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ON SOME CLASS OF NONLINEAR PROCESSES  
WITH A MEMORY IN CONTINUOUS TIMEIntroduction

In [1] we have introduced the notion of a  $(\alpha, k)$ -computation as a continuous function  $x : <0; +\infty) \rightarrow \mathbb{R}$  satisfying the condition

$$(1) \quad x(t) = \int_0^k \alpha(s)x(t-k+s)ds \quad \text{for all } t \geq k,$$

where  $\alpha : \mathbb{R} \rightarrow \mathbb{R}$  is a non-zero polynomial and  $k$  denotes a positive number.  $(\alpha, k)$ -computations seem to be useful in describing some linear processes investigated in continuous time and which are characterized by a "memory" whose length is  $k$  time units. Such processes occur in many technical, economical and biological problems, e.g. in control theory, renewal theory, in the description of cells reproduction [2] etc. Basic properties of  $(\alpha, k)$ -computations were investigated by Żakowski in [3], [4].

In this paper we introduce the notion of a CGFk-process as a continuous function  $x : <0; +\infty) \rightarrow \mathbb{R}$  satisfying some integral, generally nonlinear condition (6). CGFk-processes seem to be useful in the description of some class of continuous processes, generally nonlinear, with "memory" whose length is  $k$ . We have also considered some qualitative properties of CGFk-processes.

### 1. Basic notations and definitions

Let  $\mathbb{R}$  denote the set of all real numbers and  $k$  denote an arbitrary positive number. By  $\Delta$  and  $\Delta_1$  we denote the following sets

$$(2) \quad \Delta = \{(s, t, u) \in \mathbb{R}^3 : 0 \leq s \leq k \wedge t \geq k \wedge u \in \mathbb{R}\}$$

and

$$(3) \quad \Delta_1 = \{(t, v) \in \mathbb{R}^2 : t \geq k \wedge v \in \mathbb{R}\}.$$

Let  $F : \Delta \rightarrow \mathbb{R}$  and  $G : \Delta_1 \rightarrow \mathbb{R}$  be continuous functions. We assume that there exist positive numbers  $L_F$  and  $L_G$  such that for every  $s \in \langle 0; k \rangle$ ,  $t \geq k$ ,  $\tilde{u}, \tilde{\tilde{u}} \in \mathbb{R}$  and  $\tilde{v}, \tilde{\tilde{v}} \in \mathbb{R}$  the inequalities

$$(4) \quad |F(s, t, \tilde{u}) - F(s, t, \tilde{\tilde{u}})| \leq L_F \cdot |\tilde{u} - \tilde{\tilde{u}}|$$

and

$$(5) \quad |G(t, \tilde{v}) - G(t, \tilde{\tilde{v}})| \leq L_G \cdot |\tilde{v} - \tilde{\tilde{v}}|$$

hold.

**Definition:** The continuous function  $x : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$  fulfilling for all  $t \geq k$  the condition

$$(6) \quad x(t) = G\left(t, \int_0^k F(s, t, x(t-k+s)) ds\right)$$

is said to be CGFk-process (continuous GFk-process).

If  $G(t, v) = v$  and  $F(s, t, u) = \alpha(s) \cdot u$ , where  $\alpha : \mathbb{R} \rightarrow \mathbb{R}$  is a non-zero polynomial, then CGFk-process is a  $(\alpha, k)$ -computation introduced in [1]. In this case the condition (6) is identical with the condition (1). Consequently, the notion of a CGFk-process is a generalization of the notion of  $(\alpha, k)$ -computation. Any CGFk-process describes some real, continuous process which generally is nonlinear and has a "memory" of length  $k$ .

**R e m a r k 1.** Let  $f : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$  be an arbitrary continuous function. For every  $k > 0$  this function is a CGFk-process when for example  $F \equiv 0$  and  $G(t, v) = f(t)$  for all  $t \geq k$  and  $v \in \mathbb{R}$ .

If  $x : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$ , then the restriction of  $x$  to the set  $U \subseteq \langle 0; +\infty \rangle$  is denoted by  $x|_U$ . In particular, if  $x$  is a CGFk-process, then  $x|_{\langle 0; k \rangle}$  is called the initial state of  $x$ .

