

Wojciech Żakowski

ON SOME CLASS OF NONLINEAR PROCESSES
WITH A MEMORY IN CONTINUOUS TIMEIntroduction

In [1] we have introduced the notion of a (α, k) -computation as a continuous function $x : <0; +\infty) \rightarrow \mathbb{R}$ satisfying the condition

$$(1) \quad x(t) = \int_0^k \alpha(s)x(t-k+s)ds \quad \text{for all } t \geq k,$$

where $\alpha : \mathbb{R} \rightarrow \mathbb{R}$ is a non-zero polynomial and k denotes a positive number. (α, k) -computations seem to be useful in describing some linear processes investigated in continuous time and which are characterized by a "memory" whose length is k time units. Such processes occur in many technical, economical and biological problems, e.g. in control theory, renewal theory, in the description of cells reproduction [2] etc. Basic properties of (α, k) -computations were investigated by Żakowski in [3], [4].

In this paper we introduce the notion of a CGF k -process as a continuous function $x : <0; +\infty) \rightarrow \mathbb{R}$ satisfying some integral, generally nonlinear condition (6). CGF k -processes seem to be useful in the description of some class of continuous processes, generally nonlinear, with "memory" whose length is k . We have also considered some qualitative properties of CGF k -processes.

1. Basic notations and definitions

Let \mathbb{R} denote the set of all real numbers and k denote an arbitrary positive number. By Δ and Δ_1 we denote the following sets

$$(2) \quad \Delta = \{(s, t, u) \in \mathbb{R}^3 : 0 \leq s \leq k \wedge t \geq k \wedge u \in \mathbb{R}\}$$

and

$$(3) \quad \Delta_1 = \{(t, v) \in \mathbb{R}^2 : t \geq k \wedge v \in \mathbb{R}\}.$$

Let $F : \Delta \rightarrow \mathbb{R}$ and $G : \Delta_1 \rightarrow \mathbb{R}$ be continuous functions. We assume that there exist positive numbers L_F and L_G such that for every $s \in \langle 0; k \rangle$, $t \geq k$, $\tilde{u}, \tilde{u} \in \mathbb{R}$ and $\tilde{v}, \tilde{v} \in \mathbb{R}$ the inequalities

$$(4) \quad |F(s, t, \tilde{u}) - F(s, t, \tilde{u})| \leq L_F \cdot |\tilde{u} - \tilde{u}|$$

and

$$(5) \quad |G(t, \tilde{v}) - G(t, \tilde{v})| \leq L_G \cdot |\tilde{v} - \tilde{v}|$$

hold.

Definition: The continuous function $x : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$ fulfilling for all $t \geq k$ the condition

$$(6) \quad x(t) = G\left(t, \int_0^k F(s, t, x(t-k+s)) ds\right)$$

is said to be CGFk-process (continuous GFk-process).

If $G(t, v) = v$ and $F(s, t, u) = \alpha(s) \cdot u$, where $\alpha : \mathbb{R} \rightarrow \mathbb{R}$ is a non-zero polynomial, then CGFk-process is a (α, k) -computation introduced in [1]. In this case the condition (6) is identical with the condition (1). Consequently, the notion of a CGFk-process is a generalization of the notion of (α, k) -computation. Any CGFk-process describes some real, continuous process which generally is nonlinear and has a "memory" of length k .

Remark 1. Let $f : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$ be an arbitrary continuous function. For every $k > 0$ this function is a CGFk-process when for example $F \equiv 0$ and $G(t, v) = f(t)$ for all $t \geq k$ and $v \in \mathbb{R}$.

If $x : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$, then the restriction of x to the set $U \subseteq \langle 0; +\infty \rangle$ is denoted by $x|U$. In particular, if x is a CGFk-process, then $x| \langle 0; k \rangle$ is called the initial state of x .

If $x : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$ and $T \geq 0$ then by x_T we denote the function $\langle 0; +\infty \rangle \rightarrow \mathbb{R}$ such that

$$(7) \quad x_T(t) = x(t + T) \quad \text{for all } t \geq 0.$$

The set of all real and continuous functions on the interval $\langle 0; a \rangle$ we denote by $\mathbb{C}_{\langle 0; a \rangle}$.

