \vee x &= ((xy)'x')' = x, \\
x0 &= 0, \\
x1 &= x, \\
(x')' &= x, \\
(x \vee y)' &= ((x'y')')' = x'y' \text{ and} \\
(xy)' &= ((x')'(y')')' = x' \vee y'
\end{aligned}$$

for all  $x, y, z \in R$ ,  $(R, \vee, \cdot, ', 0, 1)$  is a bounded lattice and  $'$  an antitone involution. Hence  $0' = 1$  and  $1' = 0$  which finally for all  $x \in R$  implies

$$x \vee x' = (x'(x')')' = (x'x)' = (xx')' = 0' = 1 \text{ and } xx' = 0. \blacksquare$$

THEOREM 2.3. *The ortholattices are exactly the QN-algebras  $(R, +, \cdot, ', 0, 1)$  satisfying  $1 + 1 = 1$ .*

Proof. Let  $(L, \vee, \wedge, ', 0, 1)$  be an ortholattice and define  $(L, +, \cdot, ', 0, 1) := (L, \vee, \wedge, ', 0, 1)$ . Then  $1 + 1 = 1 \vee 1 = 1$  and all axioms of a QN-algebra are satisfied since

$$\begin{aligned}
((1 + 1)'xy)'(x'y')' &= ((1 \vee 1)' \wedge x \wedge y)' \wedge (x' \wedge y')' = \\
&= (1' \wedge x \wedge y)' \wedge (x \vee y) = \\
&= (0 \wedge x \wedge y)' \wedge (x \vee y) =
\end{aligned}$$

$$\begin{aligned}
&= 0' \wedge (x \vee y) = 1 \wedge (x \vee y) = \\
&= x \vee y = x + y
\end{aligned}$$

for all  $x, y \in L$ .

Conversely, assume  $(R, +, \cdot, ', 0, 1)$  to be a QN-algebra satisfying  $1 + 1 = 1$ . Put  $x \vee y := (x'y')'$  for all  $x, y \in R$ . According to Lemma 2.2,  $(R, \vee, \cdot, ', 0, 1)$  is an ortholattice. Because of

$$\begin{aligned}
x + y &= ((1 + 1)'xy)'(x'y')' = (1'xy)'(x'y')' = (0xy)'(x'y')' = 0'(x'y')' \\
&= 1(x'y')' = (x'y')' = x \vee y
\end{aligned}$$

for all  $x, y \in R$ ,  $(R, +, \cdot, ', 0, 1)$  is an ortholattice. ■

**THEOREM 2.4.** *The Boolean quasirings are exactly the QN-algebras  $(R, +, \cdot, ', 0, 1)$  satisfying  $1 + 1 = 0$ .*

**Proof.** First let  $(R, +, \cdot, ', 0, 1)$  be a Boolean quasiring. Define  $x \vee y := (x'y')'$  for all  $x, y \in R$ . According to Theorem 1.4,  $(R, \vee, \cdot, ', 0, 1)$  is an ortholattice and  $1 + 1 = 1' = 0$ . Moreover,

$$((1 \vee 1)'x)'(x'y')' = (0'xy)'(x \vee y) = (1xy)'(x \vee y) = (xy)'(x \vee y) = x + y$$

for all  $x, y \in R$ .

Conversely, assume  $(R, +, \cdot, ', 0, 1)$  to be a QN-algebra satisfying  $1 + 1 = 0$ . Define  $x \vee y := (x'y')'$  for all  $x, y \in R$ . According to Lemma 2.2,  $(R, \vee, \cdot, ', 0, 1)$  is an ortholattice. Put  $x \oplus y := (x \vee y)(xy)'$  for all  $x, y \in R$ . According to Theorem 1.4,  $(R, \oplus, \cdot, ', 0, 1)$  is a Boolean quasiring. Now

$$\begin{aligned}
x \oplus y &= (x \vee y)(xy)' = (x'y')'(xy)' = (xy)'(x'y')' = (1xy)'(x'y')' \\
&= (0'xy)'(x'y')' = ((1 + 1)'xy)'(x'y')' = x + y
\end{aligned}$$

for all  $x, y \in R$  and hence  $(R, +, \cdot, ', 0, 1)$  is a Boolean quasiring. ■

### 3. Mutations of QN-algebras

**DEFINITION 3.1.** *Let  $\mathcal{R} = (R, +, \cdot, ', 0, 1)$  be a QN-algebra and  $a \in R$ . Then the algebra  $\mathcal{R}_a := (R, +_a, \cdot, ', 0, 1)$  with  $x +_a y := (a'xy)'(x'y')'$  for all  $x, y \in R$  is called the  $a$ -mutation of  $\mathcal{R}$ .*

We can now prove a theorem analogous to Theorem 3 of [2].

