

Ireneusz Nabiałek

BASES OF AGGREGABLE SETS

Introduction

The notions of a computable set and a computable function are introduced in [1]. The notions of an aggregable set and a basis of an aggregable set are introduced in this paper analogously to [2]. The aim of this paper is to evolve a theory of bases of aggregable sets.

1. Aggregable sets

Let $\zeta = (T, \leq, \Theta, X, Y, Z)$ be such a system which is introduced in [1]. By \mathcal{F}_ζ we denote the set of all computable functions $f : T \times X \rightarrow Y$.

Definition 1.1. Functions $f, g \in \mathcal{F}_\zeta$ are said to be commonly computable (in symbols $f \sim g$) if there exists a computable set $F \subseteq \mathcal{F}_\zeta$ such that $f, g \in F$.

Corollary 1.1. The relation \sim is reflexive and symmetric in \mathcal{F}_ζ .

Definition 1.2. The set $F \subseteq \mathcal{F}_\zeta$ is called aggregable iff the relation \sim is transitive in F and the set F is $*$ -closed (see [1]).

Corollary 1.2. Any computable set is aggregable.

Corollary 1.3. For any aggregable set F and for any set $H \subseteq F$ the set H^* (see [1]) is aggregable.

Definition 1.2 and Corollary 1.1 imply the following corollary.

Corollary 1.4. The relation \sim is an equivalence relation in any aggregable set.

Theorem 1.1. For any aggregable set F and for any $f \in F$ the equivalence class $[f]_\sim$ is a computable set.

Proof. The proof is analogical to the proof of Theorem 2 in [2].

2. *-connected computable sets

Definition 2.1. For any $f, g \in \mathcal{F}_\tau$ we define a relation \equiv as follows (see [1])

$$(2.1) \quad (f \equiv g) \Leftrightarrow \exists_{a, b \in T} (f_a = g_b).$$

Theorem 2.1. The relation \equiv is an equivalence relation in \mathcal{F}_τ .

Proof. It is evident that the relation \equiv is reflexive and symmetric in \mathcal{F}_τ . If $f \equiv g$ and $g \equiv h$, then there exists $a, b, c, d \in T$ such that $f_a = g_b$ and $g_c = h_d$. Let $b \leq c$ and $\theta(b) = \beta$, $\theta(c) = \gamma$ (see [1]). Let $e = \beta^{-1}(c)$ and $\theta(e) = \eta$. If $f_a = g_b$, then $(f_a)_e = (g_b)_e = g \circ \tilde{\beta} \circ \tilde{\eta} = g \circ \tilde{\gamma} = g_c$, thus $f'_a = g_c$, where $f'_a = (f_a)_e$. So, if $g_c = h_d$, then $f'_a = h_d$, whence $f \equiv h$. The relation \equiv is transitive in \mathcal{F}_τ . If $c \leq b$, then the proof is analogical.

Corollary 2.1. The relation \equiv is an equivalence relation in any aggregable set.

Theorem 2.2. For any aggregable set F and for any $f, g \in F$, if $f \equiv g$, then $f \sim g$.

Proof. The proof is analogical to the proof of Theorem 4 in [2].

Corollary 2.2. For any aggregable set F and for any function $f \in F$, $[f]_{\equiv} \subseteq [f]_{\sim}$.

Theorem 2.3. For any aggregable set F and for any $f \in F$ the equivalence class $[f]_{\equiv}$ is a computable set.

Proof. The set $[f]_{\equiv}$ is \mathbb{Z} -injective, because $[f]_{\equiv} \subseteq [f]_{\sim}$ and $[f]_{\sim}$ is \mathbb{Z} -injective. If $g \in ([f]_{\equiv})^*$, then by Theorem 5.3 [1] there exists $h \in [f]_{\equiv}$ such that $g \in \{h\}^*$. Hence $g \equiv h \equiv f$ and $g \equiv f$. Since $g \in [f]_{\equiv}$, we have $([f]_{\equiv})^* \subseteq [f]_{\equiv}$.

Since $[f]_{\#} \subseteq ([f]_{\#})^*$ (see Corollary 5.1 in [1]), we infer that $([f]_{\#})^* = [f]_{\#}$, whence the set $[f]_{\#}$ is $*$ -closed.

Definition 2.2. A set $F \subseteq \mathcal{F}_\gamma$ is called $*$ -connected iff

$$(2.2) \quad \bigvee_{f, g \in F} (f \# g).$$

By Corollaries 1.2, 2.2 and Definition 2.2 it is possible to write Theorem 2.3 in the following form.

Theorem 2.4. Any aggregable set and any computable set can be represented as a sum of $*$ -connected computable sets.