If  $x : \langle 0; +\infty \rangle \rightarrow \mathbb{R}$  and  $T \geq 0$  then by  $x_T$  we denote the function  $\langle 0; +\infty \rangle \rightarrow \mathbb{R}$  such that

$$(7) \quad x_T(t) = x(t + T) \quad \text{for all } t \geq 0.$$

The set of all real and continuous functions on the interval  $\langle 0; a \rangle$  we denote by  $C_{\langle 0; a \rangle}$ .

## 2. Some properties of the CGFk-processes

It follows from the condition (6) that if  $f$  is the initial state of any CGFk-process then

$$(8) \quad f(k) = G\left(k, \int_0^k F(s, k, f(s)) ds\right).$$

**T h e o r e m 1.** If the function  $f \in C_{\langle 0; k \rangle}$  satisfies the condition (8), then there exists exactly one CGFk-process  $x$  such that  $x|_{\langle 0; k \rangle} = f$ . This CGFk-process is a limit of the sequence  $(x_n)$  of successive approximations, defined as follows:

$$(9) \quad x_0(t) = \begin{cases} f(t) & \text{for } 0 \leq t \leq k \\ f(k) & \text{for } t > k \end{cases}$$

and

$$(10) \quad x_n(t) = \begin{cases} f(t) & \text{for } 0 \leq t \leq k \\ G\left(t, \int_0^k F(s, t, x_{n-1}(t-k+s)) ds\right) & \text{for } t > k \end{cases}$$

$n = 1, 2, \dots$ . The sequence  $(x_n)$  is almost uniformly convergent on the interval  $\langle 0; +\infty \rangle$ .

**P r o o f .** Let  $\delta$  denote an arbitrary positive number greater than  $k$ . We define the metric space:

$$(11) \quad \mathbb{C}_f^{(k, \delta)} = \left\{ x \in \mathbb{C} \langle 0; \delta \rangle : x|_{\langle 0; k \rangle} = f \right\}$$

with the metric

$$(12) \quad \rho(\tilde{x}, \tilde{\tilde{x}}) = \sup_{\langle 0; \delta \rangle} (e^{\lambda t} |\tilde{x}(t) - \tilde{\tilde{x}}(t)|),$$

where  $\lambda$  is a negative number such that

$$(13) \quad L_G \cdot L_F \cdot \frac{e^{\lambda k} - 1}{\lambda} < 1.$$

The space (11) is complete. On this space we define an operator  $A$  as follows

$$(14) \quad A[x(t)] = \begin{cases} f(t) & \text{for } 0 \leq t \leq k \\ G(t, \int_0^k F(s, t, x(t-k+s)) ds) & \text{for } k \leq t \leq \delta. \end{cases}$$

On the basis of (8) we observe that the operator  $A$  transforms the space (11) into itself.

In view of (14), (4), (5) and (12) we have for every  $t \in (k; \delta)$  and for every  $\tilde{x}, \tilde{\tilde{x}} \in \mathbb{C}_f^{(k, \delta)}$ :

$$\begin{aligned} & e^{\lambda t} |A[\tilde{x}(t)] - A[\tilde{\tilde{x}}(t)]| \leq \\ & \leq L_G \cdot L_F \int_0^k e^{\lambda(k-s)} e^{\lambda(t-k+s)} |\tilde{x}(t-k+s) - \tilde{\tilde{x}}(t-k+s)| ds \leq \\ & \leq L_G \cdot L_F \cdot \rho(\tilde{x}, \tilde{\tilde{x}}) \int_0^k e^{\lambda(k-s)} ds = L_G \cdot L_F \cdot \rho(\tilde{x}, \tilde{\tilde{x}}) \cdot \frac{e^{\lambda k} - 1}{\lambda}. \end{aligned}$$

This implies that for every  $\tilde{x}, \tilde{\tilde{x}} \in C_f^{(k, \delta)}$

$$\varphi(A[\tilde{x}], A[\tilde{\tilde{x}}]) \leq L_G \cdot L_F \frac{e^{\lambda k} - 1}{\lambda} \varphi(\tilde{x}, \tilde{\tilde{x}}).$$