**THEOREM 3.2.** *Let  $\mathcal{R} = (R, +, \cdot, ', 0, 1)$  be a QN-algebra and  $a, b \in R$ . Then the following hold:*

- (i)  $1 +_a 1 = a$ .
- (ii)  $\mathcal{R}_a$  is a QN-algebra.
- (iii)  $\mathcal{R}_1$  is an ortholattice.
- (iv)  $\mathcal{R}_0$  is a Boolean quasiring.
- (v)  $\mathcal{R}_{1+1} = \mathcal{R}$ .
- (vi)  $(\mathcal{R}_a)_b = \mathcal{R}_b$ .
- (vii)  $\{\mathcal{R}_c \mid c \in R\}$  is the set of all QN-algebras with base set  $R$  having the same multiplication and the same unary operation as  $\mathcal{R}$ .
- (viii)  $\mathcal{R}$  and  $\mathcal{R}_a$  admit the same congruences.

Proof. (i)  $1 +_a 1 = (a'11)'(1'1')' = (a')'(1')' = a1 = a$ .

(ii)  $((1 +_a 1)'xy)'(x'y')' = (a'xy)'(x'y')' = x +_a y$  for all  $x, y \in R$ .

(iii) According to (ii),  $\mathcal{R}_1$  is a QN-algebra and according to (i),  $1 +_1 1 = 1$  and hence  $\mathcal{R}_1$  is an ortholattice according to Theorem 2.3.

(iv) According to (ii),  $\mathcal{R}_0$  is a QN-algebra and according to (i),  $1 +_0 1 = 0$  and hence  $\mathcal{R}_0$  is a Boolean quasiring according to Theorem 2.4.

(v)  $x +_{1+1} y = ((1 + 1)'xy)'(x'y')' = x + y$  for all  $x, y \in R$ .

(vi) Since  $\mathcal{R}$  is a QN-algebra, the same is true for  $\mathcal{R}_a = (R, +_a, \cdot', 0, 1)$  according to (ii), and  $x +_a y = (a'xy)'(x'y')'$  for all  $x, y \in R$ . Since  $\mathcal{R}_a$  is a QN-algebra, the same is true for  $(\mathcal{R}_a)_b = (R, (+_a)_b, \cdot', 0, 1)$  according to (ii), and  $x(+_a)_b y = (b'xy)'(x'y')' = x +_b y$  for all  $x, y \in R$ .

(vii) Let  $\mathcal{S} = (R, \oplus, \cdot', 0, 1)$  be a QN-algebra. Then  $x \oplus y = ((1 \oplus 1)'xy)'(x'y')' = x +_{1 \oplus 1} y$  for all  $x, y \in R$  and hence  $\mathcal{S} = \mathcal{R}_{1 \oplus 1}$ .

(viii)  $\mathcal{R}$  and  $\mathcal{R}_a$  admit the same congruences as  $(R, \cdot')$ . ■

Finally, we describe the correspondence stated in Theorem 1.4 in a simple way by means of mutations:

**THEOREM 3.3.** *The mappings  $\mathcal{R} \mapsto \mathcal{R}_0$  and  $\mathcal{R} \mapsto \mathcal{R}_1$  coincide with the mappings **R** and **L** introduced in Theorem 1.4, respectively.*

Proof. If  $\mathcal{L} = (L, \vee, \wedge, \cdot', 0, 1)$  is an ortholattice then it is a QN-algebra and  $x \vee_0 y = (0' \wedge x \wedge y)' \wedge (x' \wedge y')' = (1 \wedge x \wedge y)' \wedge (x \vee y) = (x \wedge y)' \wedge (x \vee y)$  for all  $x, y \in L$ . Hence  $\mathcal{L}_0 = \mathbf{R}(\mathcal{L})$ . If, conversely,  $\mathcal{R} = (R, +, \cdot', 0, 1)$  is a Boolean quasiring then it is a QN-algebra and

$$x +_1 y = (1'xy)'(x'y')' = (0xy)'(x'y')' = 0'(x'y')' = 1(x'y')' = (x'y')'$$

for all  $x, y \in L$  and hence  $\mathcal{R}_1 = \mathbf{L}(\mathcal{R})$ . ■

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