2. Some properties of the CGFk-processes

It follows from the condition (6) that if f is the initial state of any CGFk-process then

$$(8) \quad f(k) = G\left(k, \int_0^k F(s, k, f(s)) ds\right).$$

Theorem 1. If the function $f \in \mathbb{C}_{\langle 0; k \rangle}$ satisfies the condition (8), then there exists exactly one CGFk-process x such that $x| \langle 0; k \rangle = f$. This CGFk-process is a limit of the sequence (x_n) of successive approximations, defined as follows:

$$(9) \quad x_0(t) = \begin{cases} f(t) & \text{for } 0 \leq t \leq k \\ f(k) & \text{for } t > k \end{cases}$$

and

$$(10) \quad x_n(t) = \begin{cases} f(t) & \text{for } 0 \leq t \leq k \\ G\left(t, \int_0^k F(s, t, x_{n-1}(t-k+s)) ds\right) & \text{for } t > k \end{cases}$$

$n = 1, 2, \dots$. The sequence (x_n) is almost uniformly convergent on the interval $<0; +\infty)$.

P r o o f. Let δ denote an arbitrary positive number greater than k . We define the metric space:

$$(11) \quad C_f^{(k, \delta)} = \left\{ x \in C_{<0; \delta>} : x |_{<0; k>} = f \right\}$$

with the metric

$$(12) \quad \varrho(\tilde{x}, \tilde{\tilde{x}}) = \sup_{<0; \delta>} (e^{\lambda t} |\tilde{x}(t) - \tilde{\tilde{x}}(t)|),$$

where λ is a negative number such that

$$(13) \quad L_G \cdot L_F \cdot \frac{e^{\lambda k} - 1}{\lambda} < 1.$$

The space (11) is complete. On this space we define an operator A as follows

$$(14) \quad A[x(t)] = \begin{cases} f(t) & \text{for } 0 \leq t \leq k \\ G(t, \int_0^k F(s, t, x(t-k+s)) ds) & \text{for } k \leq t \leq \delta. \end{cases}$$

On the basis of (8) we observe that the operator A transforms the space (11) into itself.

In view of (14), (4), (5) and (12) we have for every $t \in (k; \delta)$ and for every $\tilde{x}, \tilde{\tilde{x}} \in C_f^{(k, \delta)}$:

$$\begin{aligned} & e^{\lambda t} |A[\tilde{x}(t)] - A[\tilde{\tilde{x}}(t)]| \leq \\ & \leq L_G \cdot L_F \int_0^k e^{\lambda(k-s)} e^{\lambda(t-k+s)} |\tilde{x}(t-k+s) - \tilde{\tilde{x}}(t-k+s)| ds \leq \\ & \leq L_G \cdot L_F \cdot \varrho(\tilde{x}, \tilde{\tilde{x}}) \int_0^k e^{\lambda(k-s)} ds = L_G \cdot L_F \cdot \varrho(\tilde{x}, \tilde{\tilde{x}}) \cdot \frac{e^{\lambda k} - 1}{\lambda}. \end{aligned}$$

This implies that for every $\tilde{x}, \tilde{\tilde{x}} \in C_f^{(k, \delta)}$

$$g(A[\tilde{x}], A[\tilde{\tilde{x}}]) \leq L_G \cdot L_F \frac{e^{\lambda k} - 1}{\lambda} g(\tilde{x}, \tilde{\tilde{x}}).$$

From this and from (13), applying Banach's fixed point theorem, it follows that there exists exactly one function $x_* \in C_f^{(k, \delta)}$ such that $x_* = A[x_*]$. Moreover, $x_* = \lim_{n \rightarrow \infty} x_n$, where x_n , $n = 0, 1, 2, \dots$ is defined by (9) and (10) for $t \in \langle 0; \delta \rangle$. From (12) we note that the sequence (x_n) is uniformly convergent on the interval $\langle 0; \delta \rangle$. Because δ denotes an arbitrary positive number greater than k , the proof is complete.