From this and from (13), applying Banach's fixed point theorem, it follows that there exists exactly one function  $x_* \in C_f^{(k, \delta)}$  such that  $x_* = A[x_*]$ . Moreover,  $x_* = \lim_{n \rightarrow \infty} x_n$ , where  $x_n$ ,  $n = 0, 1, 2, \dots$  is defined by (9) and (10) for  $t \in \langle 0; \delta \rangle$ . From (12) we note that the sequence  $(x_n)$  is uniformly convergent on the interval  $\langle 0; \delta \rangle$ . Because  $\delta$  denotes an arbitrary positive number greater than  $k$ , the proof is complete.

**Remark 2.** The application of the metric (9) causes that there is no constraints imposed on the numbers  $L_G$  and  $L_F$  in Lipschitz's conditions (4) and (5).

**Theorem 2.** If  $x$  is an CGPk-process, and, moreover, if:

1° there exists a continuous function

$$M_F : \{(s, t) \in \mathbb{R}^2 : 0 \leq s \leq k \wedge t > k\} \rightarrow \langle 0; +\infty \rangle$$

such that for every  $(s, t, u) \in \Delta$  (see (2)) the condition

$$(15) \quad |F(s, t, u)| \leq M_F(s, t) \cdot |u|$$

holds,

2° there exists a function

$$M_G : \langle k; +\infty \rangle \rightarrow \langle 0; +\infty \rangle$$

such that for every  $(t, v) \in \Delta_1$  the condition

$$(16) \quad |G(t, v)| \leq M_G(t) \cdot |v|$$

holds,

$3^0$  there exists a number  $t_0 \geq 0$  for which the function  $x|_{\langle t_0; t_0+k \rangle}$  is nonnegative (nonpositive) and nonzero and, moreover

$$(17) \quad M_G(t_0+k) \cdot \int_0^k M_F(s, t_0+k) ds < 1,$$

then the function  $x|_{\langle t_0; t_0+k \rangle}$  is not nondecreasing (or not nonincreasing, respectively).

*P r o o f .* If the function  $x|_{\langle t_0; t_0+k \rangle}$  is nonnegative and nonzero, then in view of (6), (15) and (16) we have

$$x(t_0+k) \leq M_G(t_0+k) \cdot \int_0^k M_F(s, t_0+k) x(t_0+s) ds.$$

Hence we get

$$(18) \quad x(t_0+k) \leq \sup_{\langle t_0; t_0+k \rangle} x \cdot M_G(t_0+k) \cdot \int_0^k M_F(s, t_0+k) ds.$$

If

$$M_G(t_0+k) = 0 \quad \text{or} \quad \int_0^k M_F(s, t_0+k) ds = 0$$

then we have  $x(t_0+k) = 0$ . This implies that the function  $x|_{\langle t_0; t_0+k \rangle}$  is not nondecreasing. In the case if

$$M_G(t_0+k) \int_0^k M_F(s, t_0+k) ds > 0$$

the inequality

$$x(t_0+k) < \sup_{\langle t_0; t_0+k \rangle} x$$

holds on the basis of (17) and (18). On the other hand, there exists a number  $c \in \langle t_0; t_0+k \rangle$  such that  $x(c) = \sup_{\langle t_0; t_0+k \rangle} x$ ,

whence it follows that  $x(t_0+k) < x(c)$ , which completes the proof for the case of nonnegative and nonzero function  $x|<t_0; t_0+k>$ . In the case when this function is nonpositive and nonzero, the proof is analogous. Q.E.D.