Theorem 2.5. Let F be an aggregable set, $G \subseteq F$ and $f \in G$. The set G is computable and $*$ -connected iff $G^* = G$ and $\{f\}^* \subseteq G \subseteq [f]_{\#}$.

Proof. The proof is analogical to the proof of Theorem 6 in [2].

Corollary 2.3. For any aggregable set F and for any function $f \in F$ the set $\{f\}^*$ is the minimal set, and the set $[f]_{\#}$ is the maximal set of all computable and $*$ -connected sets G such that $G \subseteq F$ and $f \in G$.

3. Basis of aggregable set

Definition 3.1. A set G is called a basis of an aggregable set F iff $F = \bigcup_{f \in G} [f]_{\#}$ and

$$(3.1) \quad \bigvee_{f, g \in G} \{ (f \neq g) \Rightarrow [\sim(f = g)] \}.$$

Corollary 3.1. If a set G is a basis of an aggregable set F , then $G \subseteq F$.

Let $F/\#$ be the quotient space of the equivalence relation $\#$.

Theorem 3.1. Let F be an aggregable set and $G \subseteq F$. The set G is a basis of F iff it has exactly one element in common with every set $[f]_{\#} \in F/\#$.

Proof. The proof is analogical to the proof of Theorem 7 in [2].

Theorem 3.2. If G is a basis of an aggregable set F , then for any set $H \subseteq G$ the set H^* is aggregable and H is a basis of H^* .

Proof. If $H \subseteq G$, then $H^* \subseteq G^*$ and since $G \subseteq F$ and $F^* = F$, we have $G^* \subseteq F$ and $H^* \subseteq F$. By Corollary 1.3 the set H^* is aggregable. In H the condition 3.1 is satisfied. If $g \in H^*$, then by Theorem 5.3 in [1] there exists $f \in H$ such that $g \in \{f\}^*$. If $h \in H^*$ and $h \dot{=} g$, then $h \in \{f\}^*$. If $h \notin \{f\}^*$, then there exists $\varphi \in H$ such that $h \in \{\varphi\}^*$ and $\varphi = f$, and by (3.1), $\sim(\varphi \dot{=} f)$, thus $\sim(h \dot{=} g)$. Thus if $[f]_{\dot{=}} \in H^*/_{\dot{=}}$, then $[f]_{\dot{=}} \subseteq \{f\}^*$. By Theorem 2.5 we have $\{f\}^* \subseteq [f]_{\dot{=}}$, thus $\{f\}^* = [f]_{\dot{=}}$. Hence $H^* = \bigcup_{f \in H} [f]_{\dot{=}}$ by Theorem 5.3 in [1].

Theorem 3.3. If G is a basis of an aggregable set F and $H_1 \subsetneq H_2 \subseteq G$, then $H_1^* \subsetneq H_2^*$.

Proof. If $H_1 \subsetneq H_2$, then $H_1^* \subseteq H_2^*$, thus it is sufficient to demonstrate that if $H_1 \subsetneq H_2 \subseteq G$, then $H_1^* \neq H_2^*$. Really if $H_1 \subsetneq H_2$, then there exists $f \in H_2$ such that $f \notin H_1$. By (3.1) for any function $g \in H_1$, we have $\sim(f \dot{=} g)$, thus $f \notin \{g\}^*$, whence $f \notin H_1^*$. Because $f \in H_2$ and $f \notin H_1^*$, hence $H_1^* \neq H_2^*$.

4. Free basis

Definition 4.1. Bases G_1, G_2 of an aggregable set F are called equivalent (in symbols $G_1 \approx G_2$) if $G_1^* = G_2^*$.

Corollary 4.1. Let G_F be the set of all bases of an aggregable set F . The relation \approx is an equivalence relation in G_F .

Definition 4.2. An equivalence class $[G]_{\approx} \in G_F/\approx$ is called a free basis of F and it is denoted by \tilde{G} .

Definition 4.3. Let $\tilde{G}_1, \tilde{G}_2 \in G_F/\approx$. We define a relation $<$ in G_F/\approx as follows

$$(4.1) \quad (\tilde{G}_1 < \tilde{G}_2) \iff (\tilde{G}_1^* \neq \tilde{G}_2^*).$$

Theorem 4.1. For any aggregable set F the relation $<$ is irreflexive and transitive in G_F/\approx .