Remark 2. The application of the metric (9) causes that there is no constraints imposed on the numbers L_G and L_F in Lipschitz's conditions (4) and (5).

Theorem 2. If x is an CGFk-process, and, moreover, if:

1° there exists a continuous function

$$M_F : \{(s, t) \in \mathbb{R}^2 : 0 \leq s \leq k \wedge t > k\} \rightarrow \langle 0; +\infty \rangle$$

such that for every $(s, t, u) \in \Delta$ (see (2)) the condition

$$(15) \quad |F(s, t, u)| \leq M_F(s, t) \cdot |u|$$

holds,

2° there exists a function

$$M_G : \langle k; +\infty \rangle \rightarrow \langle 0; +\infty \rangle$$

such that for every $(t, v) \in \Delta_1$ the condition

$$(16) \quad |G(t, v)| \leq M_G(t) \cdot |v|$$

holds,

3° there exists a number $t_0 \geq 0$ for which the function $x|_{[t_0; t_0+k]}$ is nonnegative (nonpositive) and nonzero and, moreover

$$(17) \quad M_G(t_0+k) \cdot \int_0^k M_F(s, t_0+k) ds < 1,$$

then the function $x|_{[t_0; t_0+k]}$ is not nondecreasing (or not nonincreasing, respectively).

P r o o f. If the function $x|_{[t_0; t_0+k]}$ is nonnegative and nonzero, then in view of (6), (15) and (16) we have

$$x(t_0+k) \leq M_G(t_0+k) \cdot \int_0^k M_F(s, t_0+k) x(t_0+s) ds.$$

Hence we get

$$(18) \quad x(t_0+k) \leq \sup_{[t_0; t_0+k]} x \cdot M_G(t_0+k) \cdot \int_0^k M_F(s, t_0+k) ds.$$

If

$$M_G(t_0+k) = 0 \quad \text{or} \quad \int_0^k M_F(s, t_0+k) ds = 0$$

then we have $x(t_0+k) = 0$. This implies that the function $x|_{[t_0; t_0+k]}$ is not nondecreasing. In the case if

$$M_G(t_0+k) \int_0^k M_F(s, t_0+k) ds > 0$$

the inequality

$$x(t_0+k) < \sup_{[t_0; t_0+k]} x$$

holds on the basis of (17) and (18). On the other hand, there exists a number $c \in [t_0; t_0+k]$ such that $x(c) = \sup_{[t_0; t_0+k]} x$,

whence it follows that $x(t_0+k) < x(c)$, which completes the proof for the case of nonnegative and nonzero function $x|_{[t_0; t_0+k]}$. In the case when this function is nonpositive and nonzero, the proof is analogous. Q.E.D.

Theorem 3. If x is an CGFk-process and the hypotheses 1⁰ and 2⁰ of Theorem 2 hold and, moreover, for every $t > k$ the condition

$$(19) \quad M_G(t) \cdot \int_0^k M_F(s, t) ds \leq 1$$

holds, then there exists a number $c \in \langle 0; k \rangle$ such that for every $t > 0$

$$(20) \quad |x(t)| \leq |x(c)|.$$

Proof. The CGFk-process x is a limit of the sequence (x_n) of successive approximations, defined by equalities (9) and (10), where $f = x|_{\langle 0; k \rangle}$. Let $M = \sup_{\langle 0; k \rangle} |x(t)|$. On the basis of (15) and (16) we have for every $t > k$ and $n = 1, 2, \dots$

$$(21) \quad |x_n(t)| \leq M_G(t) \cdot \int_0^k M_F(s, t) |x_{n-1}(t-k+s)| ds.$$

We observe (see (9)) that $|x_0(t)| \leq M$ for every $t \geq 0$. From this and from inequalities (19) and (21) applying mathematical induction we get $|x_n(t)| \leq M$ for every $t \geq 0$ and $n = 0, 1, 2, \dots$. Passing in this inequality to the limit with $n \rightarrow +\infty$ we get $|x(t)| \leq M$ for every $t \geq 0$. Obviously, there exists a number $c \in \langle 0; k \rangle$ such that $|x(c)| = M$, so we get inequality (20). Q.E.D.