**Theorem 3.** If  $x$  is an CGPk-process and the hypotheses  $1^0$  and  $2^0$  of Theorem 2 hold and, moreover, for every  $t > k$  the condition

$$(19) \quad M_G(t) \cdot \int_0^k M_F(s, t) ds \leq 1$$

holds, then there exists a number  $c \in <0; k>$  such that for every  $t > 0$

$$(20) \quad |x(t)| \leq |x(c)|.$$

**Proof.** The CGPk-process  $x$  is a limit of the sequence  $(x_n)$  of successive approximations, defined by equalities (9) and (10), where  $f = x|<0; k>$ . Let  $M = \sup_{<0; k>} |x(t)|$ . On the basis of (15) and (16) we have for every  $t > k$  and  $n = 1, 2, \dots$

$$(21) \quad |x_n(t)| \leq M_G(t) \cdot \int_0^k M_F(s, t) |x_{n-1}(t-k+s)| ds.$$

We observe (see (9)) that  $|x_0(t)| \leq M$  for every  $t \geq 0$ . From this and from inequalities (19) and (21) applying mathematical induction we get  $|x_n(t)| \leq M$  for every  $t \geq 0$  and  $n = 0, 1, 2, \dots$ . Passing in this inequality to the limit with  $n \rightarrow +\infty$  we get  $|x(t)| \leq M$  for every  $t \geq 0$ . Obviously, there exists a number  $c \in <0; k>$  such that  $|x(c)| = M$ , so we get inequality (20). Q.E.D.

**Corollary 1.** If hypotheses  $1^0$  and  $2^0$  of Theorem 2 and the condition (19) hold, then any CGPk-process is bounded on the interval  $<0; +\infty>$ .

Theorems 2 and 3 generalize some analogous theorems of the paper [3].

**Theorem 4.** If  $x$  is an CGPk-process and if there exists a finite limit

$$(22) \quad \lim_{t \rightarrow +\infty} x(t) = g$$

and, moreover, if  $F(s, t, g)$  converges for  $t \rightarrow +\infty$  to the finite limit  $\tilde{F}(s)$ , uniformly with  $s \in \langle 0; k \rangle$ , i.e. for every  $\varepsilon > 0$  there is a  $\delta > k$  such that for every  $s \in \langle 0; k \rangle$

$$(23) \quad t > \delta \Rightarrow |F(s, t, g) - \tilde{F}(s)| < \varepsilon,$$

then

$$(24) \quad g = \lim_{t \rightarrow +\infty} G\left(t, \int_0^k \tilde{F}(s) ds\right).$$

**Proof.** From the condition (22) we have for every  $t > k$

$$g + \varepsilon(t) = G\left(t, \int_0^k F(s, t, g + \varepsilon(t - k + s)) ds\right),$$

where  $\varepsilon(t) \rightarrow 0$  if  $t \rightarrow +\infty$ . Hence, in view of (4) and (5)

$$g + \varepsilon(t) = G\left(t, \int_0^k F(s, t, g) ds\right) + \mu(t),$$

where  $\mu(t) \rightarrow 0$  if  $t \rightarrow +\infty$ . The assumptions imply that the function  $\tilde{F}$  is continuous on the interval  $\langle 0; k \rangle$ . Consequently, by (5) and (23) we have

$$(25) \quad g + \varepsilon(t) = G\left(t, \int_0^k \tilde{F}(s) ds\right) + \tilde{\mu}(t),$$



where  $\tilde{\mu}(t) \rightarrow 0$  if  $t \rightarrow +\infty$ . Passing to the limit with  $t \rightarrow +\infty$  in equality (25) we get condition (24). Q.E.D.

**R e m a r k 3.** In the case if  $G(t, v) = v$  and  $F(s, t, u) = \alpha(s) \cdot u$ , where  $\alpha: \mathbb{R} \rightarrow \mathbb{R}$  is a non-zero polynomial, i.e. in the case if an CGFk-process is a  $(\alpha, k)$ -computation, the condition (24) has the form

$$g = \int_0^k \alpha(s) g ds.$$

If  $g \neq 0$ , then this equality is equivalent to the equality

$$\int_0^k \alpha(s) ds = 1$$

given in the papers [1] and [3].

### 3. Relations between CGFk-processes and k-computable functions

**T h e o r e m 5.** If  $x$  is CGFk-process and if there exists a number  $T \geq 0$  such that for every  $t \geq k$ ,  $s \in \langle 0; k \rangle$ ,  $v \in \mathbb{R}$  and  $u \in \mathbb{R}$

$$(26) \quad G(t+T, v) = G(t, v) \quad \text{and} \quad F(s, t+T, u) = F(s, t, u)$$

then the function  $x_T$  (see (7)) is also a CGFk-process.