Proof. If $\tilde{G}_1, \tilde{G}_2 \in G_F/\approx$ and $\tilde{G}_1 < \tilde{G}_2$, then $\tilde{G}_1^* \neq \tilde{G}_2^*$ and $(\tilde{G}_2^* \subseteq \tilde{G}_1^*)$, hence $\sim(\tilde{G}_2 < \tilde{G}_1)$. If $\tilde{G}_1, \tilde{G}_2, \tilde{G}_3 \in G_F/\approx$, $\tilde{G}_1 < \tilde{G}_2$ and $\tilde{G}_2 < \tilde{G}_3$, then $\tilde{G}_1^* \subseteq \tilde{G}_2^*$ and $\tilde{G}_2^* \subseteq \tilde{G}_3^*$, hence $\tilde{G}_1^* \subseteq \tilde{G}_3^*$. So we have $\tilde{G}_1 < \tilde{G}_3$.

Corollary 4.2. For any aggregable set F the relation \leq defined as follows

$$(4.2) \quad (G_1 \leq G_2) \iff (\tilde{G}_1 < \tilde{G}_2) \vee (\tilde{G}_1 = \tilde{G}_2)$$

is a partial order relation in G_F/\approx .

Theorem 4.2. A free basis \tilde{G}_0 of F is the last element in the partially ordered set $(G_F/\approx, \leq)$ if and only if $G_0^* = F$.

Proof. Let \tilde{G}_0 be the last element of $(G_F/\approx, \leq)$ and $G_0^* \neq F$. Since G_0 is a basis of F , we have $G_0^* \subseteq F$ and if $G_0^* \neq F$, then $G_0^* \subseteq F$. Hence there exists $f \in F$ such that $f \notin G_0^*$. By Definition 1.1 there exists exactly one function $\varphi \in G_0$ such that $\varphi^* = f$. Let H_0 denote the set $(G_0 - \{\varphi\}) \cup \{f\}$. H_0 is clearly a basis of F . We have $G_0 = H \cup \{\varphi\}$ and $H_0 = H \cup \{f\}$, where $H = G_0 - \{\varphi\}$. Whence $G_0^* = H^* \cup \{\varphi\}^*$ and $H_0^* = H^* \cup \{f\}^*$ (see [1]). Because for every $\tilde{G} \in G_F/\approx$ we have $G^* \subseteq G_0^*$, we infer that $H_0^* \subseteq G_0^*$, hence $\{f\}^* \subseteq \{\varphi\}^*$ and $f \in \{\varphi\}^*$. Since $\{\varphi\}^* \subseteq G_0^*$, we have $f \in G_0^*$. Hence if \tilde{G}_0 is the last element of $(G_F/\approx, \leq)$, then $G_0^* = F$. If $G_0^* = F$, then for any $\tilde{G} \in G_F/\approx$, we have $G^* \subseteq G_0^*$, because $G^* \subseteq F$.

Definition 4.4. The last element \tilde{G}_0 in $(G_F/\sim, \leq)$ is called a proper basis of F .

Corollary 4.3. A free basis \tilde{G}_0 is a proper basis of F iff $G_0^* = F$.

Theorem 3.2 implies the following corollary.

Corollary 4.4. If G is a basis of F , then for any set $H \leq G$, H is a proper basis of H^* .

Corollary 4.5. Any aggregable set has at most one proper basis.

5. Free functions

By Definition 3.1 and the condition $\{f\}^* \subseteq [f]_{\sim}$ we have the following corollary.

Corollary 5.1. For any function $g \in \{f\}^*$ the set $\{g\}$ is a basis of $\{f\}^*$.

Definition 5.1. Any free basis $\{\tilde{g}\}$ of $\{f\}^*$ is called a free function of $\{f\}^*$.

We denote by \tilde{g} the free function $\{\tilde{g}\}$ of $\{f\}^*$ and by f/\sim we denote the set of all free functions of $\{f\}^*$.

Theorem 5.1. If a function $f \in \mathcal{F}_Z$ is Z -injective then any free function $\tilde{g} \in F/\sim$ has exactly one representation g .

Proof. If $\tilde{g} \in f/\sim$, then $g \in \{f\}^*$ and $\{g\}^* \subseteq \{f\}^*$. If $h \in \tilde{g}$, then $\{h\}^* = \{g\}^*$ (see Definitions 5.1 and 4.2). If $\{h\}^* = \{g\}^* \subseteq \{f\}^*$, then $g, h \in \{f\}^*$ and if the function f is Z -injective, then the functions g, h are Z -injective, too (see [1]). If $\{h\}^* = \{g\}^*$, then $h \in \{g\}^*$ and $g \in \{h\}^*$. If $h \in \{g\}^*$, then there exists $a \in T$ such that $h = g_a$, and if $g \in \{h\}^*$, then there exists $b \in T$ such that $g = h_b$. Hence $h = (h_b)_a$. Since the function $h = g_a$ is Z -injective, we obtain $a = b = 0$ and $h = g$.