Corollary 1. If hypotheses 1⁰ and 2⁰ of Theorem 2 and the condition (19) hold, then any CGFk-process is bounded on the interval $\langle 0; +\infty \rangle$.

Theorems 2 and 3 generalize some analogous theorems of the paper [3].

Theorem 4. If x is an CGFk-process and if there exists a finite limit

$$(22) \quad \lim_{t \rightarrow +\infty} x(t) = g$$

and, moreover, if $F(s, t, g)$ converges for $t \rightarrow +\infty$ to the finite limit $\tilde{F}(s)$, uniformly with $s \in \langle 0; k \rangle$, i.e. for every $\varepsilon > 0$ there is a $\delta > k$ such that for every $s \in \langle 0; k \rangle$

$$(23) \quad t > \delta \Rightarrow |F(s, t, g) - \tilde{F}(s)| < \varepsilon,$$

then

$$(24) \quad g = \lim_{t \rightarrow +\infty} G\left(t, \int_0^k \tilde{F}(s) ds\right).$$

Proof. From the condition (22) we have for every $t > k$

$$g + \varepsilon(t) = G\left(t, \int_0^k F(s, t, g + \varepsilon(t-k+s)) ds\right),$$

where $\varepsilon(t) \rightarrow 0$ if $t \rightarrow +\infty$. Hence, in view of (4) and (5)

$$g + \varepsilon(t) = G\left(t, \int_0^k F(s, t, g) ds\right) + \mu(t),$$

where $\mu(t) \rightarrow 0$ if $t \rightarrow +\infty$. The assumptions imply that the function \tilde{F} is continuous on the interval $\langle 0; k \rangle$. Consequently, by (5) and (23) we have

$$(25) \quad g + \varepsilon(t) = G\left(t, \int_0^k \tilde{F}(s) ds\right) + \tilde{\mu}(t),$$

where $\tilde{\mu}(t) \rightarrow 0$ if $t \rightarrow +\infty$. Passing to the limit with $t \rightarrow +\infty$ in equality (25) we get condition (24). Q.E.D.

R e m a r k 3. In the case if $G(t, v) = v$ and $F(s, t, u) = \alpha(s) \cdot u$, where $\alpha : \mathbb{R} \rightarrow \mathbb{R}$ is a non-zero polynomial, i.e. in the case if an CGFk-process is a (α, k) -computation, the condition (24) has the form

$$g = \int_0^k \alpha(s) g ds.$$

If $g \neq 0$, then this equality is equivalent to the equality

$$\int_0^k \alpha(s) ds = 1$$

given in the papers [1] and [3].

3. Relations between CGFk-processes and k-computable functions

T h e o r e m 5. If x is CGFk-process and if there exists a number $T > 0$ such that for every $t > k$, $s \in \langle 0; k \rangle$, $v \in \mathbb{R}$ and $u \in \mathbb{R}$

$$(26) \quad G(t+T, v) = G(t, v) \quad \text{and} \quad F(s, t+T, u) = F(s, t, u)$$

then the function x_T (see (7)) is also a CGFk-process.

P r o o f . The function x_T is obviously continuous. In view of (6) we have for every $t > k$

$$x(t+T) = G\left(t+T, \int_0^k F(s, t+T, x(t+T-k+s)) ds\right).$$

Hence, in virtue of (26) and (7)

$$x_T(t) = G\left(t, \int_0^k F(s, t, x_T(t-k+s)) ds\right)$$

for all $t > k$, which completes the proof.