**P r o o f .** The function  $x_T$  is obviously continuous. In view of (6) we have for every  $t \geq k$

$$x(t+T) = G\left(t+T, \int_0^k F(s, t+T, x(t+T-k+s)) ds\right).$$

Hence, in virtue of (26) and (7)

$$x_T(t) = G\left(t, \int_0^k F(s, t, x_T(t-k+s)) ds\right)$$

for all  $t \geq k$ , which completes the proof.

**Theorem 6.** If  $x_1$  and  $x_2$  are CGFk-processes and the condition (26) holds for  $T = T_1$  and  $T = T_2$ , then

$$(27) \quad \left[ (x_1 | \langle T_1; T_1+k \rangle)_{T_1} = (x_2 | \langle T_2; T_2+k \rangle)_{T_2} \right] \Rightarrow \\ \Rightarrow \left[ (x_1 | \langle T_1; +\infty \rangle)_{T_1} = (x_2 | \langle T_2; +\infty \rangle)_{T_2} \right].$$

**Proof.** According to (7) we have

$$h_1(t) = (x_1 | \langle T_1; T_1+k \rangle)_{T_1}(t) = x_1(t+T_1) \quad \text{for } t \in \langle 0; k \rangle$$

$$h_2(t) = (x_2 | \langle T_2; T_2+k \rangle)_{T_2}(t) = x_2(t+T_2) \quad \text{for } t \in \langle 0; k \rangle$$

$$H_1(t) = (x_1 | \langle T_1; +\infty \rangle)_{T_1}(t) = x_1(t+T_1) \quad \text{for } t \in \langle 0; +\infty \rangle$$

and

$$H_2(t) = (x_2 | \langle T_2; +\infty \rangle)_{T_2}(t) = x_2(t+T_2) \quad \text{for } t \in \langle 0; +\infty \rangle.$$

It follows from Theorem 5 that the function  $H_1$  is an CGFk-process with the initial state  $h_1$  and the function  $H_2$  is an CGFk-process with the initial state  $h_2$ . Consequently, in virtue of Theorem 1, the equality  $h_1 = h_2$  implies  $H_1 = H_2$ . Q.E.D.

If the conditions (26) hold for all  $T > 0$ , i.e. if the functions  $F$  and  $G$  are constant with respect to the variable  $t$  then from Theorem 6 we get for any CGFk-process  $x$  the following conditions

$$(28) \quad \left[ (x | \langle T_1; T_1+k \rangle)_{T_1} = (x | \langle T_2; T_2+k \rangle)_{T_2} \right] \Rightarrow \\ \Rightarrow \left[ (x | \langle T_1; +\infty \rangle)_{T_1} = (x | \langle T_2; +\infty \rangle)_{T_2} \right]$$

for every  $T_1 \geq 0$  and  $T_2 \geq 0$ . This condition is a necessary and sufficient condition of  $k$ -computability in the sense of the paper [5]. Consequently, we have

**C o r o l l a r y 2.** If the conditions (26) hold for all  $T \geq 0$ , then any CGFk-process  $x$  is  $k$ -computable.

We observe that in the linear case, if  $G(t, v) = v$  and  $F(s, t, u) = \alpha(s) \cdot u$  where  $\alpha: \mathbb{R} \rightarrow \mathbb{R}$  is a non-zero polynomial, the conditions (26) obviously hold for all  $T \geq 0$ . Hence any  $(\alpha, k)$ -computation is  $k$ -computable.

The  $k$ -computable functions possess many interesting properties (see [5]). In the special case when the function  $F$  and  $G$  are constant with respect to the variable  $t$  it follows from Corollary 2 that CGFk-processes possess also these properties.

Note that investigations, given in papers [6] and [7], concerning  $(Z, Q)$ -computability also concern  $k$ -computability because any  $k$ -computable function is  $(Z, Q)$ -computable in the case if  $Z = \langle 0; k \rangle$  and  $Q = \langle 0; +\infty \rangle$ .

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