Theorem 5.2. If a free function $\tilde{g} \in f/\sim$ has more than one representation, then the computable function f is Z -periodic (see [1]).

Proof. By Theorem 5.1 if there exists a free function $\tilde{g} \in f/\sim$ such that \tilde{g} has more than one representation, then

f is not Z -injective. Because f is the computable function, thus f is Z -periodic (see Theorem 6.2 in [1]).

Corollary 5.2. If a free function \tilde{g} has more than one representation, then g is Z -periodic.

Theorem 5.3. For any function $f \in \mathcal{F}_Z$ the relation \leq defined by (4.2) is an ordering relation in f/\sim .

Proof. By Corollary 4.2 the relation \leq is a partially ordering relation in f/\sim . It is enough to prove that the relation \leq is connective in f/\sim . If $\tilde{g}, \tilde{h} \in f/\sim$ then by Definition 5.1 we have $g, h \in \{f\}^*$ and by Theorem 6.4 from [1], $g \in \{h\}^*$ or $h \in \{g\}^*$. If $g \in \{h\}^*$, then $\{g\}^* \subseteq \{h\}^*$, thus $\tilde{g} \leq \tilde{h}$. Analogously, if $h \in \{g\}^*$, then $\tilde{h} \leq \tilde{g}$.

Theorem 5.4. For any $f \in \mathcal{F}_Z$ the free function \tilde{f} is the last element in $(f/\sim, \leq)$.

Proof. It is enough to demonstrate that for any $\tilde{g} \in f/\sim$, $\tilde{g} \leq \tilde{f}$. Because $\tilde{g} \in f/\sim$, thus $g \in \{f\}^*$, hence $\{g\}^* \subseteq \{f\}^*$ and $\tilde{g} \leq \tilde{f}$.

Theorem 5.5. Let F be an aggregable set and $f \in F$. If the set $[f]_{\sim}$ has a proper basis \tilde{G} , then \tilde{G} is a free function.

Proof. Any basis of $[f]_{\sim}$ is a one-element set $\{g\}$. If the set $[f]_{\sim}$ has a proper basis, then there exists $g \in [f]_{\sim}$ such that $\{g\}^* = [f]_{\sim}$ and the free function \tilde{g} is a proper basis of $[f]_{\sim}$.

Theorem 5.6. An aggregable set F has a proper basis \tilde{G} iff for every $f \in F$, the set $[f]_{\sim}$ has a proper basis \tilde{g} .

Proof. If \tilde{G} is a proper basis of F , then for every $f \in F$ there exists exactly one function $g \in \tilde{G}$ such that $g \in [f]_{\sim}$. Let $G_1 = \tilde{G} - \{g\}$, thus $\tilde{G} = G_1 \cup \{g\}$ and because \tilde{G} is a proper basis of F , then $G_1^* \cup \{g\}^* = F$. Because $F = \bigcup_{h \in G_1} [h]_{\sim} \cup [g]_{\sim}$, hence

$$(5.1) \quad G_1^* \cup \{g\}^* = \bigcup_{h \in G_1} [h]_{\equiv} \cup [g]_{\equiv}.$$

Because the set $[g]_{\equiv}$ is separate in G_1^* and $\bigcup_{h \in G_1} [h]_{\equiv}$, we have

$$(5.2) \quad \{g\}^* \cap [g]_{\equiv} = [g]_{\equiv}.$$

By (5.2) we have $[g]_{\equiv} \subseteq \{g\}^*$, and because $\{g\}^* \subseteq [g]_{\equiv}$ (see Theorem 2.5) we obtain $\{g\}^* = [g]_{\equiv}$. Hence the free function \tilde{g} is a proper basis of $[g]_{\equiv}$. Since $g \in [f]_{\equiv}$, we have $[g]_{\equiv} = [f]_{\equiv}$, and \tilde{g} is a proper basis of $[f]_{\equiv}$.

Corollary 5.3. Let F be an aggregable set and G be a basis of F . The set G is a representation of a proper basis \tilde{G} of F if and only if $F = \bigcup_{f \in G} \{f\}^*$ and if any only if $\tilde{G} = \bigcup_{f \in G} \tilde{f}$.

BIBLIOGRAPHY

- [1] I. N a b i a ł e k: Deterministic computability, Demonstratio Math. 9 (1976) 681-689.
- [2] I. N a b i a ł e k: Z-aggregable sets of functions, Demonstratio Math. 8 (1975) 491-495.

INSTITUTE OF MATHEMATICS, TECHNICAL UNIVERSITY OF WARSAW

Received February 14, 1976.