Theorem 6. If x_1 and x_2 are CGFk-processes and the condition (26) holds for $T = T_1$ and $T = T_2$, then

$$(27) \quad \left[(x_1 | \langle T_1; T_1+k \rangle)_{T_1} = (x_2 | \langle T_2; T_2+k \rangle)_{T_2} \right] \Rightarrow$$

$$\Rightarrow \left[(x_1 | \langle T_1; +\infty \rangle)_{T_1} = (x_2 | \langle T_2; +\infty \rangle)_{T_2} \right].$$

Proof. According to (7) we have

$$h_1(t) = (x_1 | \langle T_1; T_1+k \rangle)_{T_1}(t) = x_1(t+T_1) \quad \text{for } t \in \langle 0; k \rangle$$

$$h_2(t) = (x_2 | \langle T_2; T_2+k \rangle)_{T_2}(t) = x_2(t+T_2) \quad \text{for } t \in \langle 0; k \rangle$$

$$H_1(t) = (x_1 | \langle T_1; +\infty \rangle)_{T_1}(t) = x_1(t+T_1) \quad \text{for } t \in \langle 0; +\infty \rangle$$

and

$$H_2(t) = (x_2 | \langle T_2; +\infty \rangle)_{T_2}(t) = x_2(t+T_2) \quad \text{for } t \in \langle 0; +\infty \rangle.$$

It follows from Theorem 5 that the function H_1 is an CGFK-process with the initial state h_1 and the function H_2 is an CGFk-process with the initial state h_2 . Consequently, in virtue of Theorem 1, the equality $h_1 = h_2$ implies $H_1 = H_2$. Q.E.D.

If the conditions (26) hold for all $T \geq 0$, i.e. if the functions F and G are constant with respect to the variable t then from Theorem 6 we get for any CGFk-process x the following conditions

$$(28) \quad \left[(x | \langle T_1; T_1+k \rangle)_{T_1} = (x | \langle T_2; T_2+k \rangle)_{T_2} \right] \Rightarrow$$

$$\Rightarrow \left[(x | \langle T_1; +\infty \rangle)_{T_1} = (x | \langle T_2; +\infty \rangle)_{T_2} \right]$$

for every $T_1 > 0$ and $T_2 > 0$. This condition is a necessary and sufficient condition of k -computability in the sense of the paper [5]. Consequently, we have

Corollary 2. If the conditions (26) hold for all $T > 0$, then any CGF k -process x is k -computable.

We observe that in the linear case, if $G(t, v) = v$ and $F(s, t, u) = \alpha(s) \cdot u$ where $\alpha: \mathbb{R} \rightarrow \mathbb{R}$ is a non-zero polynomial, the conditions (26) obviously hold for all $T > 0$. Hence any (α, k) -computation is k -computable.

The k -computable functions possess many interesting properties (see [5]). In the special case when the function F and G are constant with respect to the variable t it follows from Corollary 2 that CGF k -processes possesses also these properties.

Note that investigations, given in papers [6] and [7], concerning (Z, Q) -computability also concern k -computability because any k -computable function is (Z, Q) -computable in the case if $Z = \langle 0; k \rangle$ and $Q = \langle 0; +\infty \rangle$.

REFERENCES

- [1] W. Żakowski : On a concept of (α, k) -computation, Bull. Acad. Polon. Sci., Ser. Sci. Techn. 27 (1979) 841-844.
- [2] M. Kimmel : Integro-differential equations in cellular population dynamics, Proc. Internat. Conf. "Functional Differential Systems and Related Topics", Błażejewko, 19-26 May, 1979.
- [3] W. Żakowski : (α, k) -computations in continuous time, Demonstratio Math., 12 (1979) 1071-1087.
- [4] W. Żakowski : Defining polynomials for the functions $\sin t$ and $\cos t$, [in Polish], ICS PAS Reports, 371 (1979).
- [5] B. Konikowska : Continuous machines, Information and Control 22 (1973) 353-372.

- [6] W. Łukowski : A generalization of the notions of a machine and computability, *Information and Control* 33 (1977) 166-176.
- [7] E. Stankiewicz : On some properties of (Z, Q) -machine and its computations, *Bull. Acad. Polon. Sci., Ser. Sci. Math. Astron. Phys.* 25 (1977) 205-208.

INSTITUTE OF MATHEMATICS, TECHNICAL UNIVERSITY OF WARSAW

Received August 11th, 1980.